

Using context and model categories to define directed homotopies

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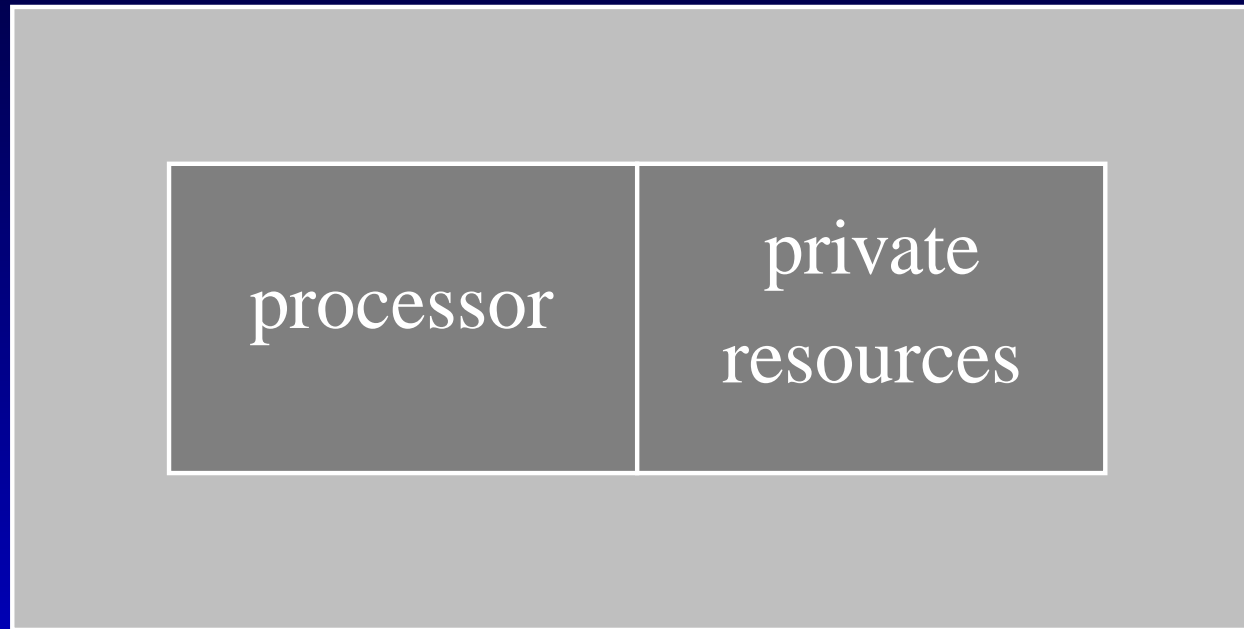
<http://igat.epfl.ch/bubenik/>

1. Introduction

- A concurrent system
- A model for concurrency
- Equivalences - undirected and directed
- The Problem

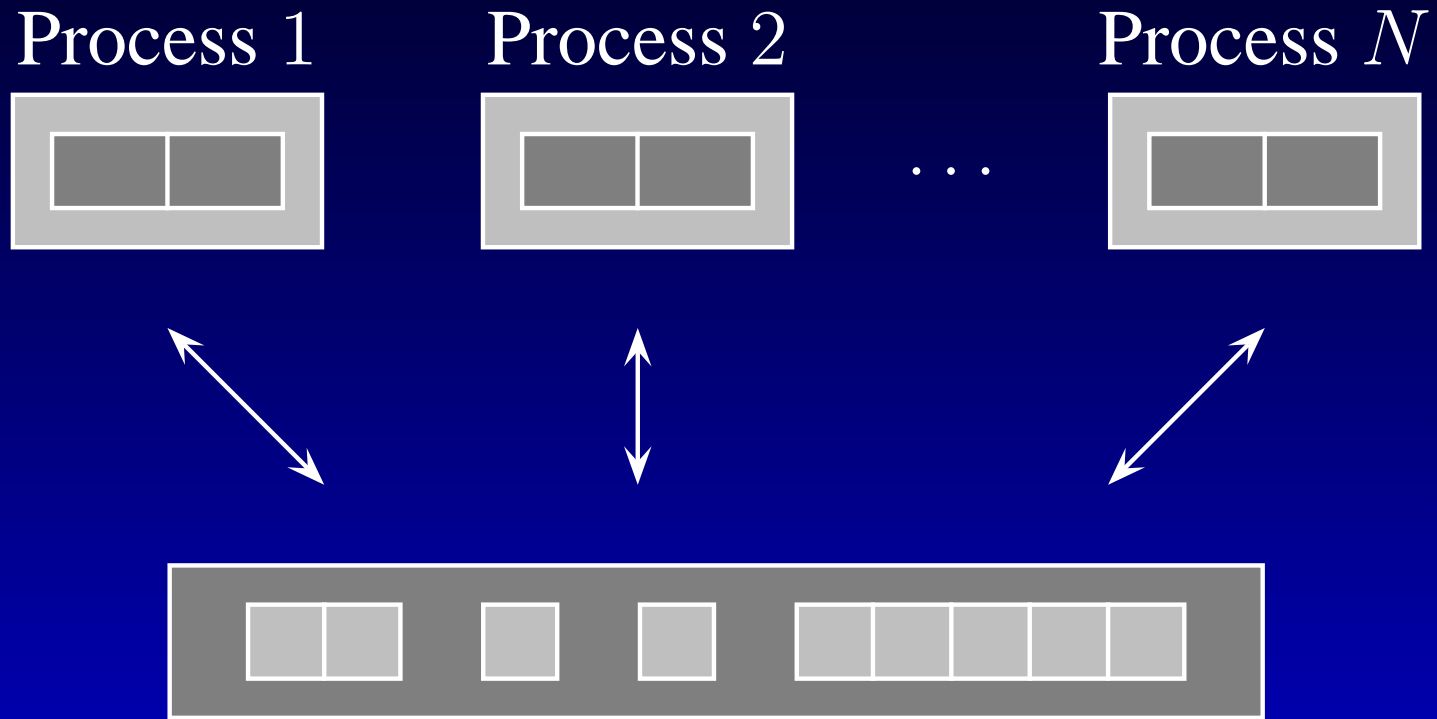
A nonconcurrent system

Process



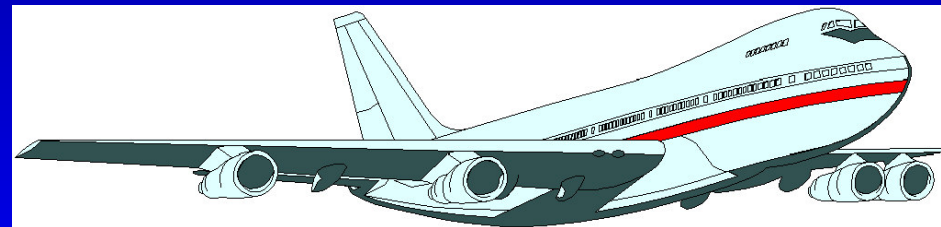
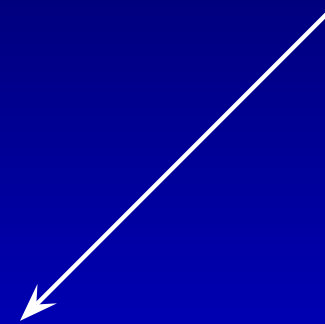
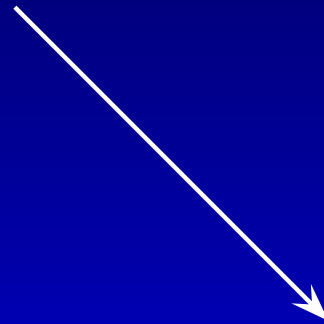
A process with its own private resources

A concurrent system



Several processes with shared resources

Example: Internet database





Example: dual core processors

Intel (Q3 2005)


AMD (Q3 2005)

digital ENTERPRISE
multi-core transition

	2005	2006	FUTURE
	EM64T Larger cache (2M) Faster FSB DDR2 Power Management (DBS) PCI-Express* I/O	DUAL CORE Lower Power Cores Enhanced Memory (FBD) Virtualization (VT) iAMT I/O Packet Acceleration Advanced Storage Controllers	MULTI-CORE 2 or more Advanced memory, virtualization, RAS and manageability
	DUAL CORE Multi-Threading Pellston Foxton	DUAL CORE Multi-Threading Pellston Foxton	MULTI-CORE 4 or more Advanced memory, virtualization, RAS and manageability

DUAL/MULTI-CORE RAMP
ALIGNED with PLATFORM FEATURES

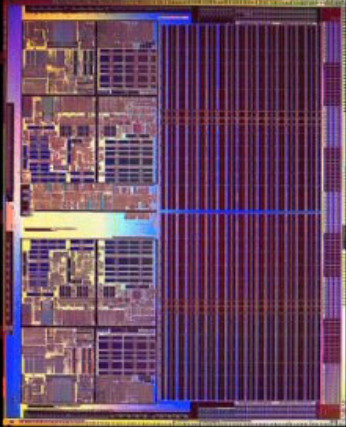
Source: Intel roadmap. *Other names and brands may be claimed as the property of others.

Mid-2005: AMD Opteron™ Processor Dual-Core 


90nm Process
Approximately same die size as 130nm single-core AMD Opteron™ processor*
~205 million transistors*

95 Watt Power Envelope
Fits into 90nm power infrastructure

940 Socket Compatible
All that's needed is a BIOS upgrade
Compatible with all motherboards designed to our 90nm specification



*Based on current revisions of the design
11/12/2004


2004 Analyst Day

A concurrent system

Example:

2 processes using 2 shared resources a and b which can only be used by one process at a time

Notation:

Px - a process locks resource x

Vx - a process releases resource x

Program of the first process: $Pa \quad Pb \quad Vb \quad Va$

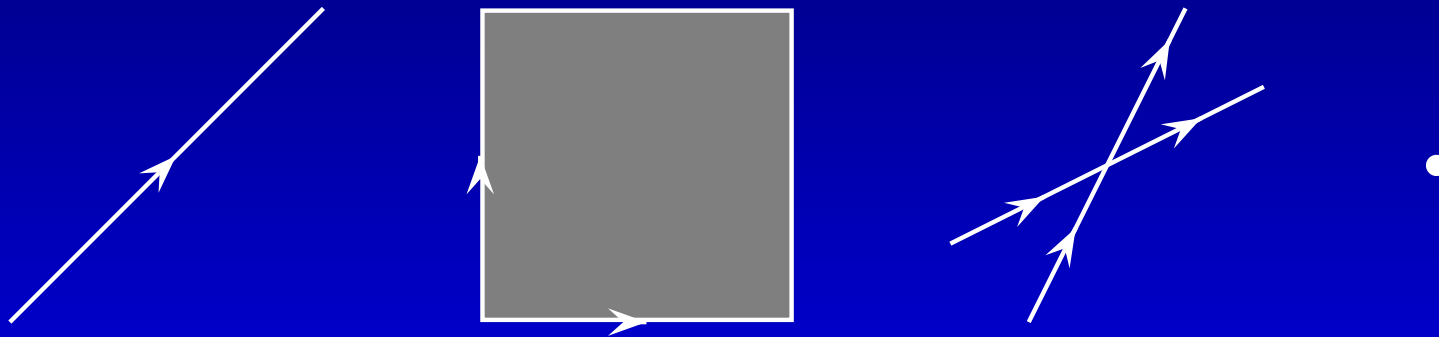
Program of the second process: $Pb \quad Pa \quad Va \quad Vb$

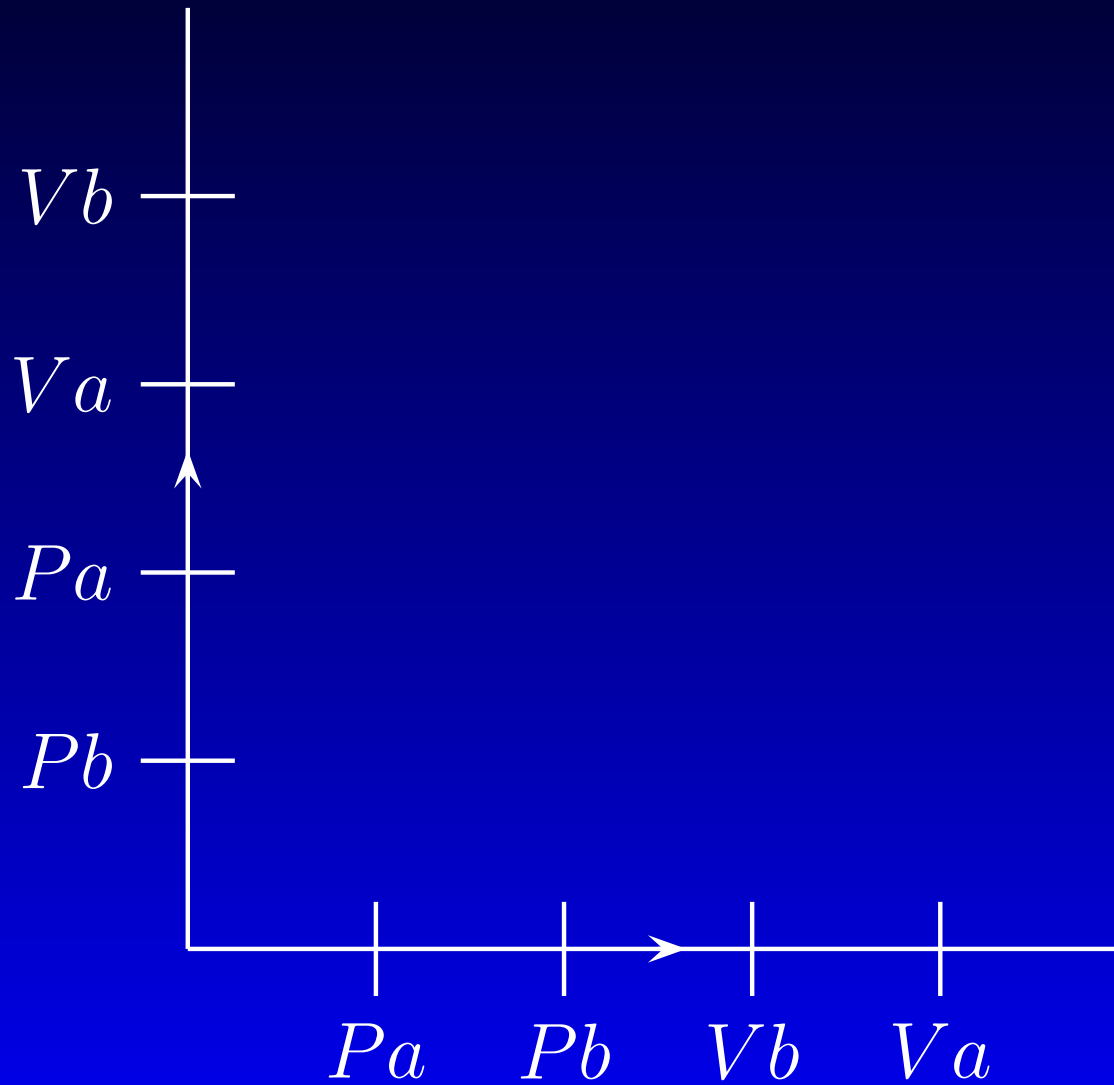
A model for concurrency

Concurrent systems can be modeled by subspaces of \mathbb{R}^n together with a partial order.

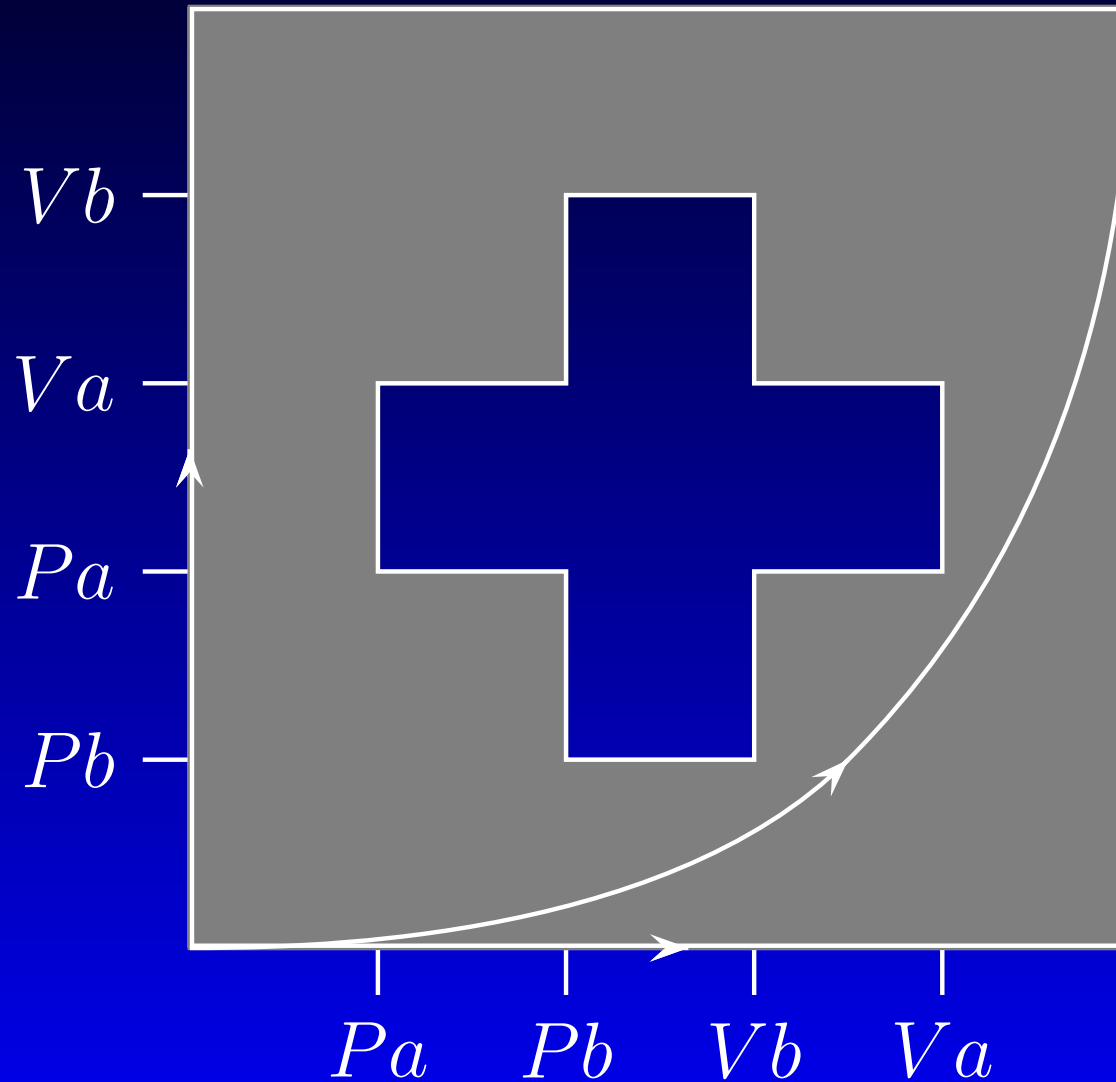
Definition: A **po-space** is a topological space U with a partial order \leq which is a closed subset of $U \times U$.

Examples:

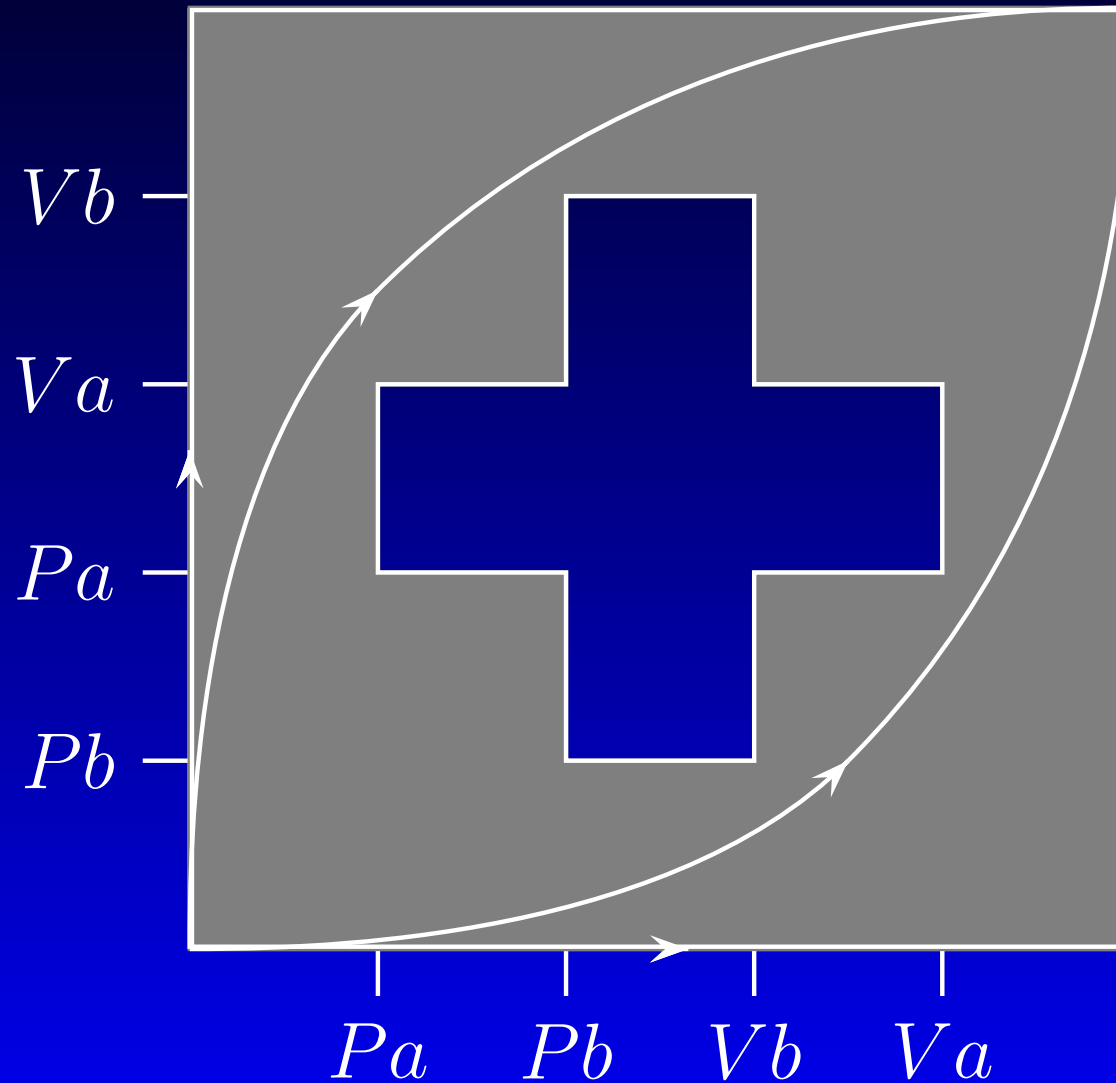




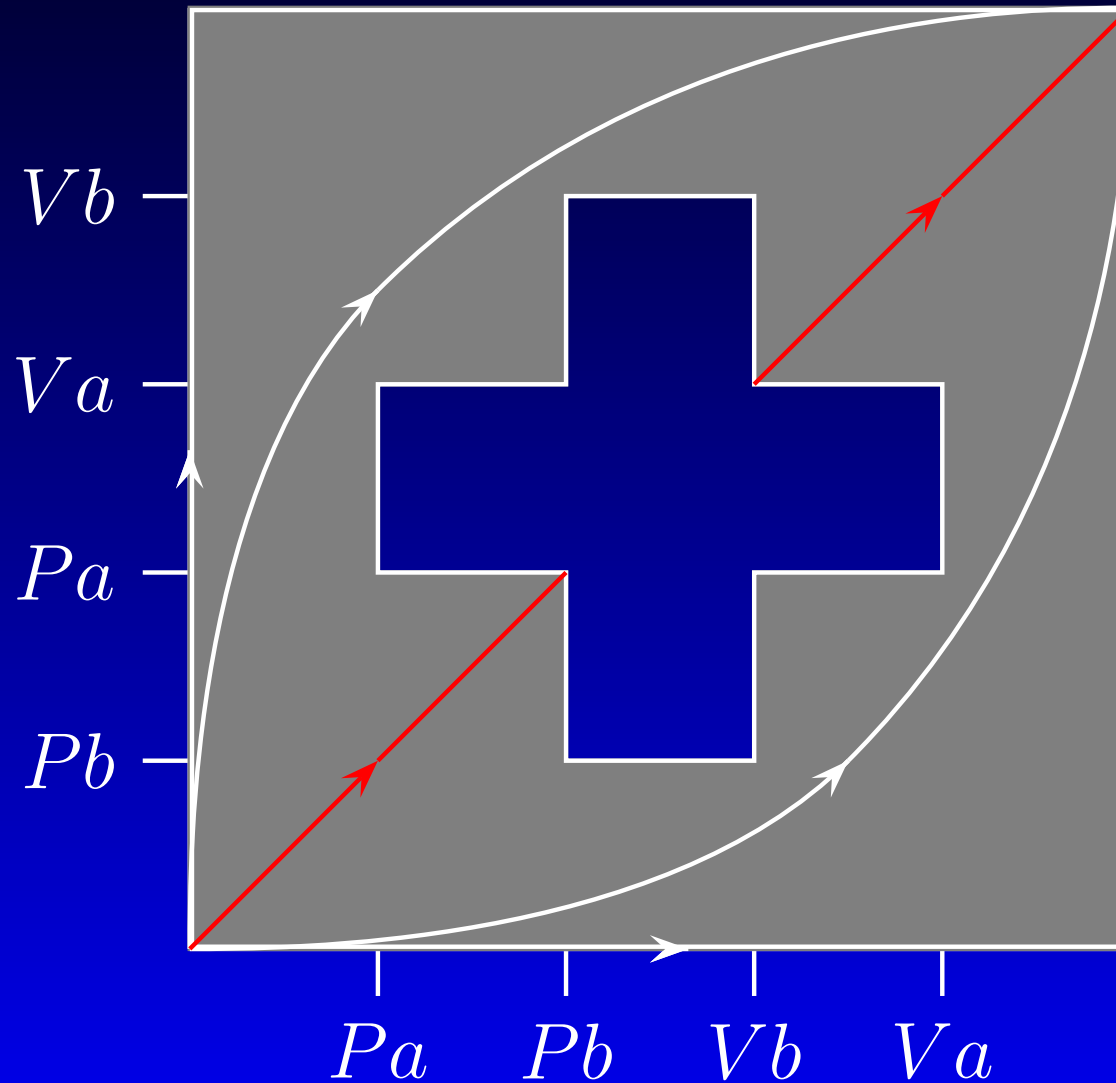
The Swiss flag



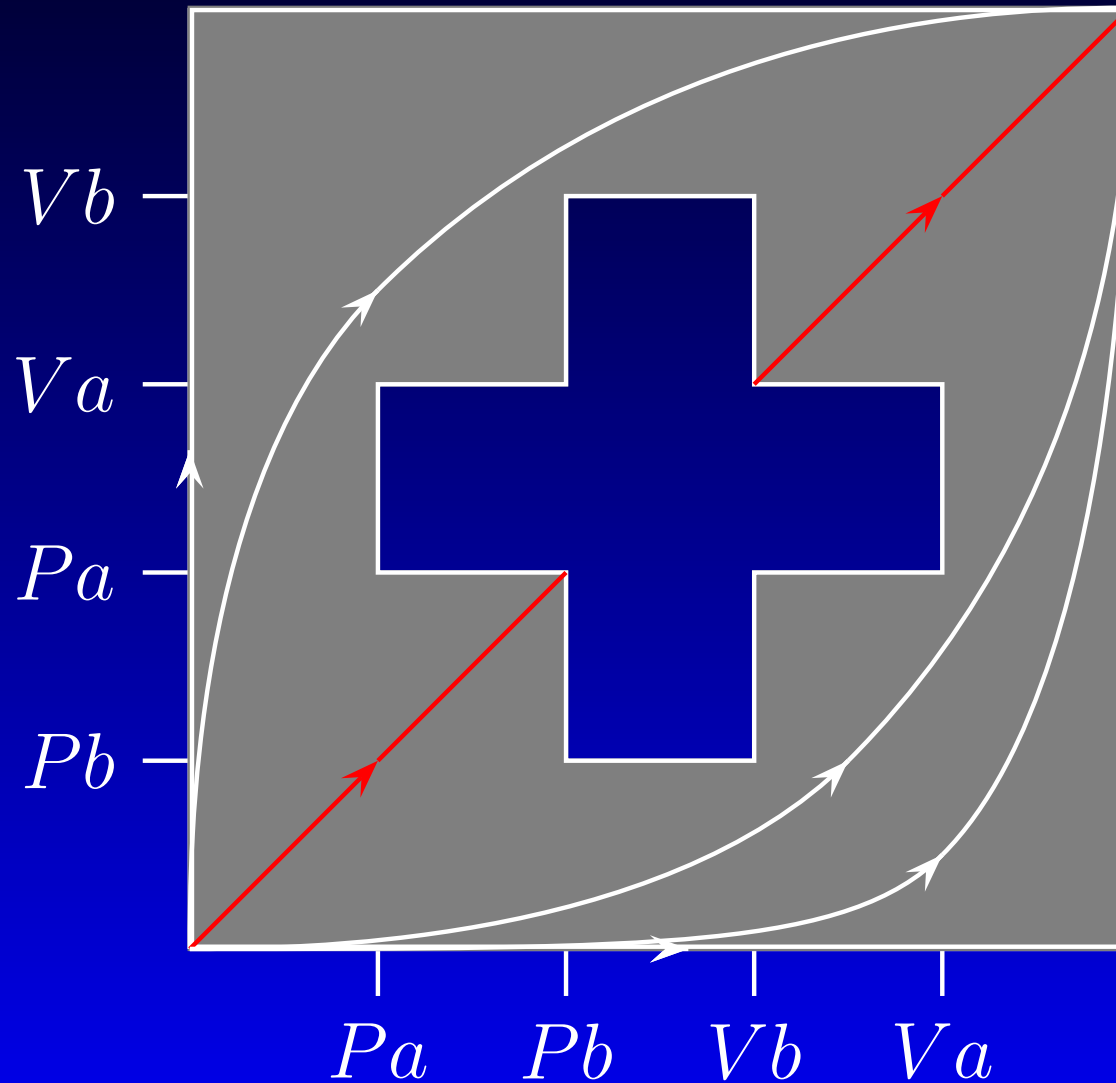
The Swiss flag



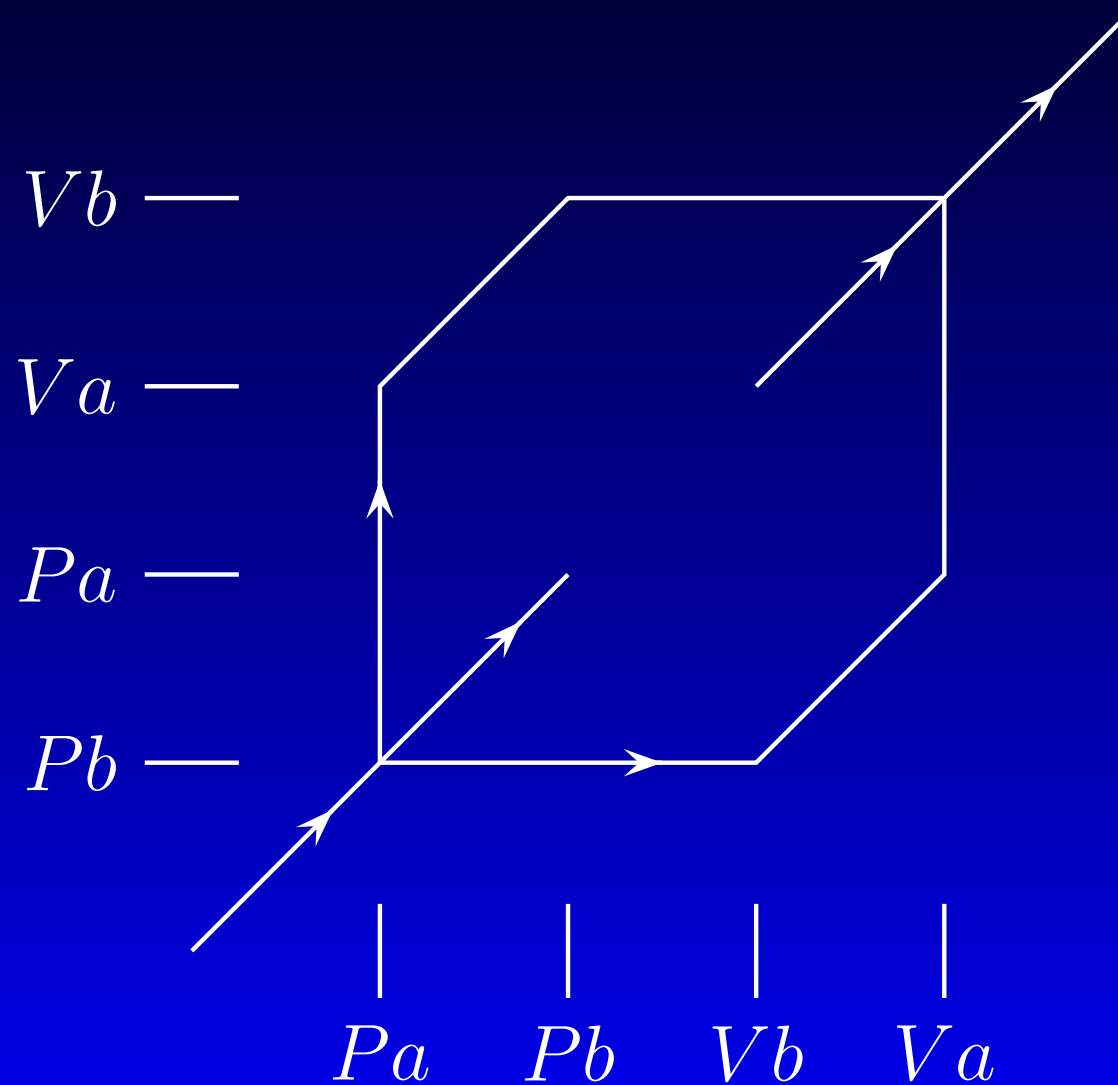
The Swiss flag



The Swiss flag



A sub-po-space of the Swiss flag



Goal

Develop a framework for concurrency where equivalences are accounted for.

We would like equivalences that allow a piece-by-piece analysis of po-spaces.

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We would like equivalences that allow a piece-by-piece analysis of po-spaces.

Idea: Use algebraic topology.

Undirected equivalences

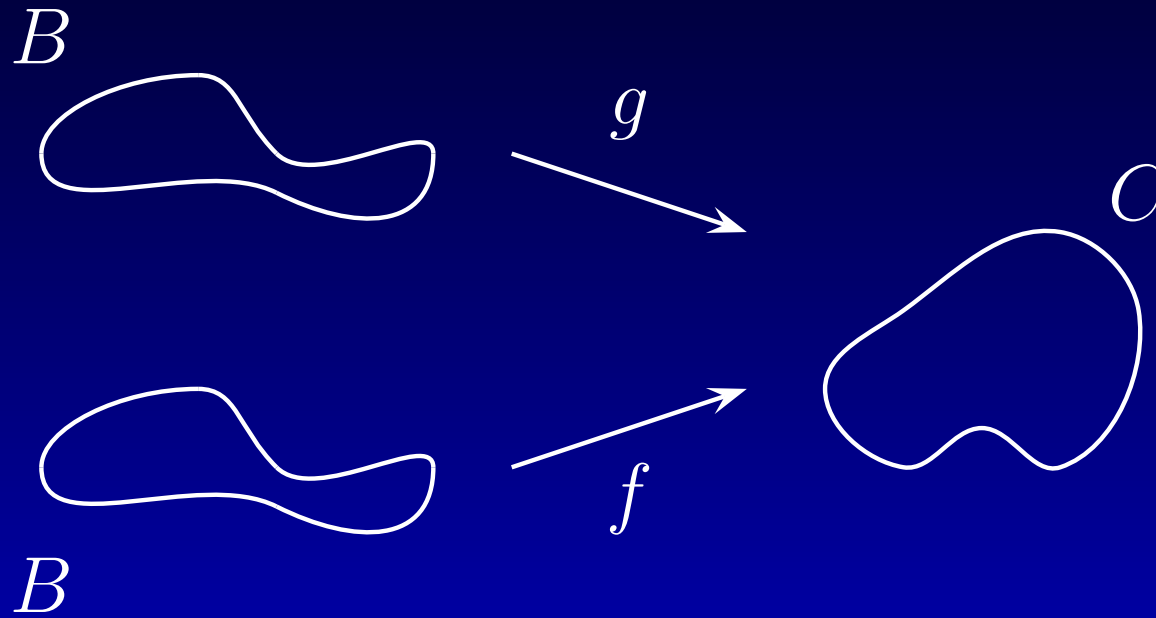
Working with (undirected) spaces and continuous maps, equivalences are defined using homotopies.

Maps are equivalent if there is a homotopy between them.

Spaces are equivalent if there maps between them whose compositions are homotopic to the identity map.

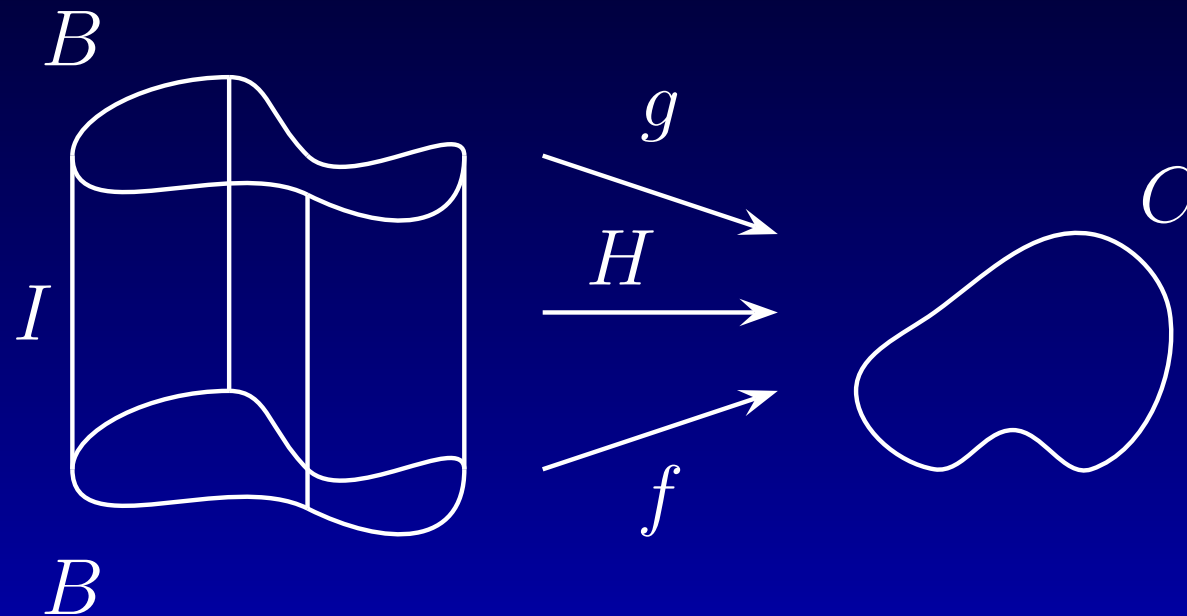
Undirected equivalences

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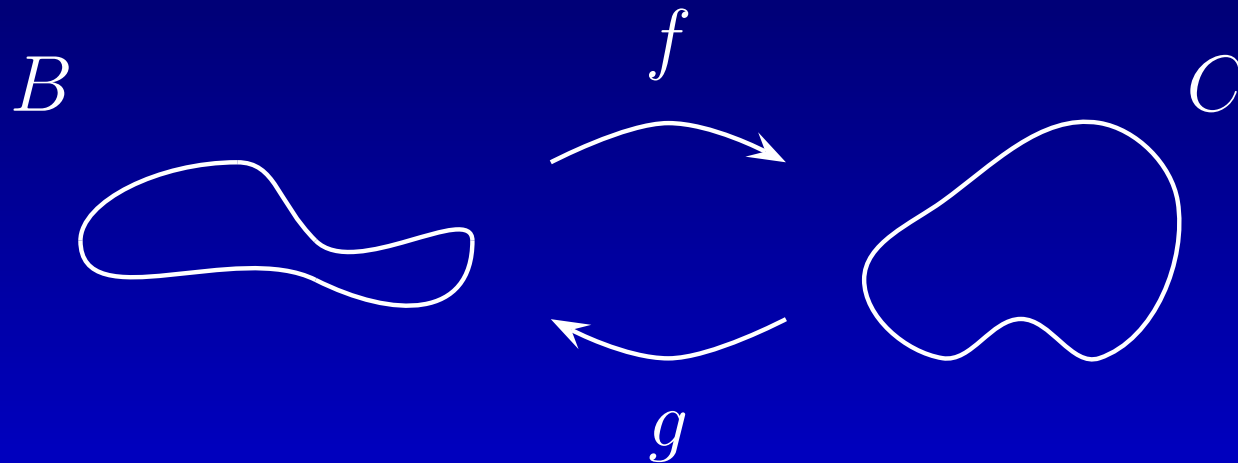


a **homotopy** between f and g is a continuous map $H : B \times I \rightarrow C$ restricting to f and g . This is an equivalence relation. Write $H : f \xrightarrow{\sim} g$.

Undirected equivalences

Definition: Spaces B, C are **homotopy equivalent** if there are maps $f : B \rightleftarrows C : g$ such that

$$g \circ f \simeq \text{Id}_B \text{ and } f \circ g \simeq \text{Id}_C .$$



Directed spaces

Definition:

- A **po-space** is a topological space U with a partial order \leq which is a closed subset of $U \times U$.
- A **directed map (dimap)** is a continuous map $f : U_1 \rightarrow U_2$ between po-spaces such that

$$x \leq y \implies f(x) \leq f(y).$$

- Let **PoSpc** be the category whose objects are po-spaces and morphisms are dimaps.

Directed spaces

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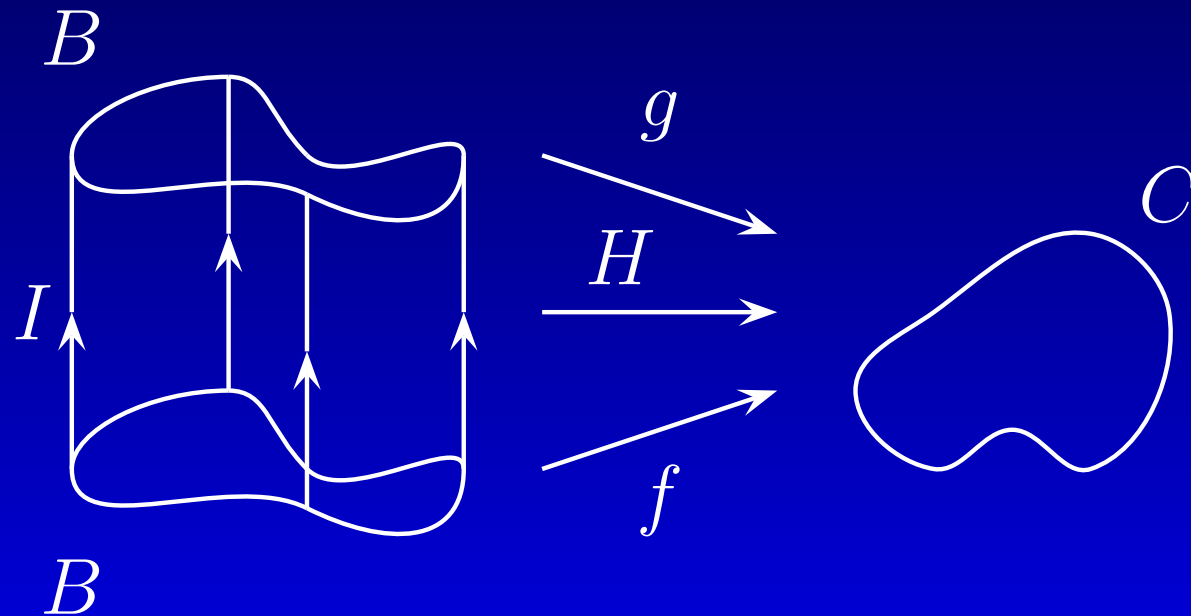
- Let **PoSpc** be the category whose objects are po-spaces and morphisms are dimaps.

Remark: Subspaces and products of po-spaces inherit a po-space structure.

Directed equivalences

Definition:

- A **directed homotopy (dihomotopy)** between dimaps $f, g : B \rightarrow C$ is a dimap $H : B \times \vec{I} \rightarrow C$ restricting to f and g . Write $H : f \rightarrow g$.



Directed equivalences

Definition:

- Write $f \simeq g$ if there is a chain of dihomotopies

$$f \rightarrow f_1 \leftarrow f_2 \rightarrow \dots \leftarrow f_n \rightarrow g.$$

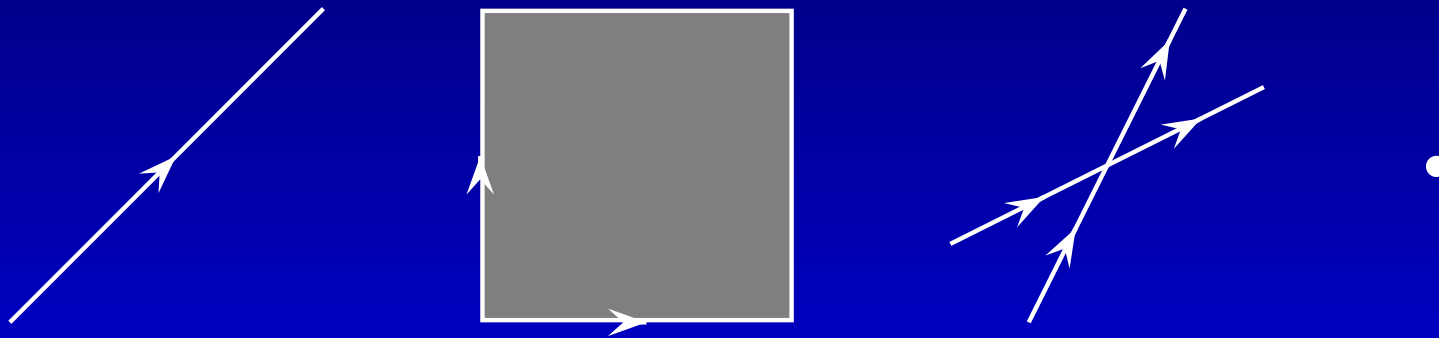
- Po-Spaces B, C are **dihomotopy equivalent** if there are dimaps $f : B \rightleftarrows C : g$ such that

$$g \circ f \simeq \text{Id}_B \text{ and } f \circ g \simeq \text{Id}_C .$$

The Problem

Recall: We wanted to use dihomotopy equivalences to provide equivalences of concurrent systems.

However all of the following spaces are dihomotopy equivalent.



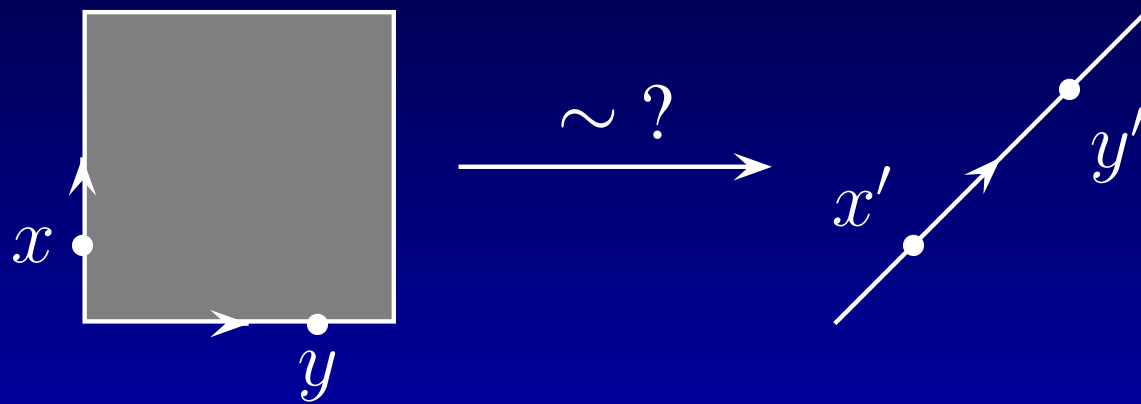
Thus, a stronger notion of equivalence is needed.

2. Using context

- A basic example
- A solution
- Equivalences using context
- Non-equivalences using directed paths
- Piece-by-piece example

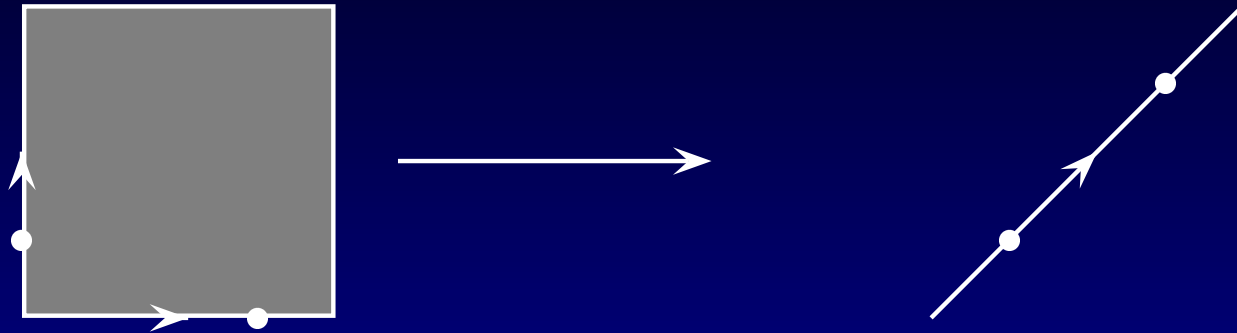
A basic example

Question: Is there an equivalence between $\vec{I} \times \vec{I}$ and \vec{I} ?

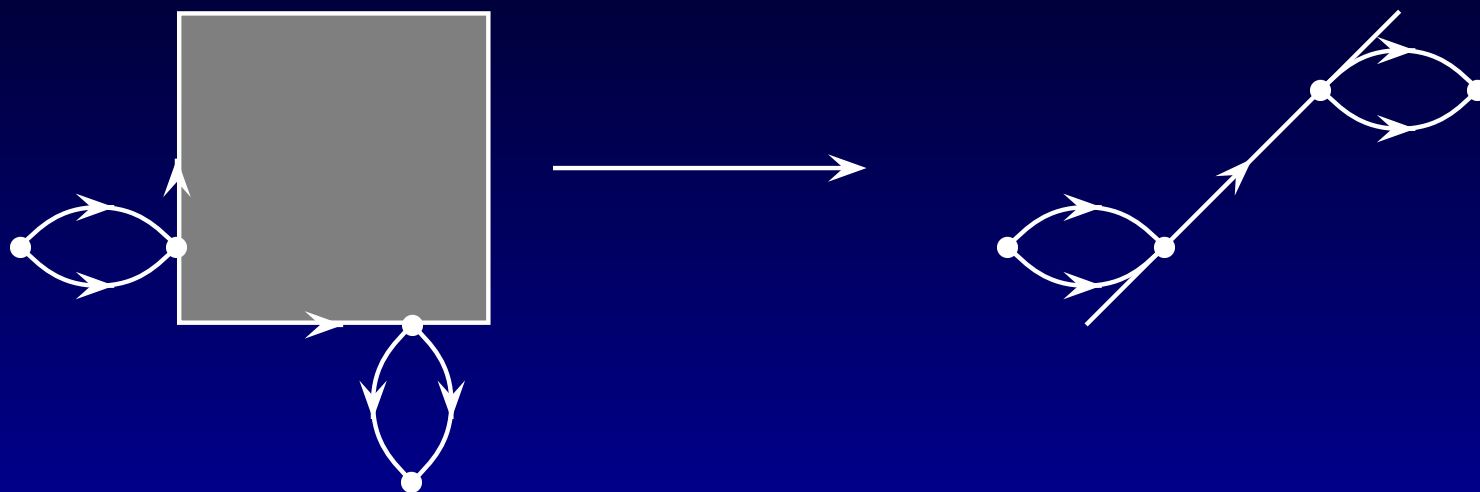


Our guide: If this map is an equivalence, then so is the map obtained by attaching po-spaces to this example. (Formally, we want the set of equivalences to be closed under pushouts with inclusions.)

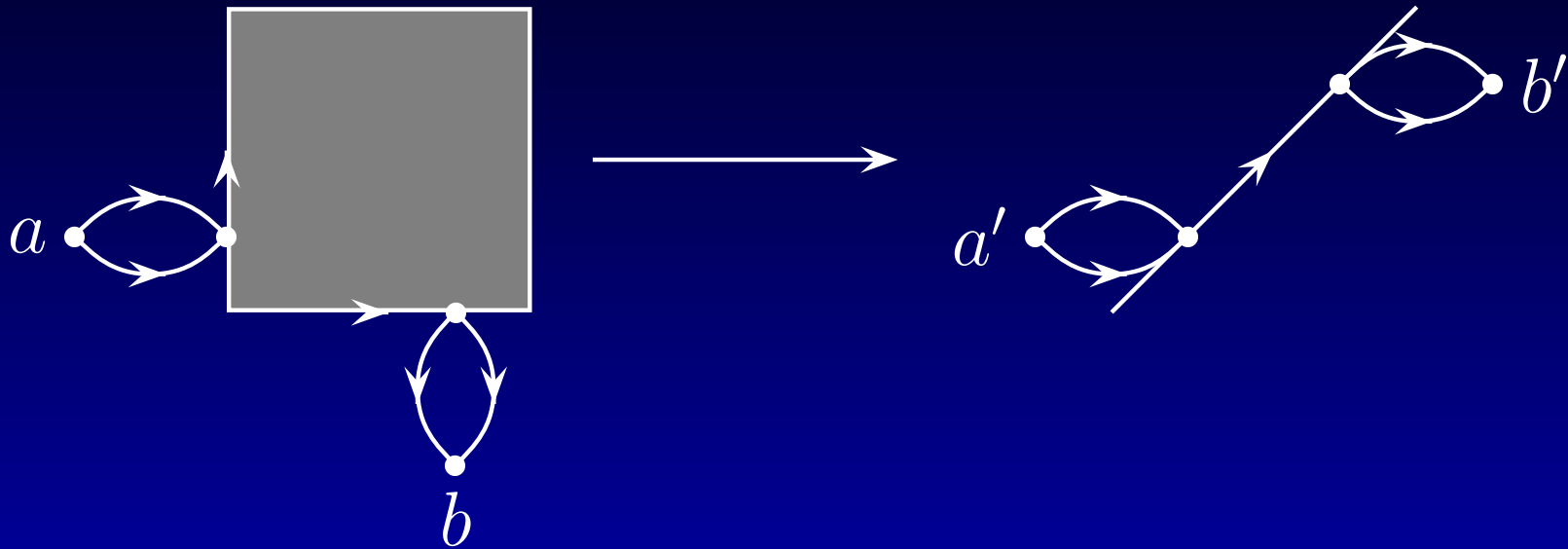
Adding to the basic example



Adding to the basic example



Adding to the basic example



There is no execution path from a to b , while there are such paths from a' to b' .
Thus, this map should not be an equivalence.

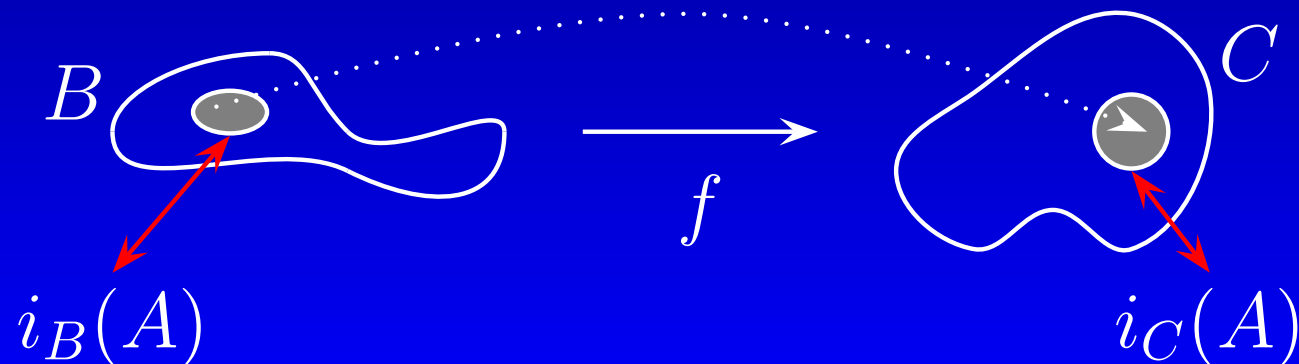
A Solution

Instead of working with just po-spaces work with po-spaces together with **context**.

Definition: Choose a po-space A (called the **context**). This choice will depend on the attachments one wants to consider.

Consider po-spaces B together with a dimap $i_B : A \rightarrow B$ and consider morphisms which are dimaps such that

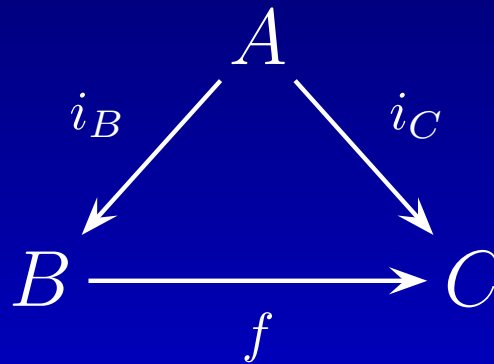
$$f(i_B(a)) = i_C(a) \text{ for all } a \in A$$



$A \downarrow \mathbf{PoSpc}$

Definition: Given a pospace A , let $A \downarrow \mathbf{PoSpc}$ be the category whose

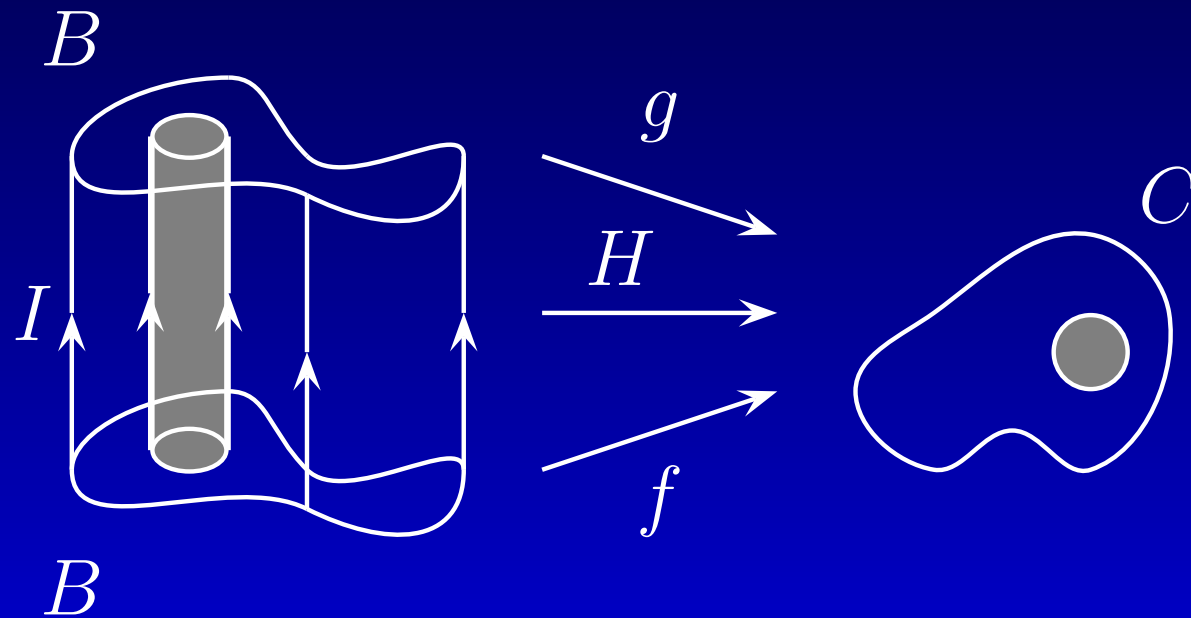
- objects are dimaps $\iota_B : A \rightarrow B$,
- morphisms are dimaps $f : B \rightarrow C$ such that $f \circ \iota_B = \iota_C$



Equivalences using context

Definition:

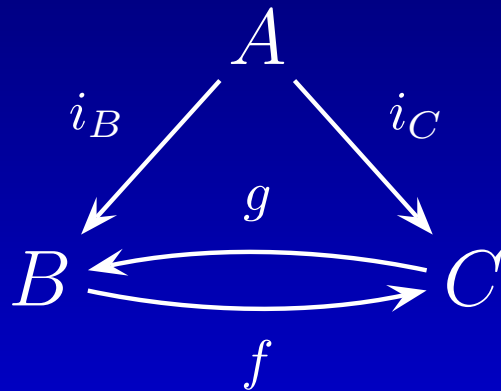
- A **dihomotopy** between $f, g : B \rightarrow C$ in the context of A is a dihomotopy $H : f \rightarrow g \text{ rel } A$.



Equivalences using context

Definition:

- Write $f \simeq g$ if there is a chain of dihomotopies $f \rightarrow f_1 \leftarrow f_2 \rightarrow \dots \leftarrow f_n \rightarrow g$.
- $i_B : A \rightarrow B, i_C : A \rightarrow C$ are **dihomotopy equivalent** if there are dimaps



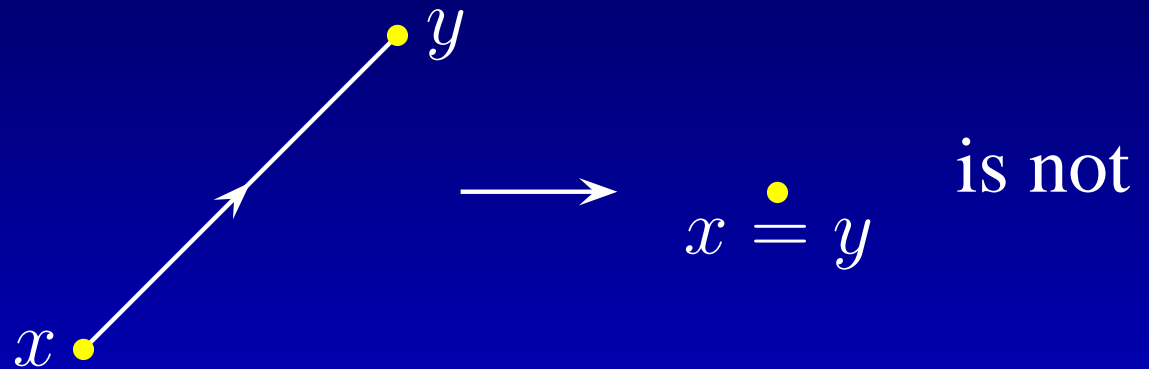
such that $g \circ f \simeq \text{Id}_B$
and $f \circ g \simeq \text{Id}_C$.

Example

Let $A = \{x, y\}$ with $x \leq y$.

Then po-spaces under the context A are just po-spaces with two marked points, one of which is after the other.

In this category

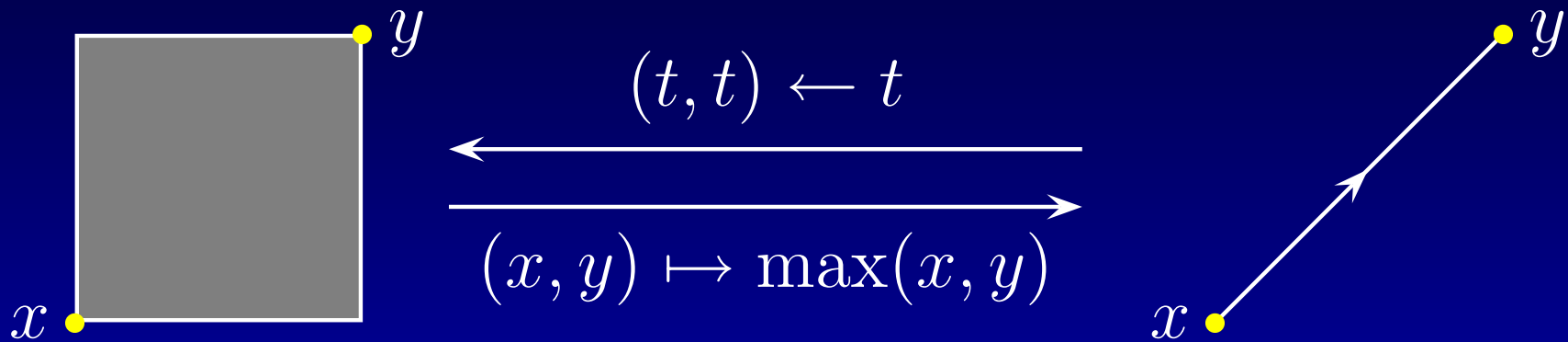


is not

a dihomotopy equivalence since there is no dimap in the reverse direction.

Example of an equivalence

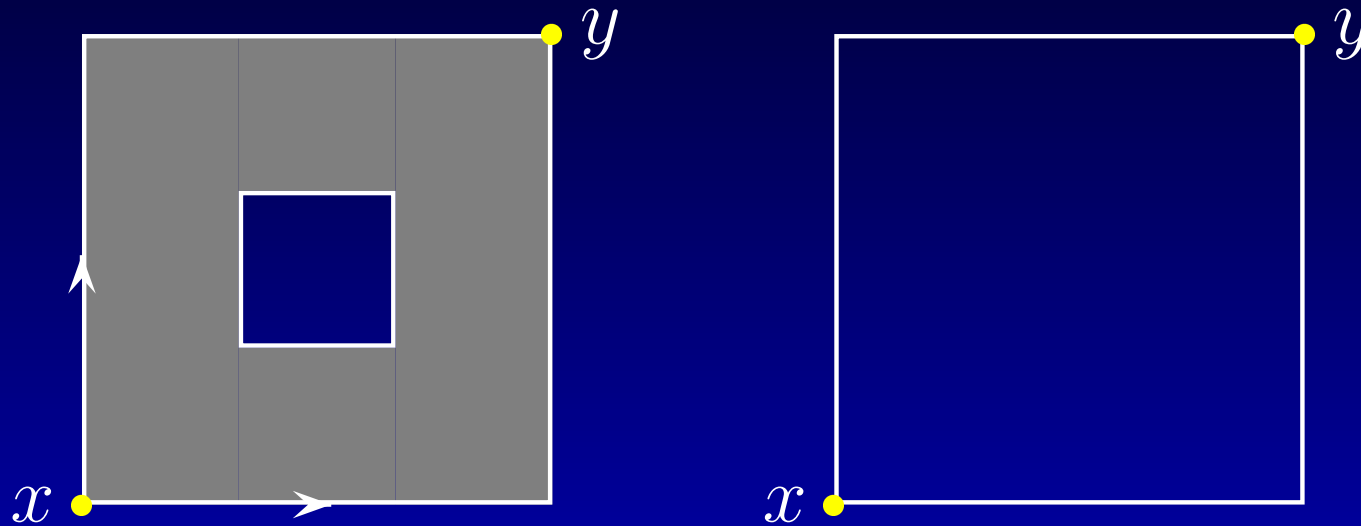
In the same context (of two marked points) the dimaps



give a dihomotopy equivalence.

Another equivalence

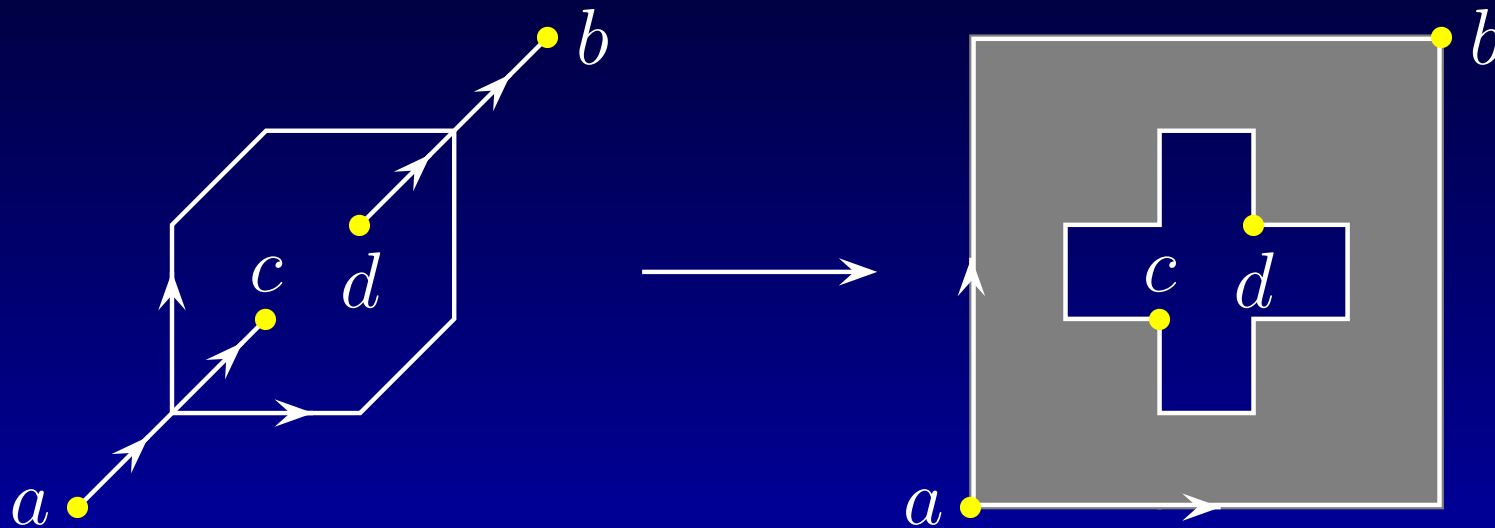
$\vec{I} \times \vec{I}$ with a square removed and two marked points



is dihomotopy equivalent to its boundary.

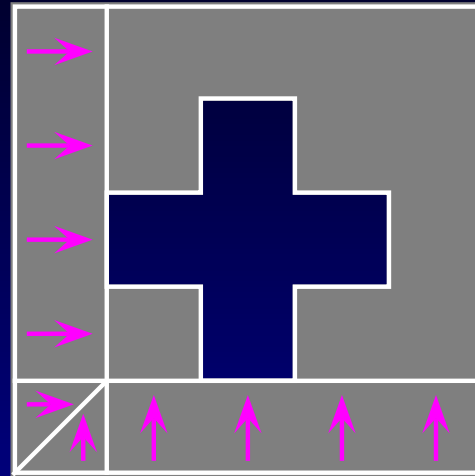
Context for the Swiss flag

Let $A = \{a, b, c, d\}$. Then the inclusion

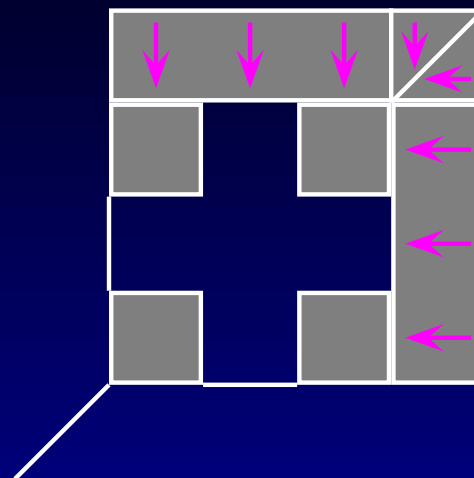
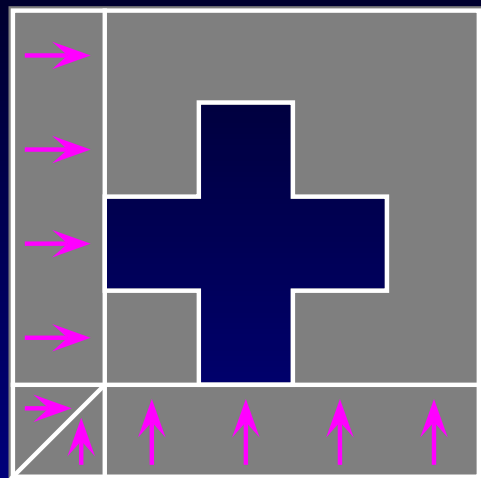


is a dihomotopy equivalence in the context of the four marked points.

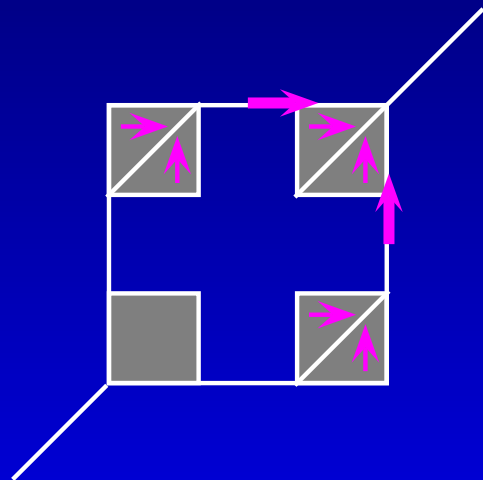
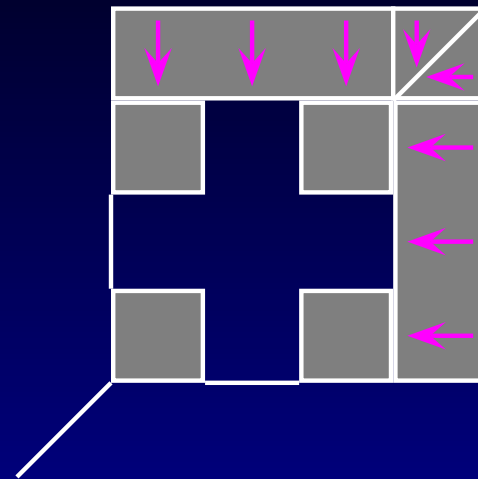
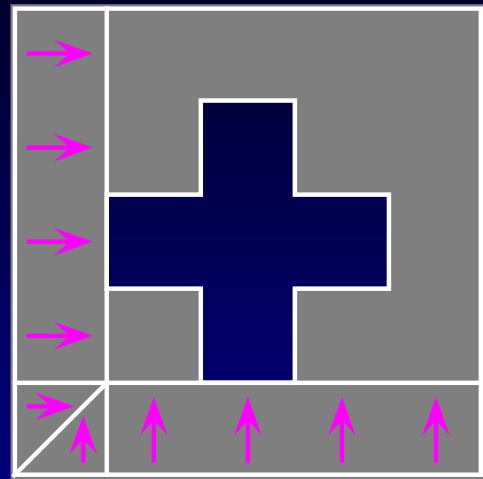
Sketch of the proof



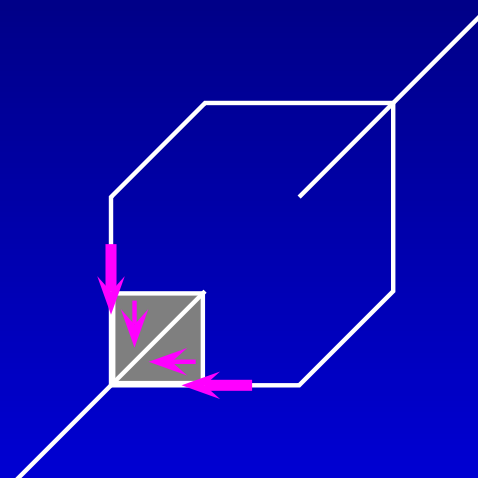
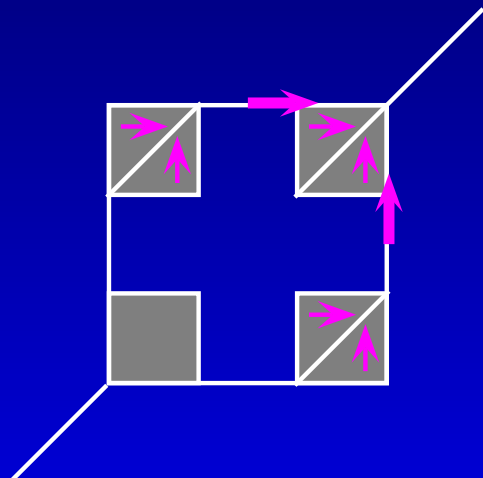
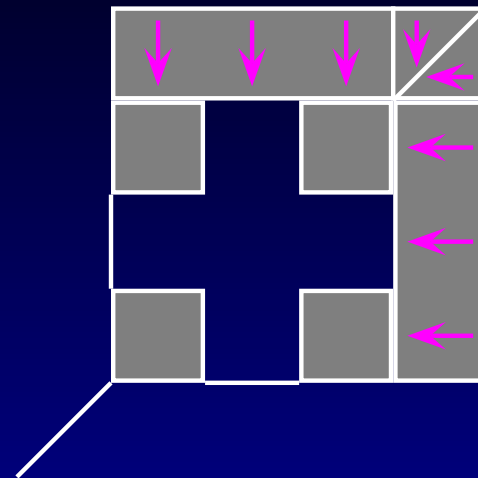
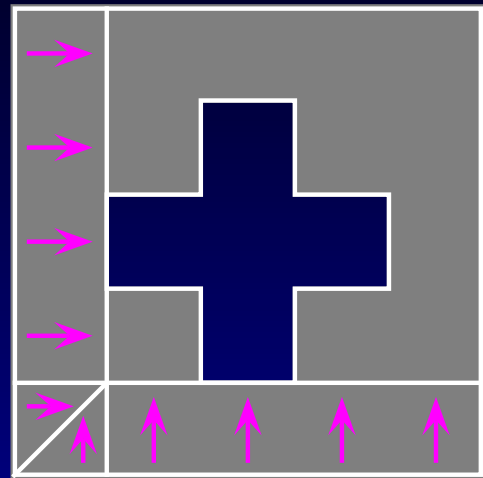
Sketch of the proof



Sketch of the proof



Sketch of the proof



Directed paths

Definition: Let $x, y \in$ the po-space B .

- A **dipath** is a dimap $\vec{I} \rightarrow B$.
- Dipaths are **dihomotopy equivalent** if they are so in the context of their endpoints.
- Let $\vec{\pi}_1(B)(x, y)$ be the set of dihomotopy equivalence classes of dipaths from x to y .

Dipaths in equivalent spaces

Notation: Given a context A and po-spaces B, C together with $i_B : A \rightarrow B$ and $i_C : A \rightarrow C$, if $x \in A$ denote $i_B(x)$ by x_B and $i_C(x)$ by x_C .

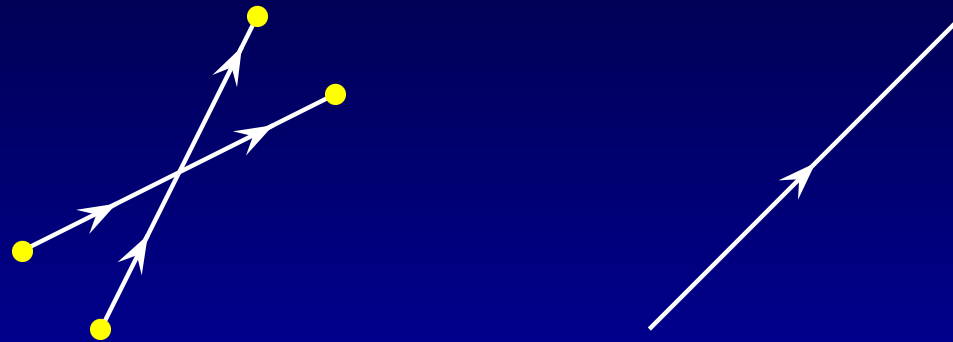
Proposition: Given a dimap $f : B \rightarrow C$ respecting the context and $x, y \in A$ there is an induced map

$$\vec{\pi}_1(f)(x, y) : \vec{\pi}_1(B)(x_B, y_B) \rightarrow \vec{\pi}_1(C)(x_C, y_C).$$

If f is a dihomotopy equivalence then it is an isomorphism.

Example

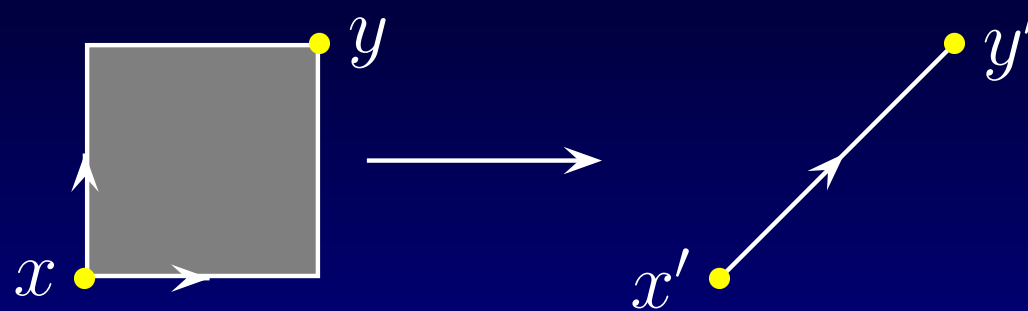
In the context of its four endpoints



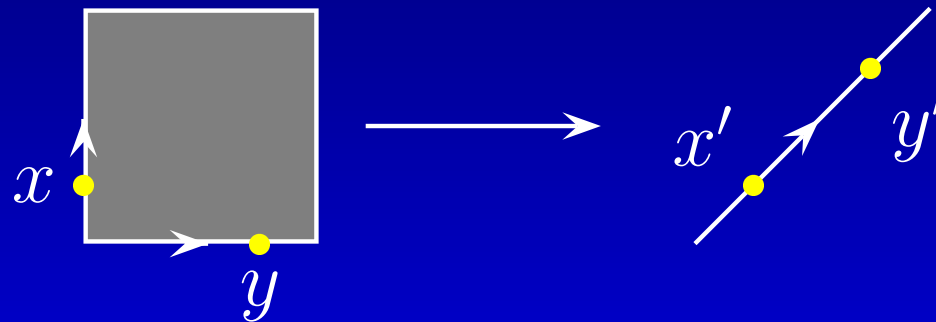
the left hand po-space is not dihomotopy equivalent to \vec{I} .

Another example

Recall that there is a dihomotopy equivalence



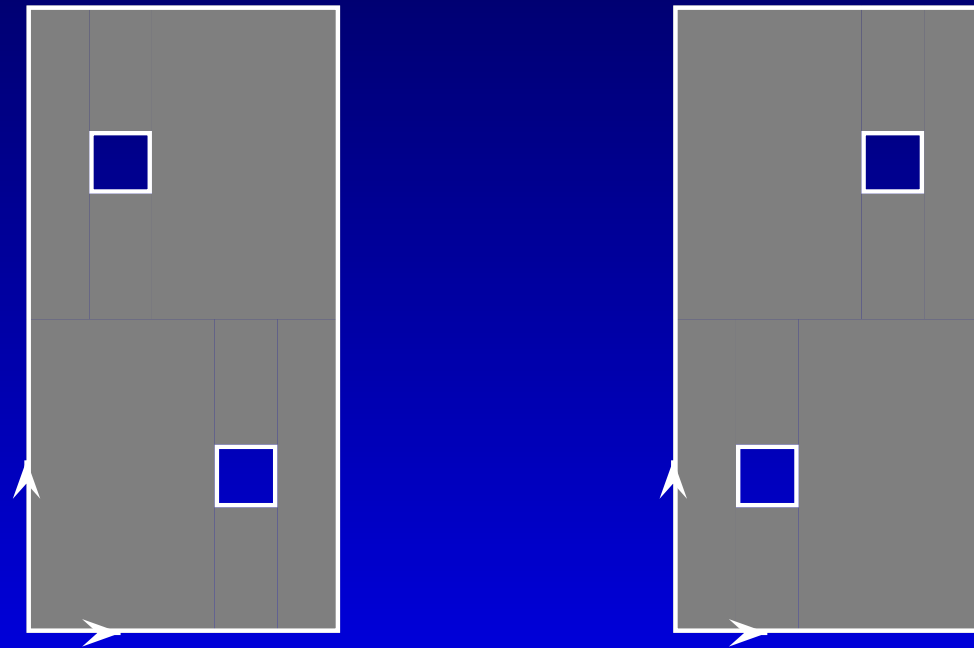
However there is no dihomotopy equivalence



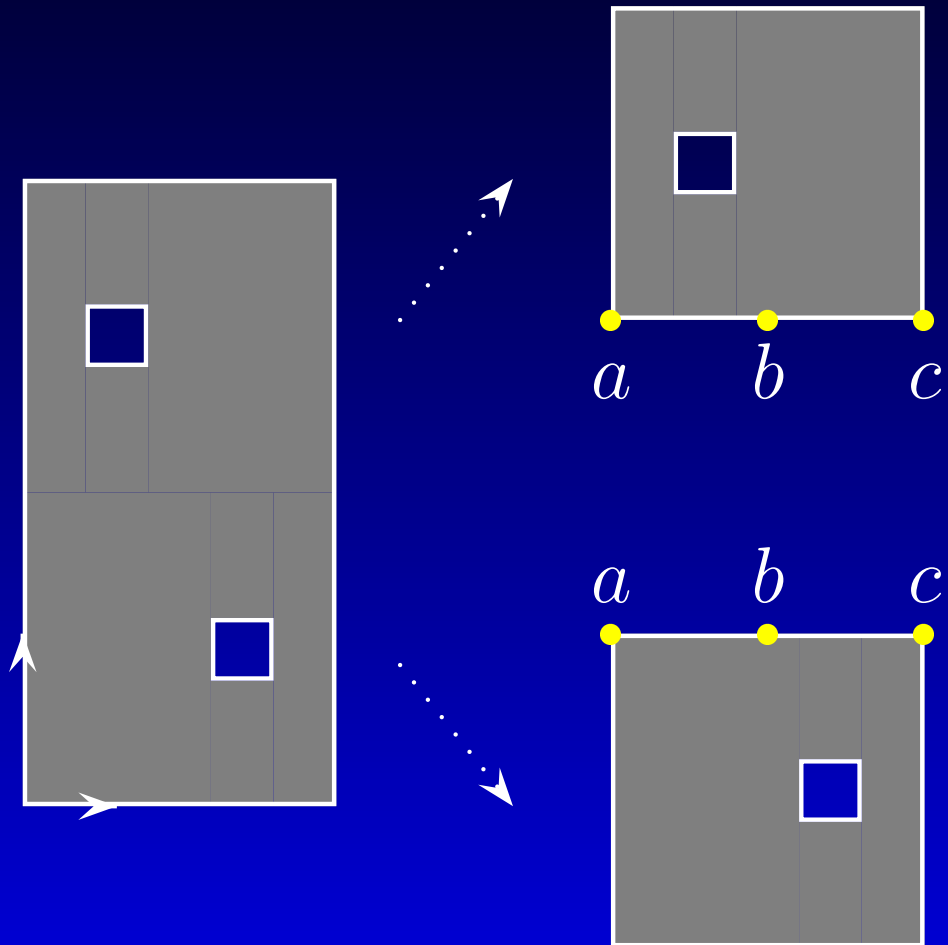
Equivalence depends on the context!

Compound examples

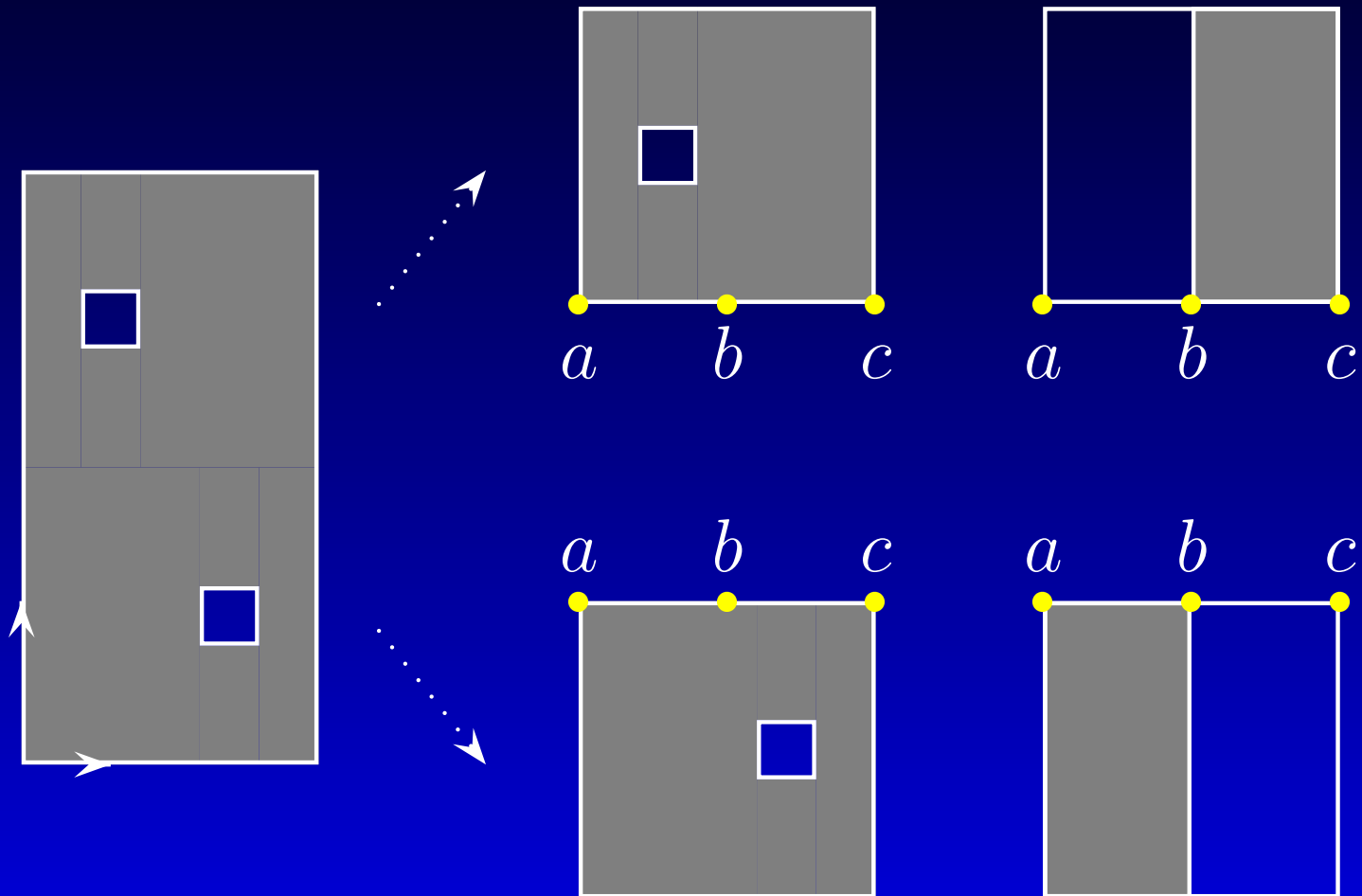
We would like to find equivalent po-spaces to the following examples by analyzing them piece-by-piece.



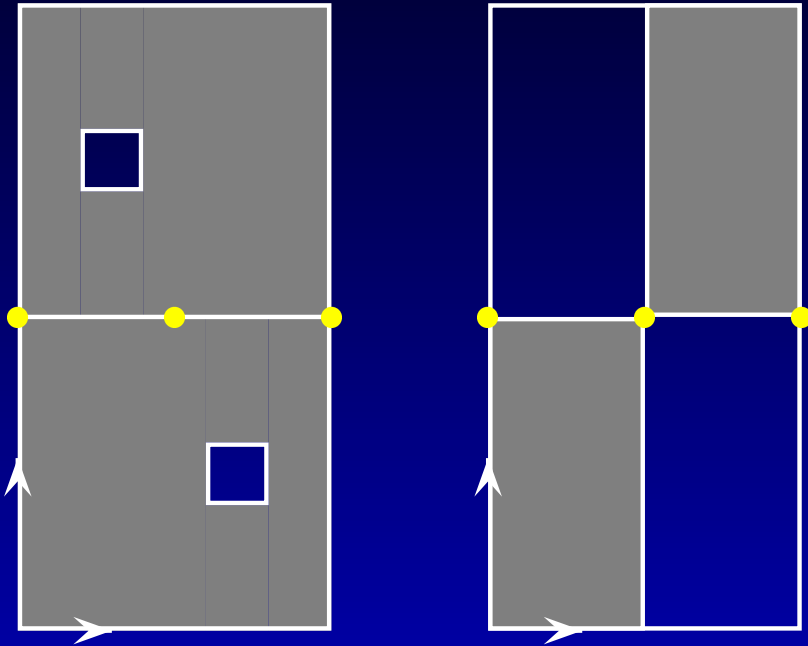
Equivalences of pieces



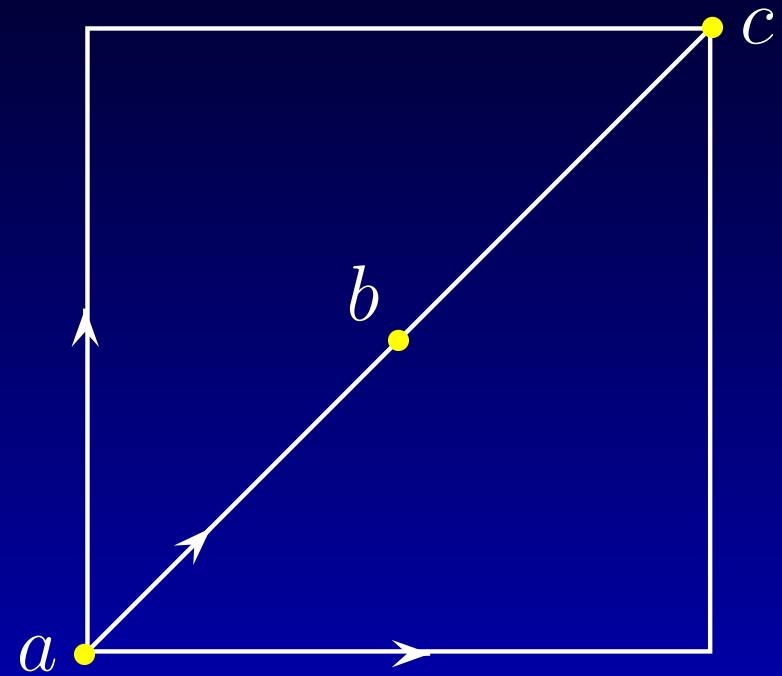
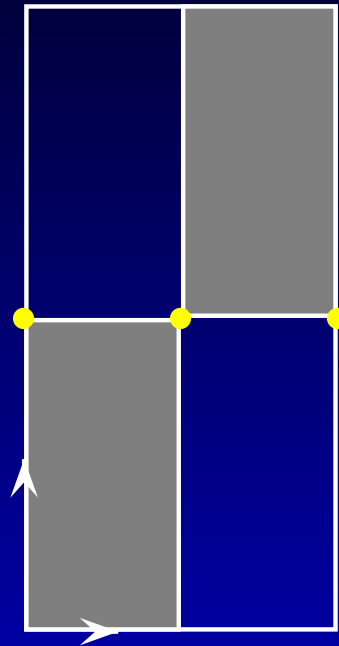
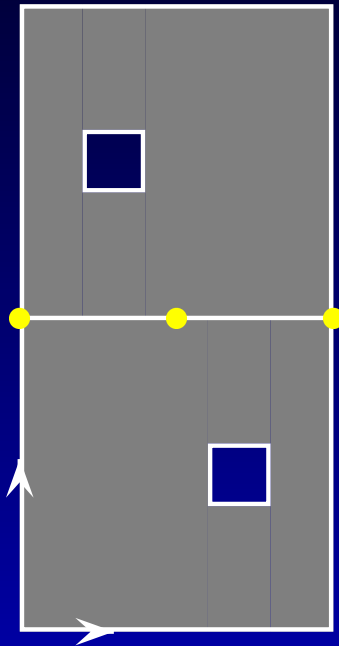
Equivalences of pieces



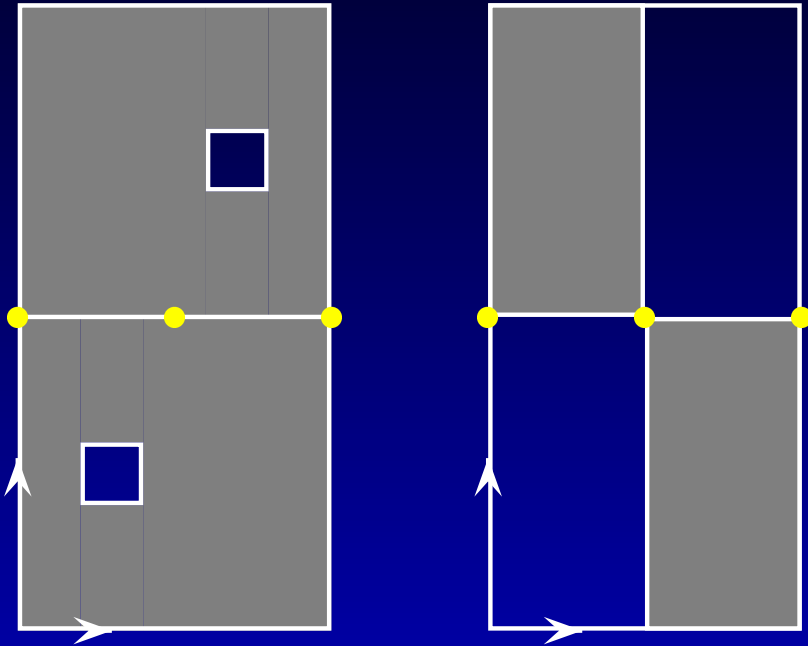
Patching the pieces together



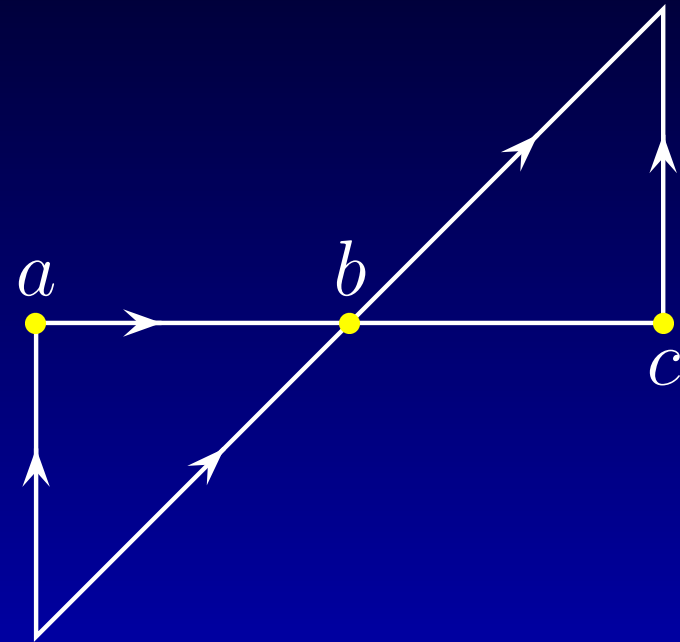
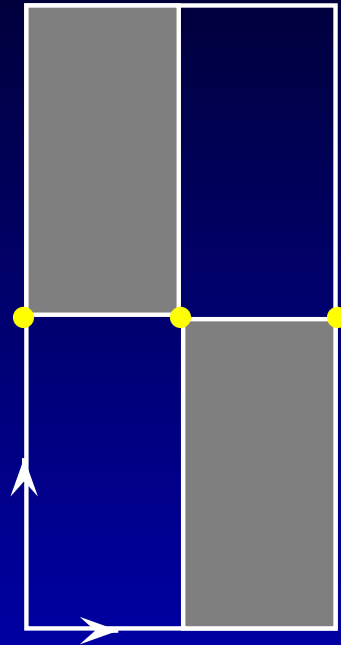
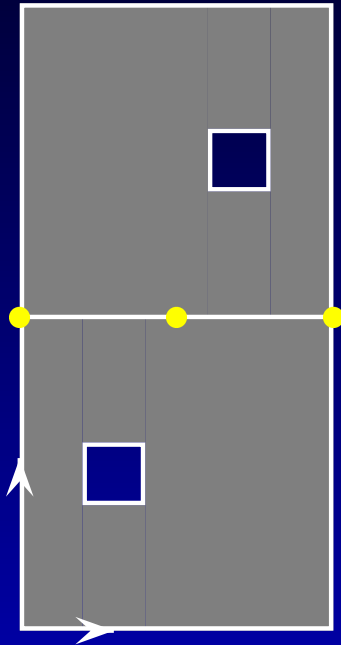
Patching the pieces together



Second example



Second example



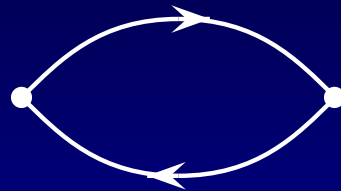
3. Modeling systems with loops

(joint work with Kris Worytkiewicz)

- Local po-spaces
- Model categories
- Equivalences for concurrent systems which may have loops

A more general model

We would like to model execution loops such as



These cannot be modeled by po-spaces.

However they can be modeled by **local po-spaces**.

Local po-spaces

Definition:

- An **order atlas** is a open cover of po-spaces with compatible partial orders.
- A **local po-space** is a topological space together with an equivalence class of order atlases.
- A morphism of local po-spaces is a continuous map which respects the orders.

Equivalences of local po-spaces

Just as with po-spaces, we can define local po-spaces under some context A , and we consider morphisms which respect the context.

We can also define dihomotopy equivalences using context exactly the same way as with po-spaces.

Enter some machinery

A powerful framework for studying equivalences is given by **model categories**.

Definition:

A **model category** is a category (with all small limits and colimits) and with three distinguished classes of morphisms: weak equivalences, cofibrations, and fibrations satisfying four simple axioms.

The structure of a model category allows one to apply the machinery of homotopy theory.

Model category axioms

M0. \mathcal{C} has all small limits and small colimits

M1. 2 out of 3

M2. retracts

M3. lifting property

M4. factorization

A model for concurrent systems

Theorem [B-Worytkiewicz]:

The category of local po-spaces under a context A embeds into a model category such that

- the weak equivalences are the dihomotopy equivalences rel A ,
- the cofibrations are the monomorphisms, and
- pushouts of weak equivalence with cofibrations are weak equivalences

Sheaves, Simplicial presheaves

Definition:

- The category of **presheaves** $\mathbf{Set}^{\mathbf{LoPospc}^{\text{op}}}$ has as objects contravariant functors from $\mathbf{LoPospc}$ to \mathbf{Set} and has as morphisms natural transformations

Sheaves, Simplicial presheaves

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- The category of **presheaves** $\mathbf{Set}^{\mathbf{LoPospc}^{\text{op}}}$ has as objects contravariant functors from $\mathbf{LoPospc}$ to \mathbf{Set} and has as morphisms natural transformations
- Similarly there is a category of **simplicial presheaves** $\mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$

Sheaves, Simplicial presheaves

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- Similarly there is a category of **simplicial presheaves** $\mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$
- There is a Yoneda embedding
 $\mathbf{LoPospc} \hookrightarrow \mathbf{Set}^{\mathbf{LoPospc}^{\text{op}}} \hookrightarrow \mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$

Sheaves, Simplicial presheaves

Definition:

- The category of **presheaves** $\mathbf{Set}^{\mathbf{LoPospc}^{\text{op}}}$ has as objects contravariant functors from $\mathbf{LoPospc}$ to \mathbf{Set} and has as morphisms natural transformations
- Similarly there is a category of **simplicial presheaves** $\mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$
- There is a Yoneda embedding $\mathbf{LoPospc} \hookrightarrow \mathbf{Set}^{\mathbf{LoPospc}^{\text{op}}} \hookrightarrow \mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$
- The **sheaves** $\mathbf{Shv}(\mathbf{LoPospc})$ are the presheaves which are compatible with the topology

Sketch of the proof

Theorem[Jardine]:

Let \mathbf{C} be a small category with a Grothendieck topology. Then $\mathbf{sSet}^{\mathbf{C}^{op}}$ the category of simplicial presheaves on \mathbf{C} has a (proper, simplicial) model structure in which

- the cofibrations are the monomorphisms, and
- the weak equivalences are the local weak equivalences.

Furthermore, if the Grothendieck topos $\mathit{Shv}(\mathbf{C})$ has enough points then the local weak equivalences are the stalkwise equivalences.

Grothendieck topology

Proposition[B-W]: $\mathbf{LoPospc}$ has a Grothendieck topology given by **open directed covers**.

Enough points

Let Z be a local po-space. **Definition:** The category of directed étale bundles over Z has

- objects: dimaps $E \rightarrow Z$ which are local homeomorphisms
- morphisms: maps $E_1 \rightarrow E_2$ such that

$$\begin{array}{ccc} E_1 & \xrightarrow{\quad} & E_2 \\ & \searrow & \swarrow \\ & Z & \end{array} \quad \text{commutes}$$

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Theorem[B-W]: There is an equivalence between $\mathbf{Shv}(Z)$ and $\mathbf{Etale}(Z)$.

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Theorem[B-W]: There is an equivalence between $\mathbf{Shv}(Z)$ and $\mathbf{Etale}(Z)$.

Corollary: $\mathbf{Shv}(\mathbf{LoPospc})$ has enough points.

Stalkwise equivalences

Using the above results, it follows from Jardine's theorem that $\mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$ has a model structure in which the weak equivalences are the stalkwise equivalences.

Stalkwise equivalences

Using the above results, it follows from Jardine's theorem that $\mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$ has a model structure in which the weak equivalences are the stalkwise equivalences.

Proposition[B-W]: The stalkwise equivalences coming from $\mathbf{LoPospc}$ are the isomorphisms.

Adding context

Given context $A \in \mathbf{LoPospc}$ there is a induced model structure on

$$A \downarrow \mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$$

Finally one can localize with respect to the dihomotopy equivalences $\text{rel } A$ to obtain the main theorem. \square

4. Summary

- It would be useful to have a robust notion of equivalence in models of concurrency.
- To allow a piece-by-piece analysis we would like equivalences that remain equivalences even after additions are made to the model.
- Using **context** provides such equivalences.

Summary

- Using po-spaces, local po-spaces, context, and model categories, we have a good mathematical framework for studying concurrent systems.
- In particular, using equivalences, this framework should allow for a piece-by-piece analysis of concurrent systems.

Future Work

Theoretical: Consider all possible contexts in a single framework. (Stay tuned for Kris' talk!)

Practical: Use the current theoretical framework for analyzing real-world examples.

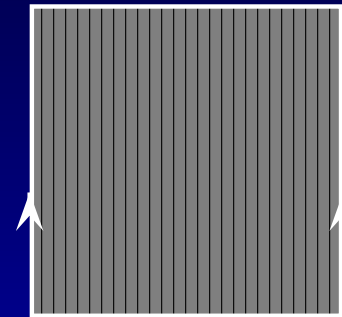
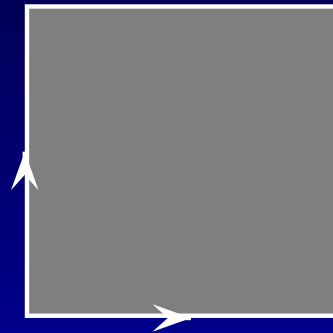
Acknowledgments

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This talk and associated preprints are available at <http://igat.epfl.ch/bubenik/>

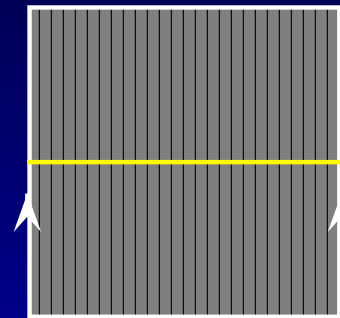
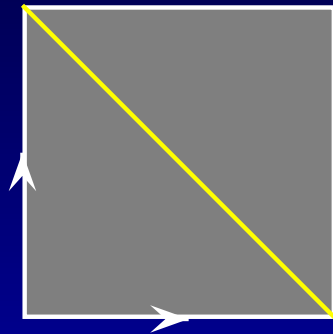
Non-discrete context

Take $\vec{I} \times \vec{I}$ and $I \times \vec{I}$



Non-discrete context

Take $\vec{I} \times \vec{I}$ and $I \times \vec{I}$



and glue them together along the yellow lines.