

# Using context and model categories to define directed homotopies

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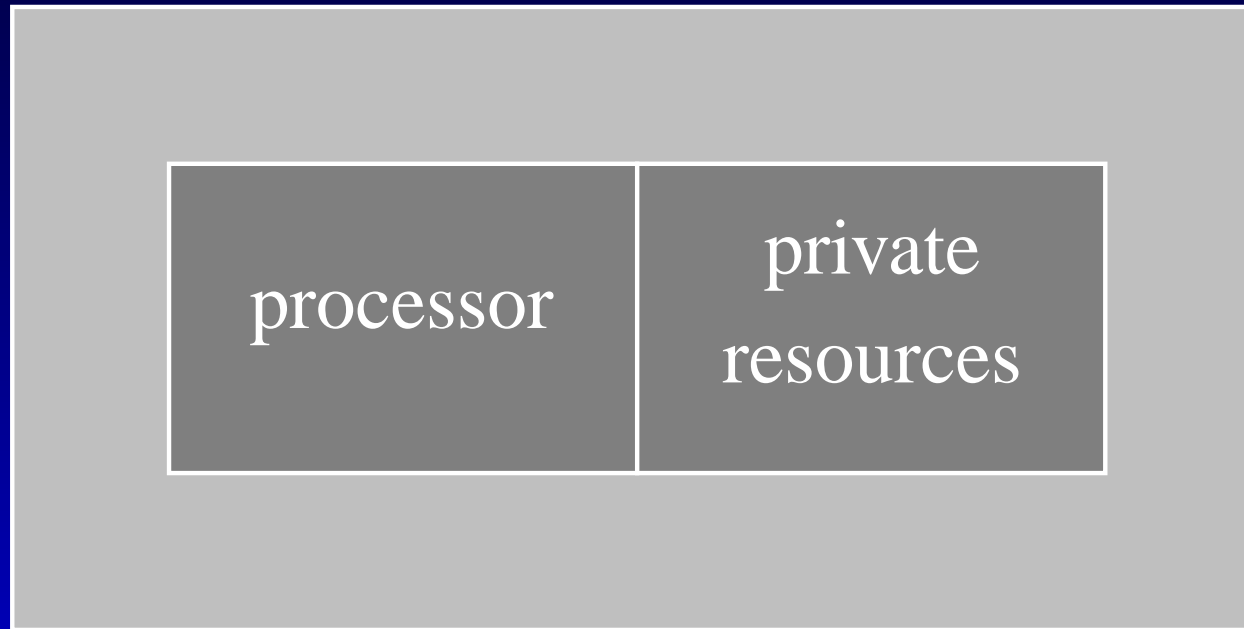
<http://igat.epfl.ch/bubenik/>

# 1. Introduction

- A concurrent system
- A model for concurrency
- Equivalences - undirected and directed
- The Problem

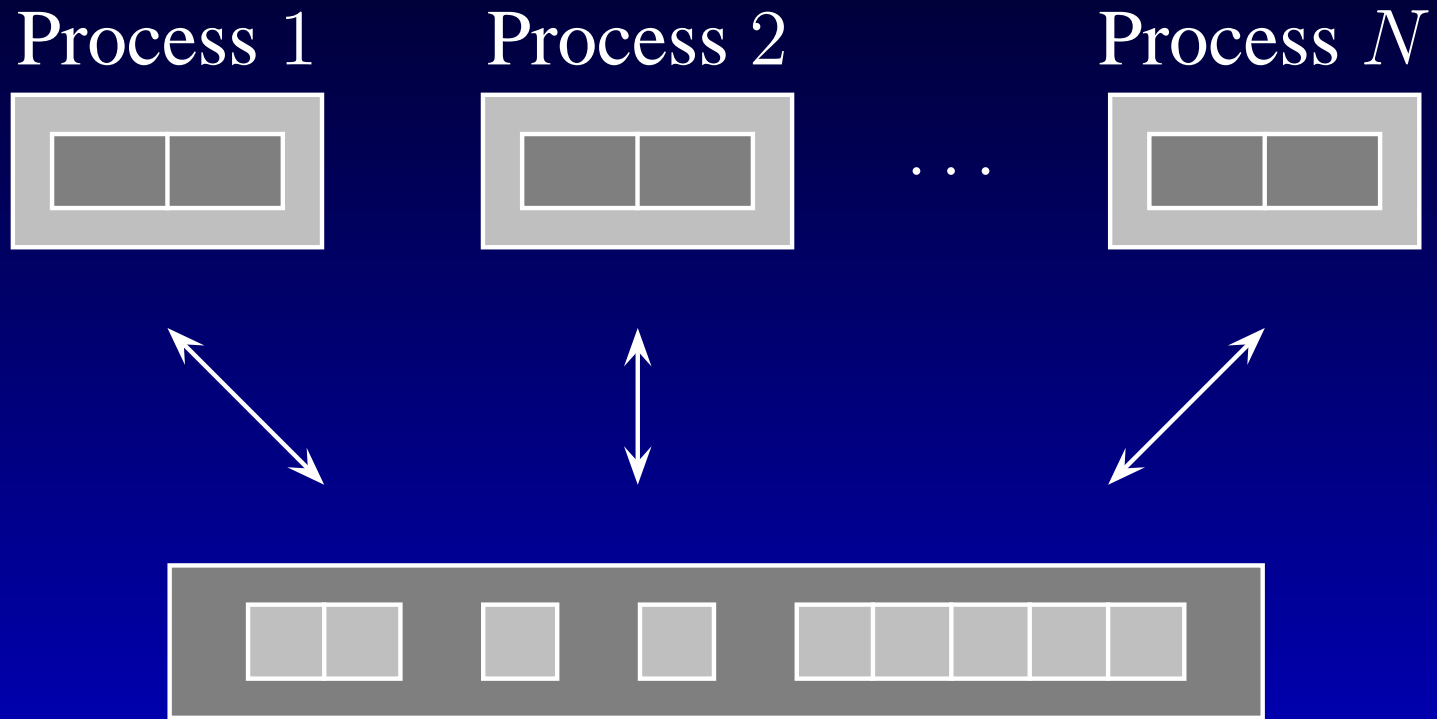
# A nonconcurrent system

Process



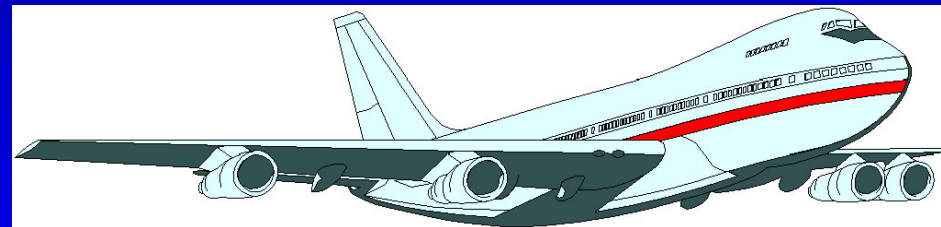
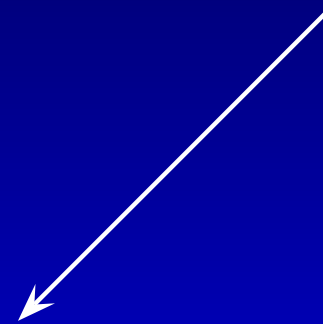
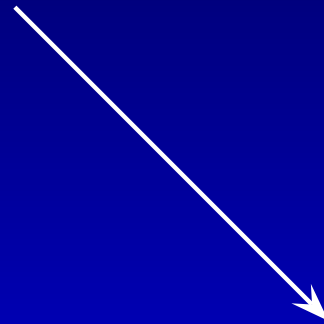
A process with its own private resources

# A concurrent system



Several processes with shared resources

# Example: Internet database





# Example: dual core processors

Intel (Q3 2005)


AMD (Q3 2005)

**digital ENTERPRISE**  
multi-core transition

	2005	2006	FUTURE
	EM64T Larger cache (2M) Faster FSB DDR2 Power Management (DBS) PCI-Express* I/O	<b>DUAL CORE</b> Lower Power Cores Enhanced Memory (FBD) Virtualization (VT) iAMT I/O Packet Acceleration Advanced Storage Controllers	<b>MULTI-CORE</b> 2 or more  Advanced memory, virtualization, RAS and manageability
	<b>DUAL CORE</b> Multi-Threading Pellston Foxton	<b>DUAL CORE</b> Multi-Threading Pellston Foxton	<b>MULTI-CORE</b> 4 or more  Advanced memory, virtualization, RAS and manageability

**DUAL/MULTI-CORE RAMP**  
ALIGNED with PLATFORM FEATURES

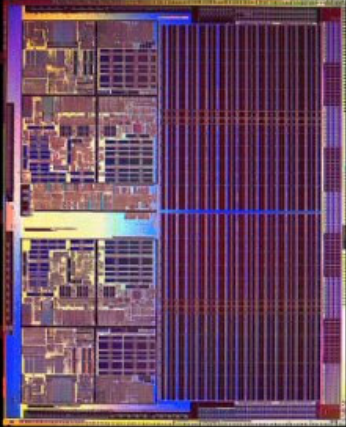
Source: Intel roadmap. \*Other names and brands may be claimed as the property of others.

Mid-2005: AMD Opteron™ Processor Dual-Core 


**90nm Process**  
Approximately same die size as 130nm single-core AMD Opteron™ processor\*  
~205 million transistors\*

**95 Watt Power Envelope**  
Fits into 90nm power infrastructure

**940 Socket Compatible**  
All that's needed is a BIOS upgrade  
Compatible with all motherboards designed to our 90nm specification



\*Based on current revisions of the design  
11/12/2004

  
2004 Analyst Day

# A concurrent system

## Example:

2 processes using 2 shared resources  $a$  and  $b$  which can only be used by one process at a time

## Notation:

$Px$  - a process locks resource  $x$

$Vx$  - a process releases resource  $x$

Program of the first process:  $Pa \quad Pb \quad Vb \quad Va$

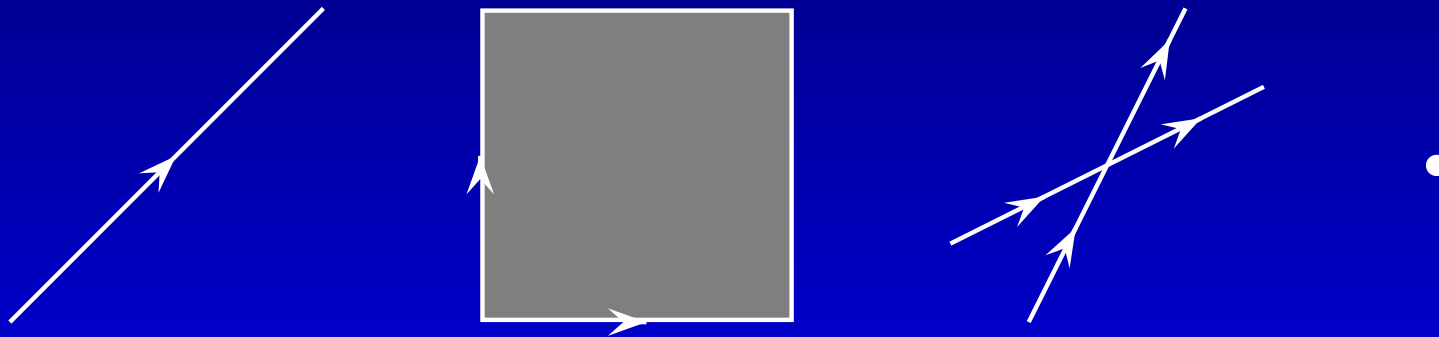
Program of the second process:  $Pb \quad Pa \quad Va \quad Vb$

# A model for concurrency

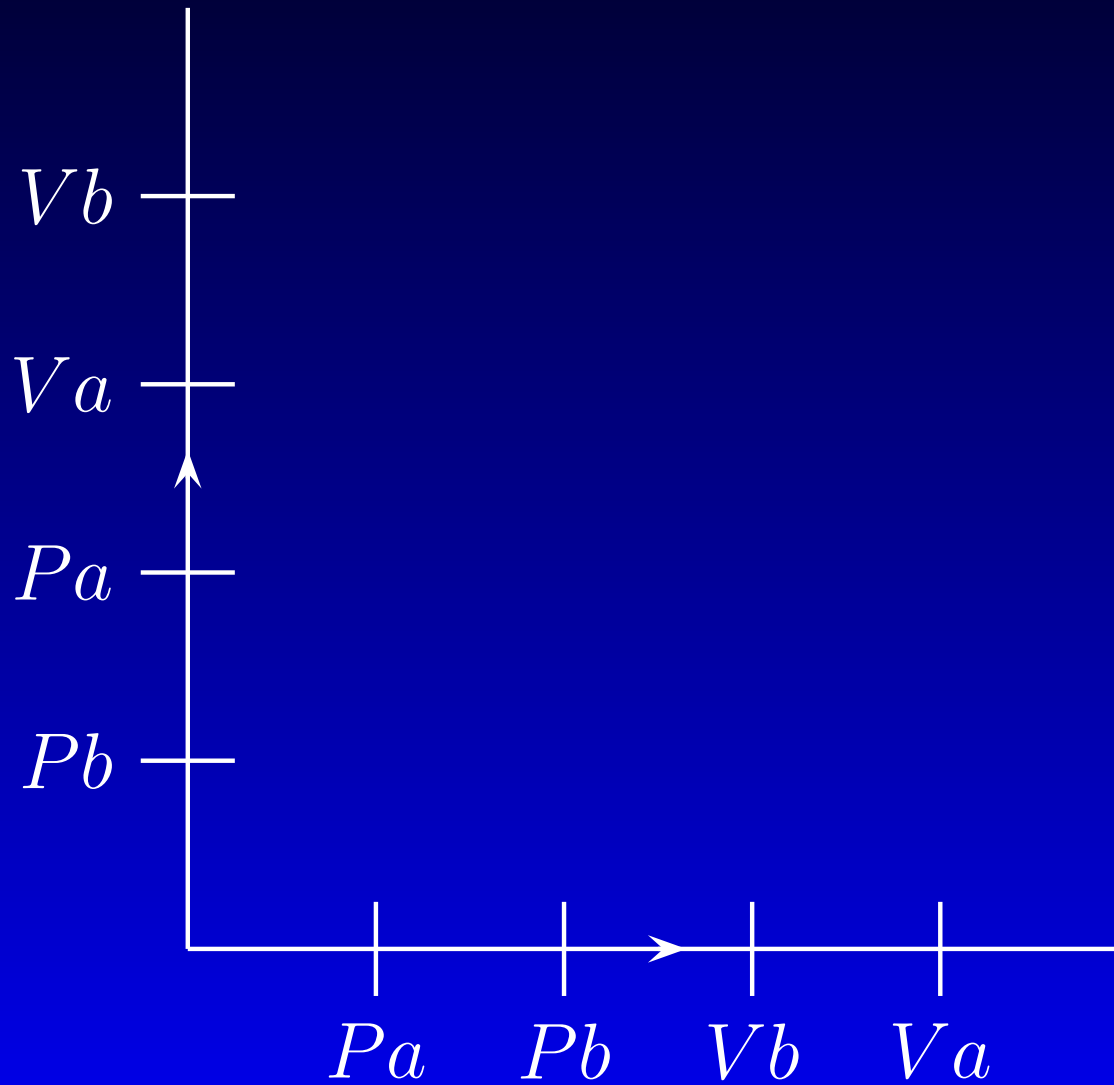
Concurrent systems can be modeled by subspaces of  $\mathbb{R}^n$  together with a partial order.

**Definition:** A **po-space** is a topological space  $U$  with a partial order  $\leq$  which is a closed subset of  $U \times U$ .

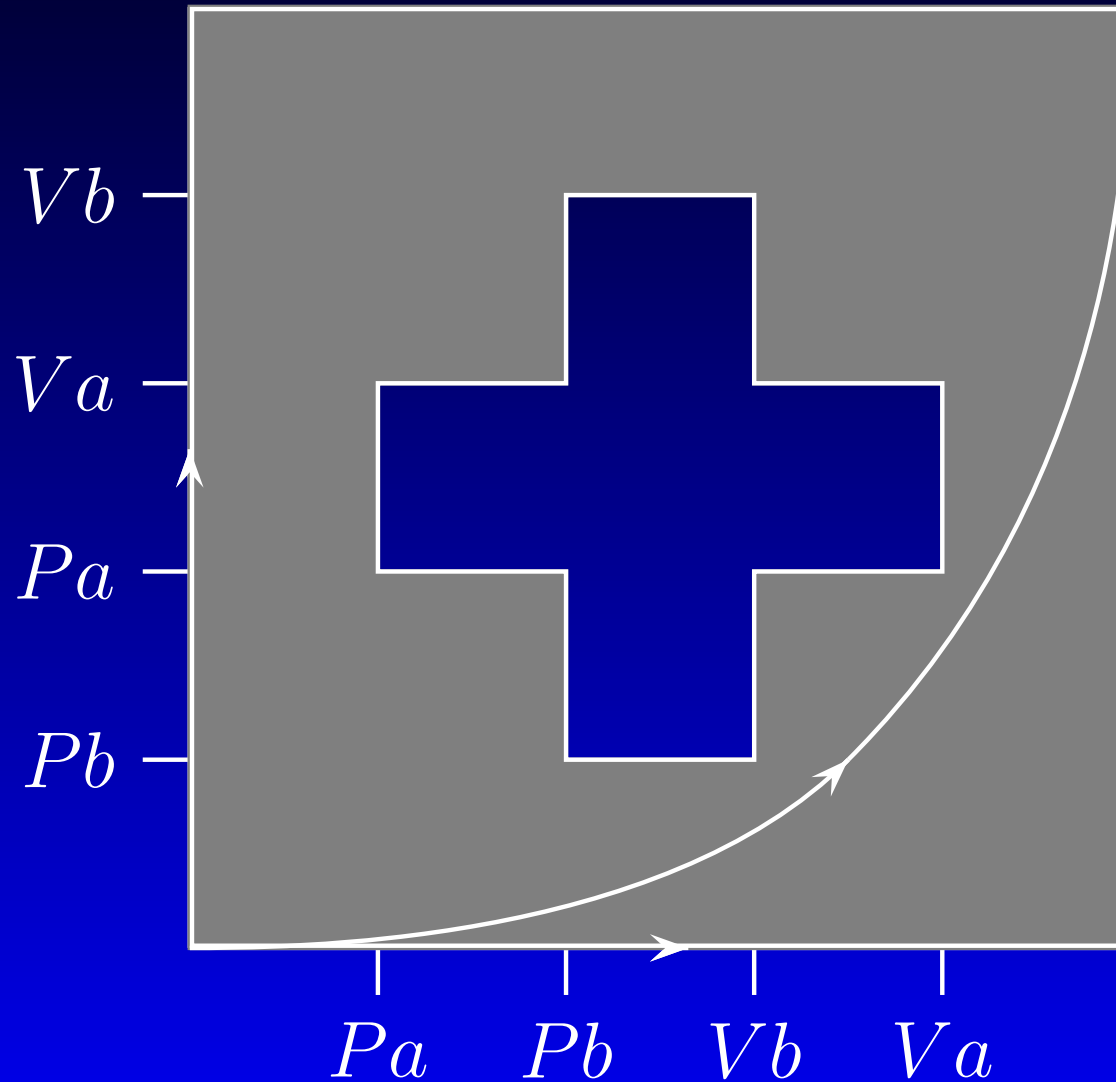
**Examples:**



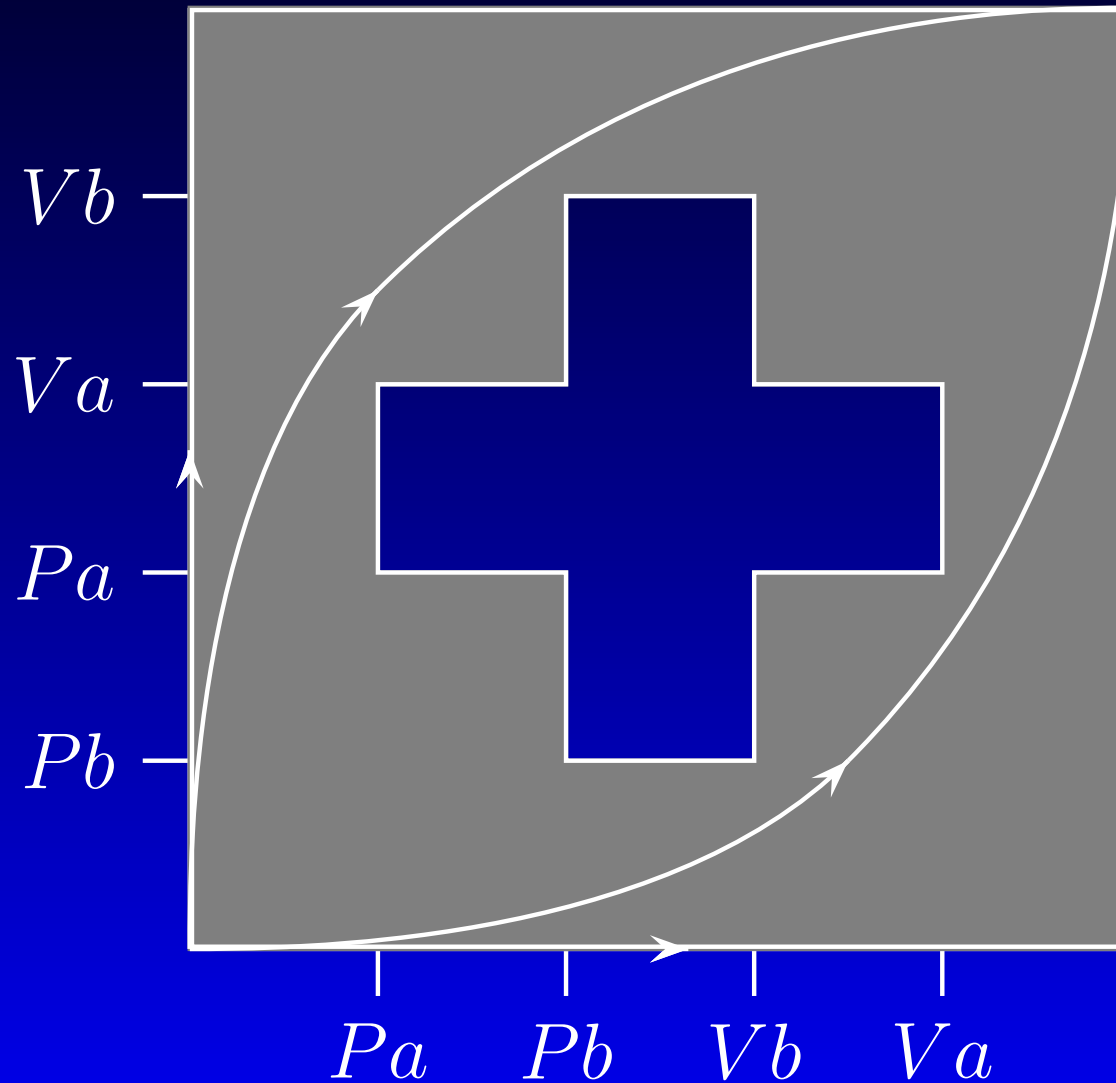




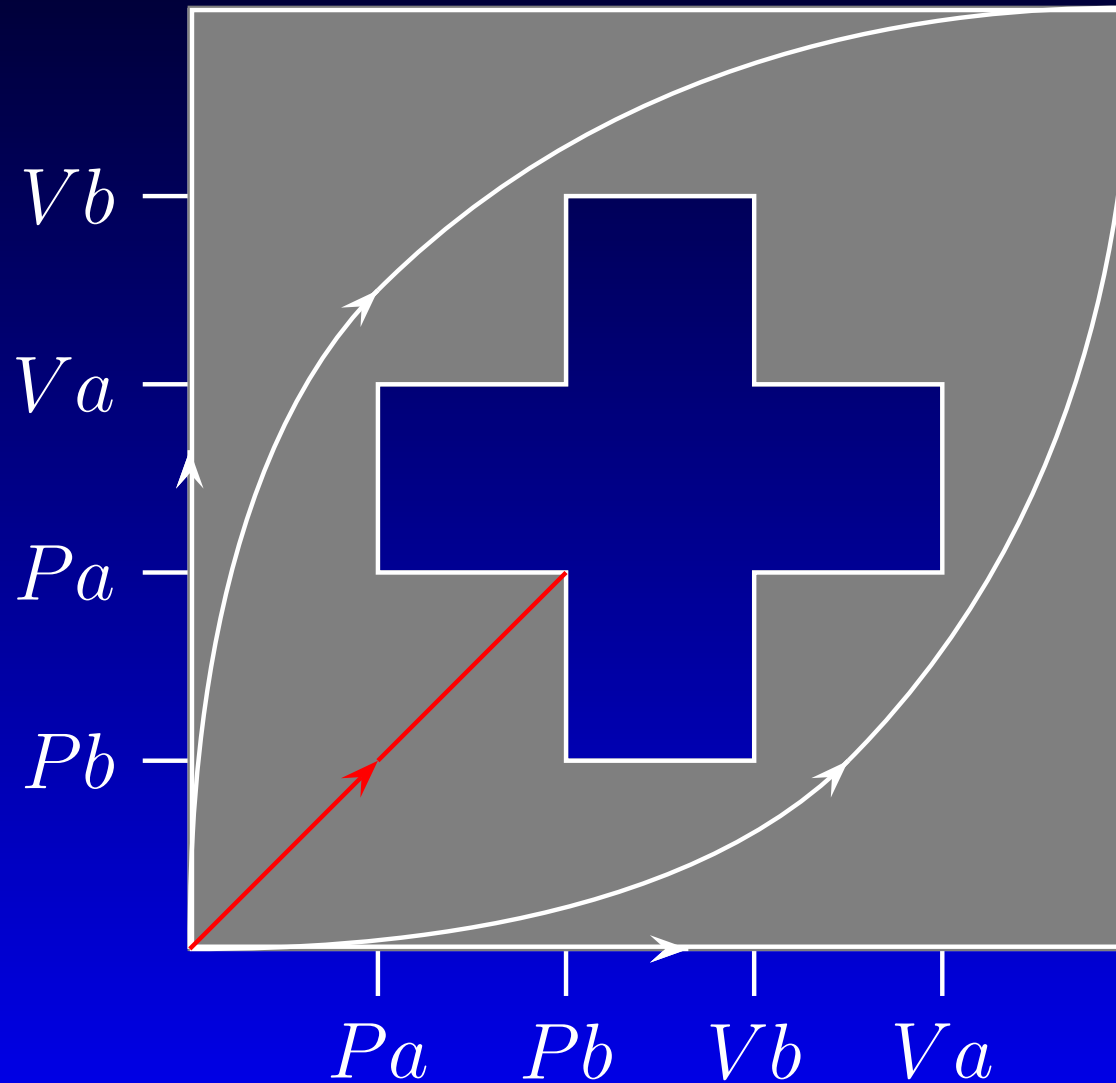
# The Swiss flag



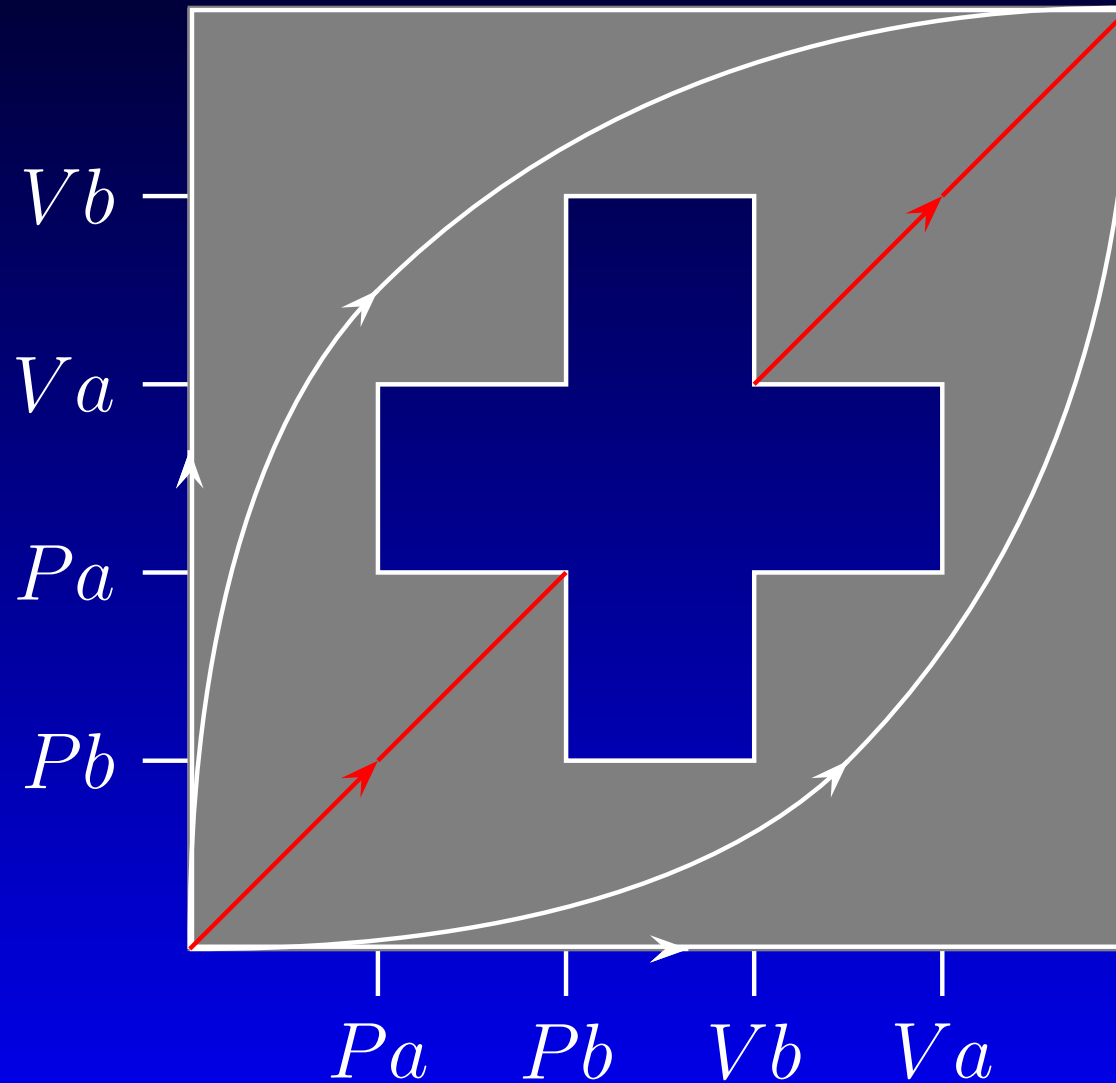
# The Swiss flag



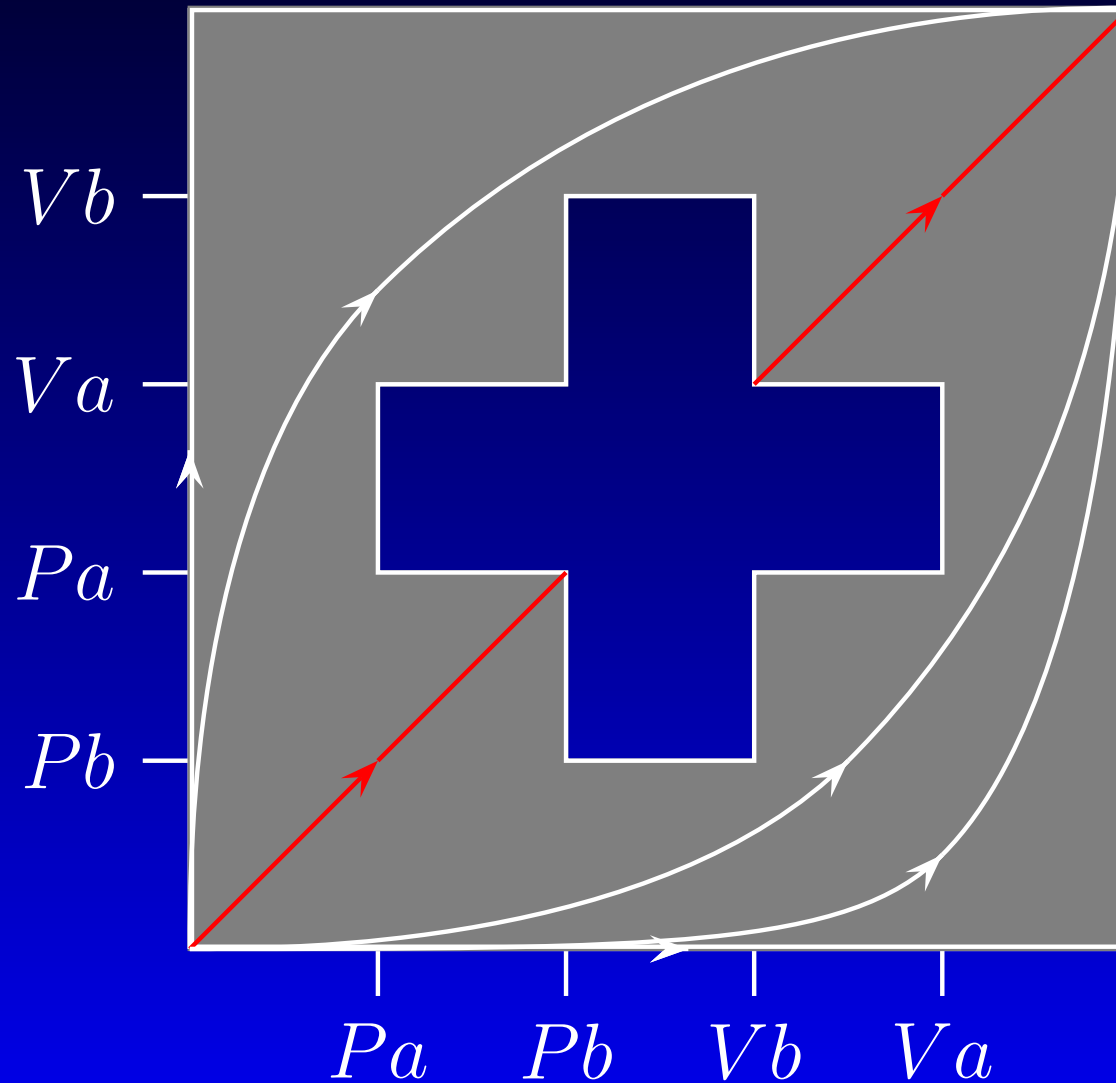
# The Swiss flag



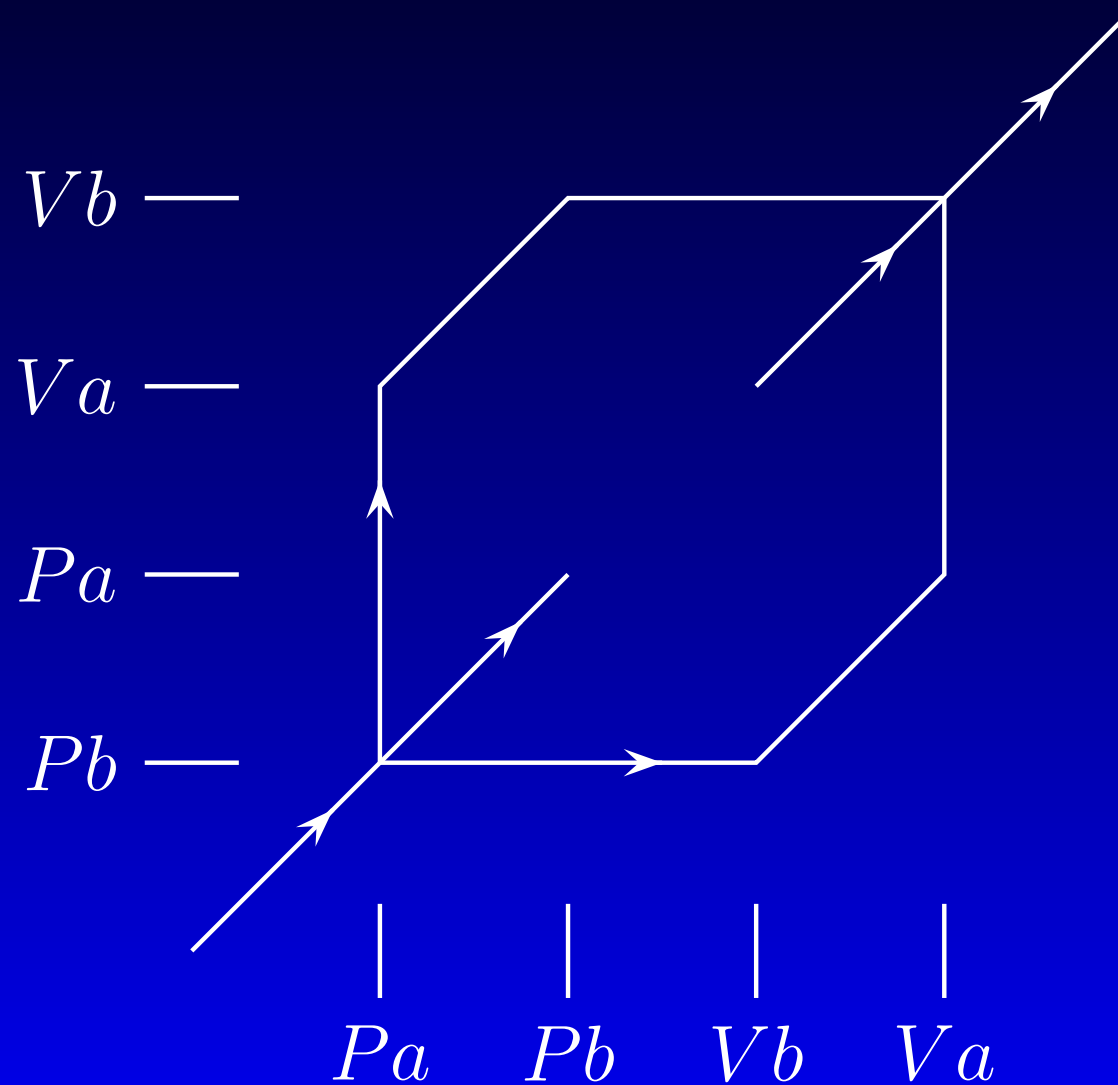
# The Swiss flag



# The Swiss flag



# A sub-po-space of the Swiss flag



# Goal

Develop a framework for concurrency where equivalences are accounted for.

We would like equivalences that allow a piece-by-piece analysis of po-spaces.



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Develop a framework for concurrency where equivalences are accounted for.

We would like equivalences that allow a piece-by-piece analysis of po-spaces.

**Idea:** Use algebraic topology.

# Undirected equivalences

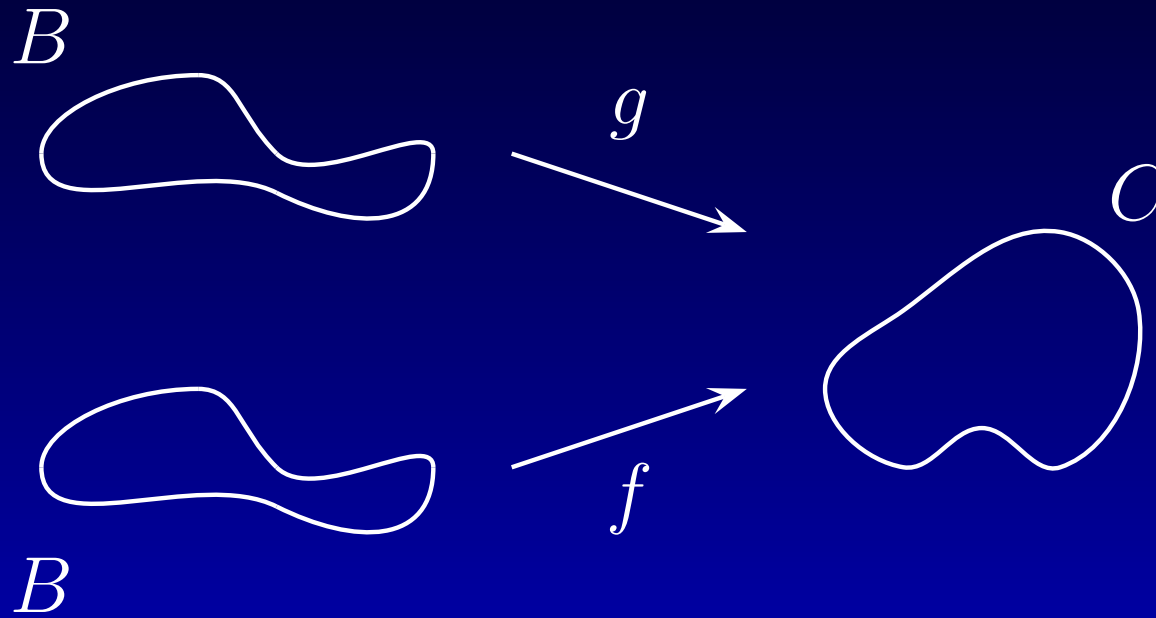
Working with (undirected) spaces and continuous maps, equivalences are defined using homotopies.

Maps are equivalent if there is a homotopy between them.

Spaces are equivalent if there maps between them whose compositions are homotopic to the identity map.

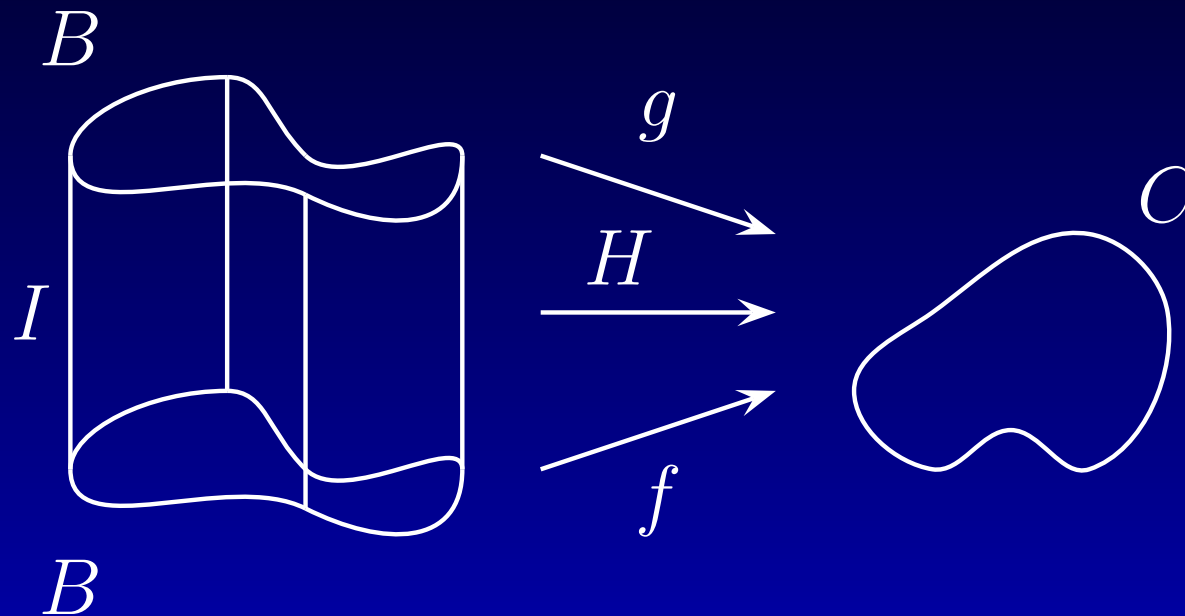
# Undirected equivalences

**Definition:** Given continuous maps  $f, g : B \rightarrow C$ ,



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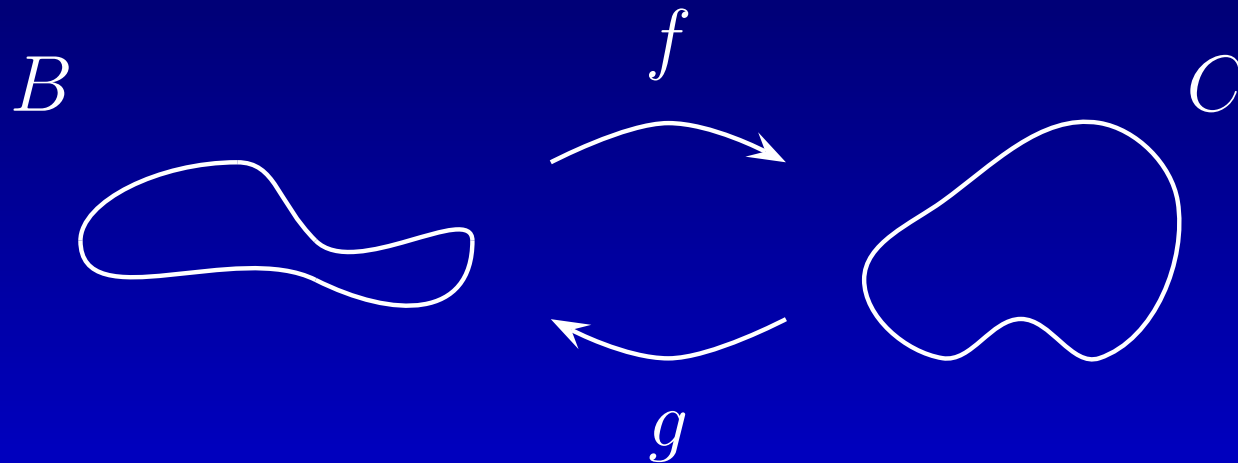


a **homotopy** between  $f$  and  $g$  is a continuous map  $H : B \times I \rightarrow C$  restricting to  $f$  and  $g$ . This is an equivalence relation. Write  $H : f \xrightarrow{\sim} g$ .

# Undirected equivalences

**Definition:** Spaces  $B, C$  are **homotopy equivalent** if there are maps  $f : B \rightleftarrows C : g$  such that

$$g \circ f \simeq \text{Id}_B \text{ and } f \circ g \simeq \text{Id}_C .$$



# Directed spaces

## Definition:

- A **po-space** is a topological space  $U$  with a partial order  $\leq$  which is a closed subset of  $U \times U$ .
- A **directed map (dimap)** is a continuous map  $f : U_1 \rightarrow U_2$  between po-spaces such that

$$x \leq y \implies f(x) \leq f(y).$$

- Let **PoSpc** be the category whose objects are po-spaces and morphisms are dimaps.

# Directed spaces

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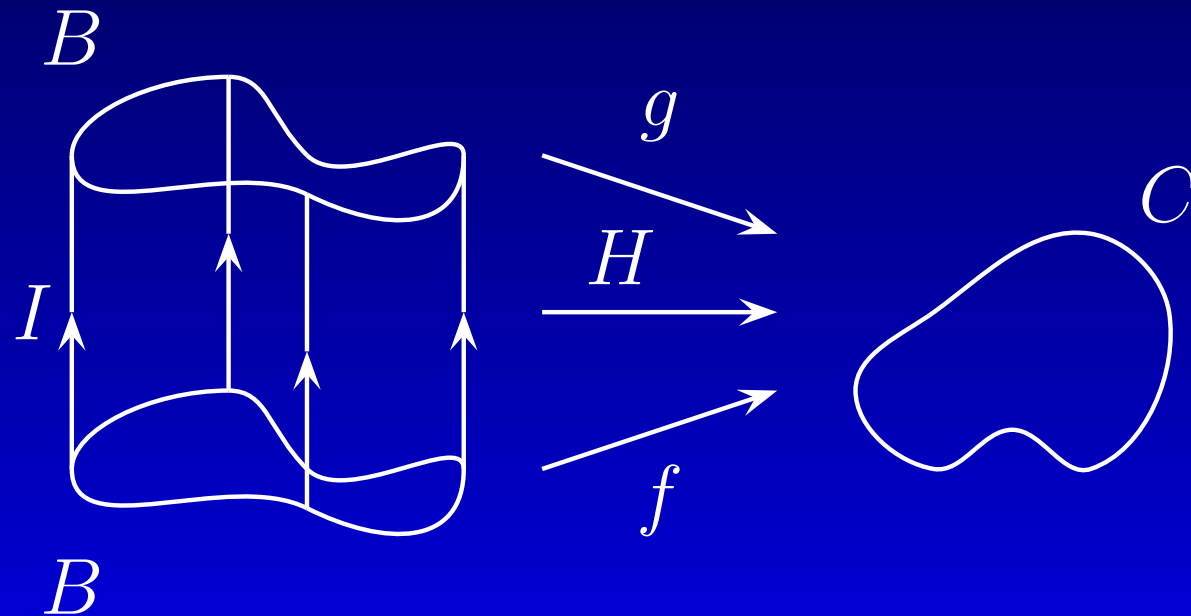
- Let **PoSpc** be the category whose objects are po-spaces and morphisms are dimaps.

**Remark:** Subspaces and products of po-spaces inherit a po-space structure.

# Directed equivalences

## Definition:

- A **directed homotopy (dihomotopy)** between dimaps  $f, g : B \rightarrow C$  is a dimap  $H : B \times \vec{I} \rightarrow C$  restricting to  $f$  and  $g$ . Write  $H : f \rightarrow g$ .





# Directed equivalences

## Definition:

- Write  $f \simeq g$  if there is a chain of dihomotopies

$$f \rightarrow f_1 \leftarrow f_2 \rightarrow \dots \leftarrow f_n \rightarrow g.$$

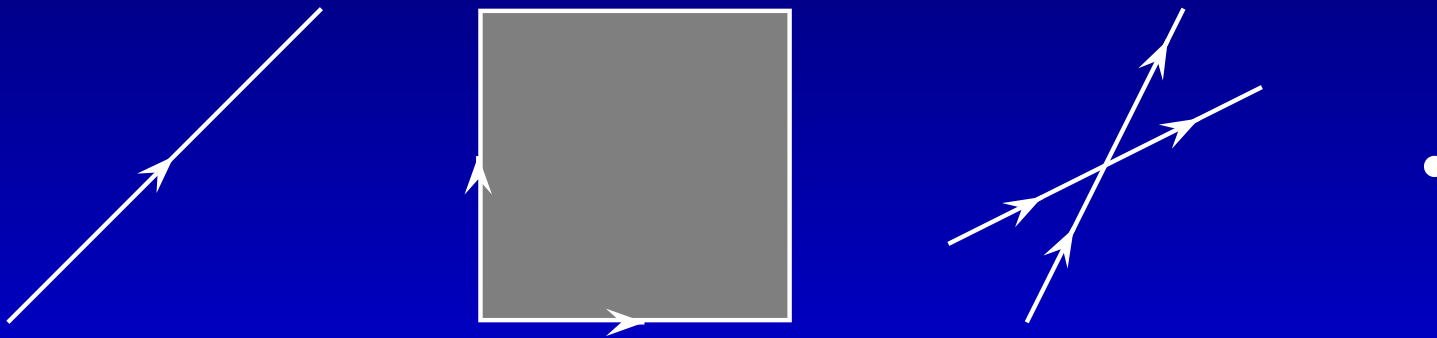
- Po-Spaces  $B, C$  are **dihomotopy equivalent** if there are dimaps  $f : B \rightleftarrows C : g$  such that

$$g \circ f \simeq \text{Id}_B \text{ and } f \circ g \simeq \text{Id}_C .$$

# The Problem

**Recall:** We wanted to use dihomotopy equivalences to provide equivalences of concurrent systems.

However all of the following spaces are dihomotopy equivalent.



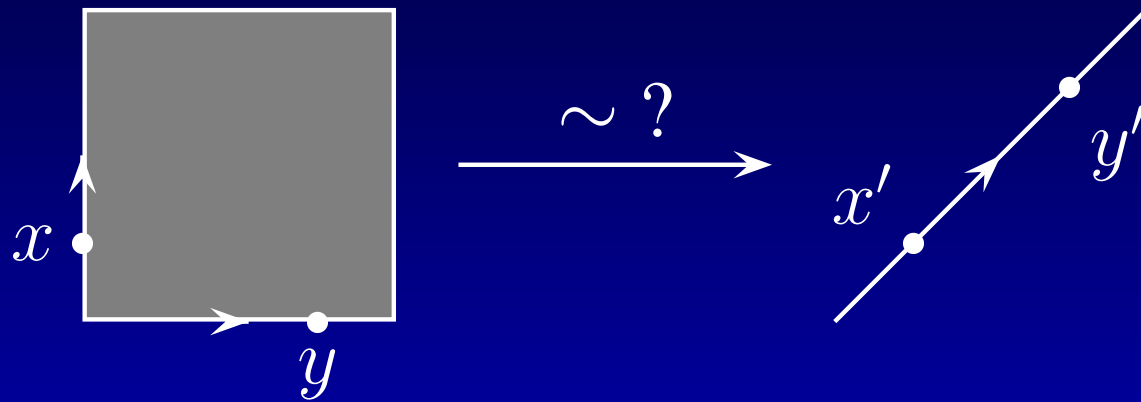
Thus, a stronger notion of equivalence is needed.

## 2. Using context

- A basic example
- A solution
- Equivalences using context
- Non-equivalences using directed paths
- Piece-by-piece example

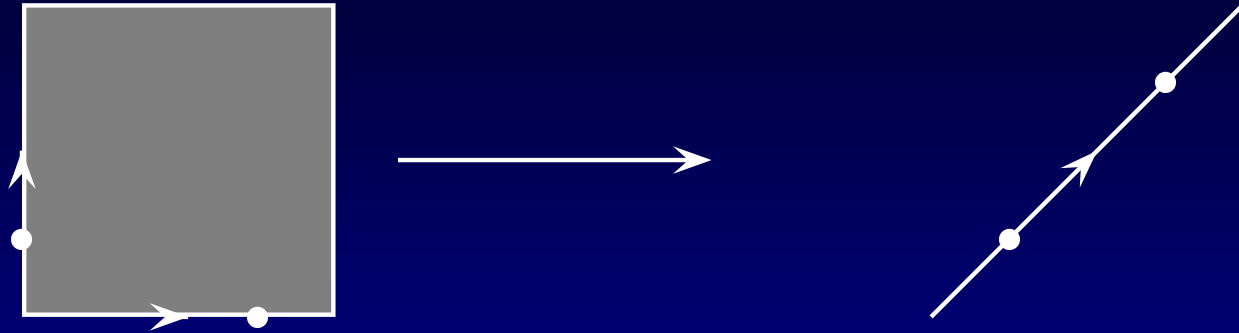
# A basic example

**Question:** Is there an equivalence between  $\vec{I} \times \vec{I}$  and  $\vec{I}$ ?

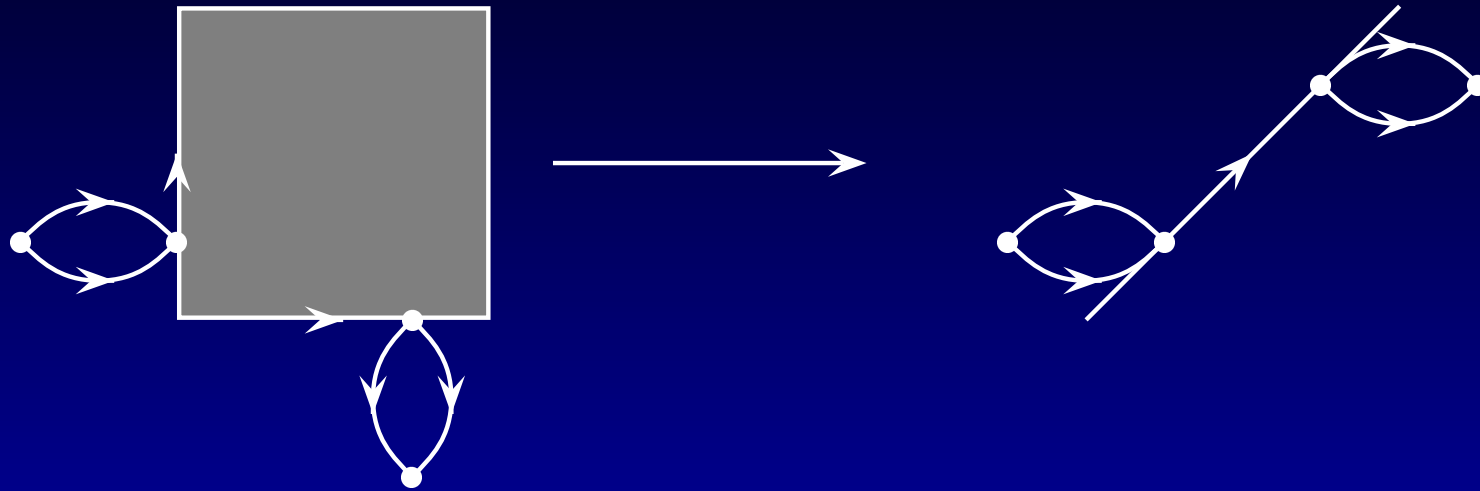


**Our guide:** If this map is an equivalence, then so is the map obtained by attaching po-spaces to this example. (Formally, we want the set of equivalences to be closed under pushouts with inclusions.)

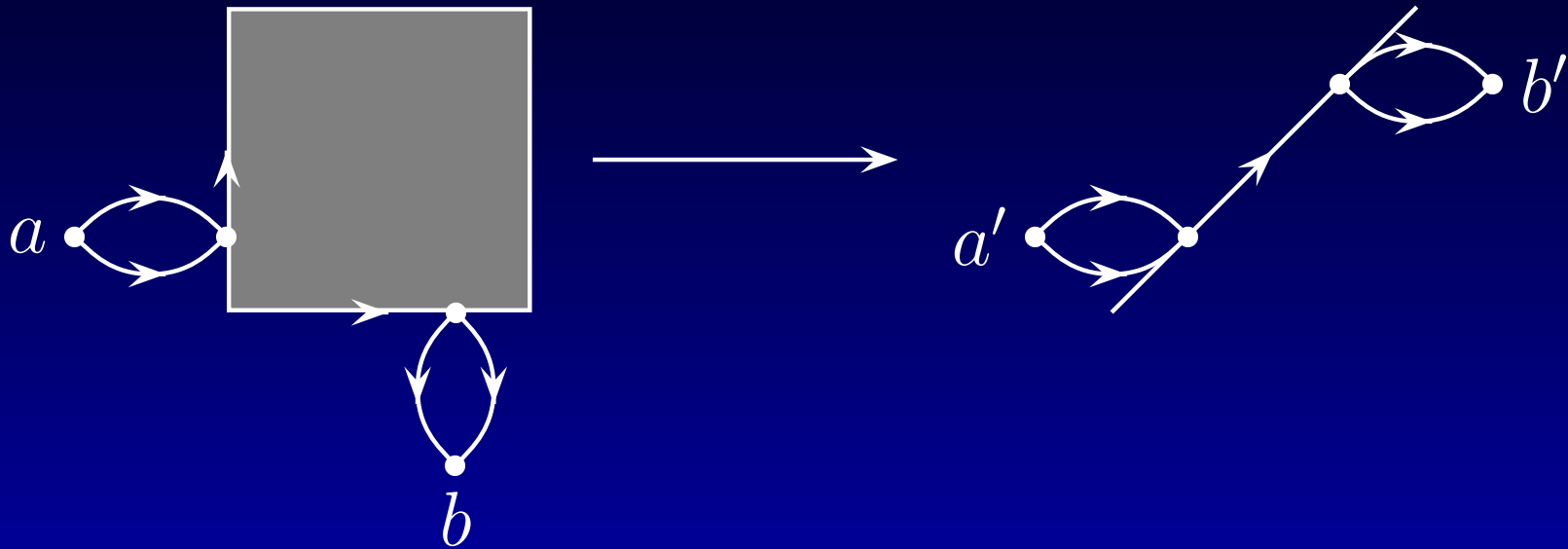
# Adding to the basic example



# Adding to the basic example



# Adding to the basic example



There is no execution path from  $a$  to  $b$ , while there are such paths from  $a'$  to  $b'$ .

Thus, this map should not be an equivalence.

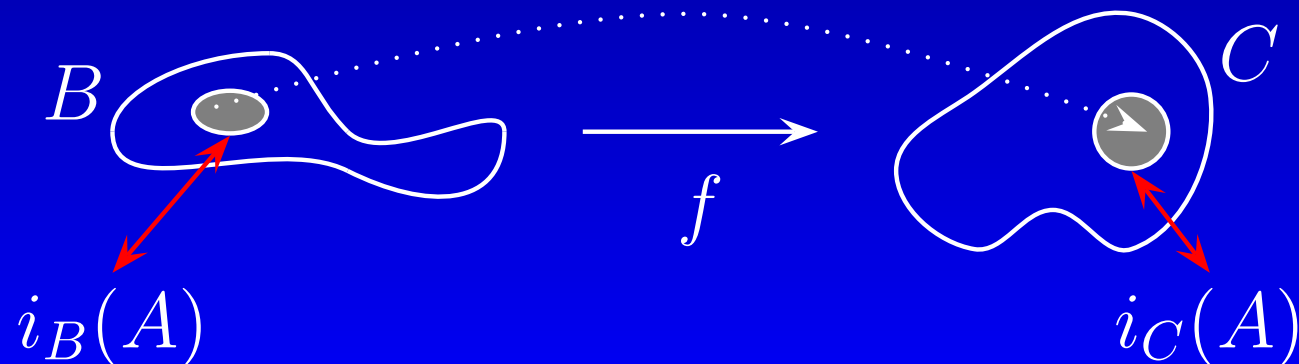
# A Solution

Instead of working with just po-spaces work with po-spaces together with **context**.

**Definition:** Choose a po-space  $A$  (called the **context**). This choice will depend on the attachments one wants to consider.

Consider po-spaces  $B$  together with a dimap  $i_B : A \rightarrow B$  and consider morphisms which are dimaps such that

$$f(i_B(a)) = i_C(a) \text{ for all } a \in A$$

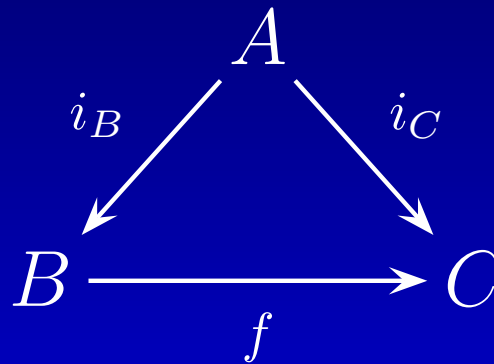




# $A \downarrow \mathbf{PoSpc}$

**Definition:** Given a pospace  $A$ , let  $A \downarrow \mathbf{PoSpc}$  be the category whose

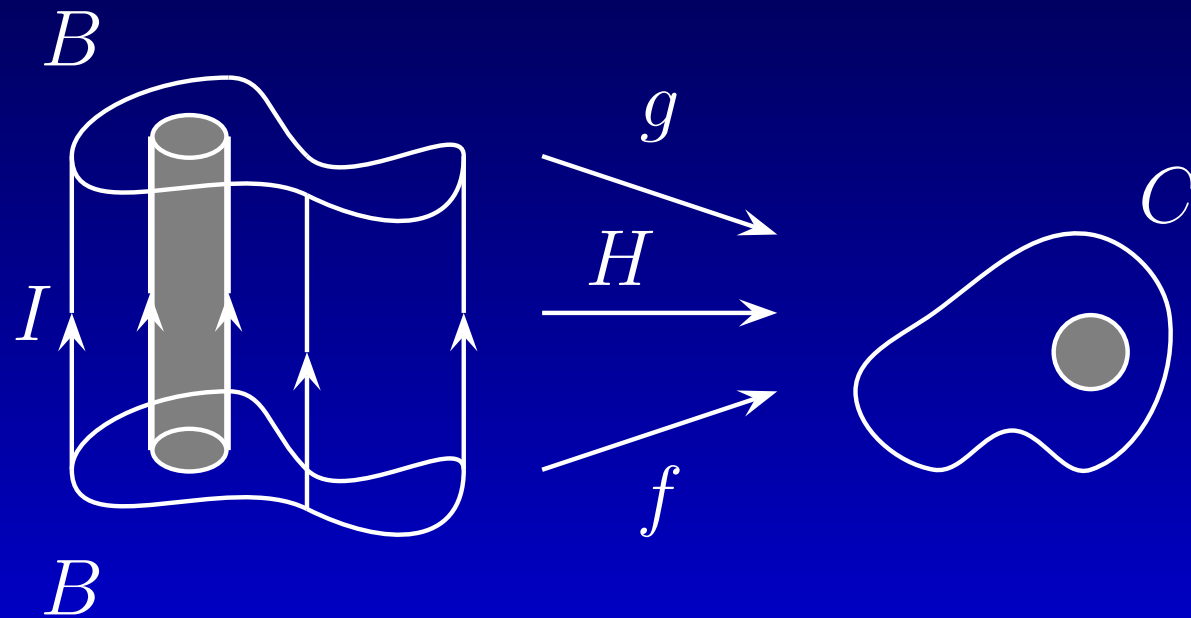
- objects are dimaps  $\iota_B : A \rightarrow B$ ,
- morphisms are dimaps  $f : B \rightarrow C$  such that  $f \circ \iota_B = \iota_C$



# Equivalences using context

## Definition:

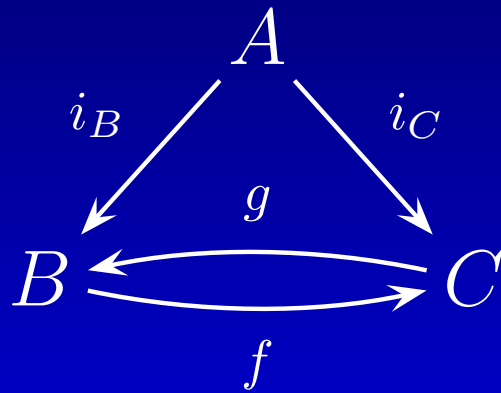
- A **dihomotopy** between  $f, g : B \rightarrow C$  in the context of  $A$  is a dihomotopy  $H : f \rightarrow g \text{ rel } A$ .



# Equivalences using context

## Definition:

- Write  $f \simeq g$  if there is a chain of dihomotopies  $f \rightarrow f_1 \leftarrow f_2 \rightarrow \dots \leftarrow f_n \rightarrow g$ .
- $i_B : A \rightarrow B, i_C : A \rightarrow C$  are **dihomotopy equivalent** if there are dimaps

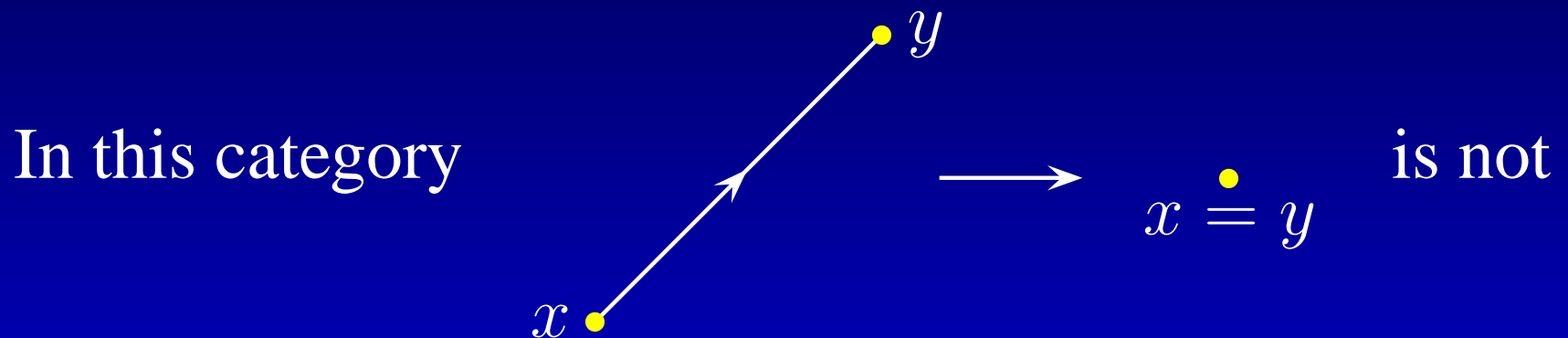


such that  $g \circ f \simeq \text{Id}_B$   
and  $f \circ g \simeq \text{Id}_C$ .

# Example

Let  $A = \{x, y\}$  with  $x \leq y$ .

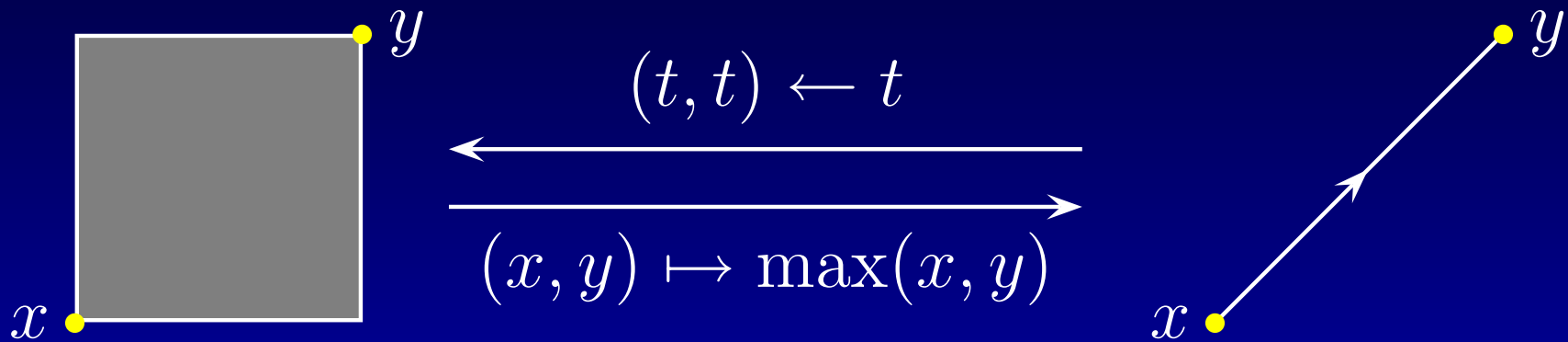
Then po-spaces under the context  $A$  are just po-spaces with two marked points, one of which is after the other.



a dihomotopy equivalence since there is no dimap in the reverse direction.

# Example of an equivalence

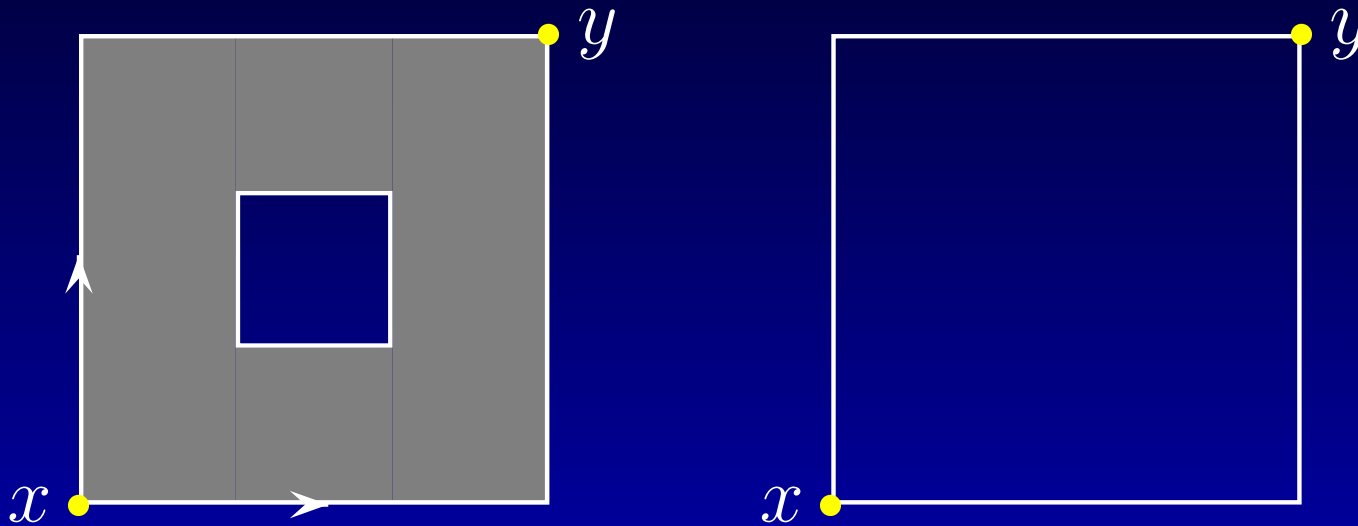
In the same context (of two marked points) the dimaps



give a dihomotopy equivalence.

# Another equivalence

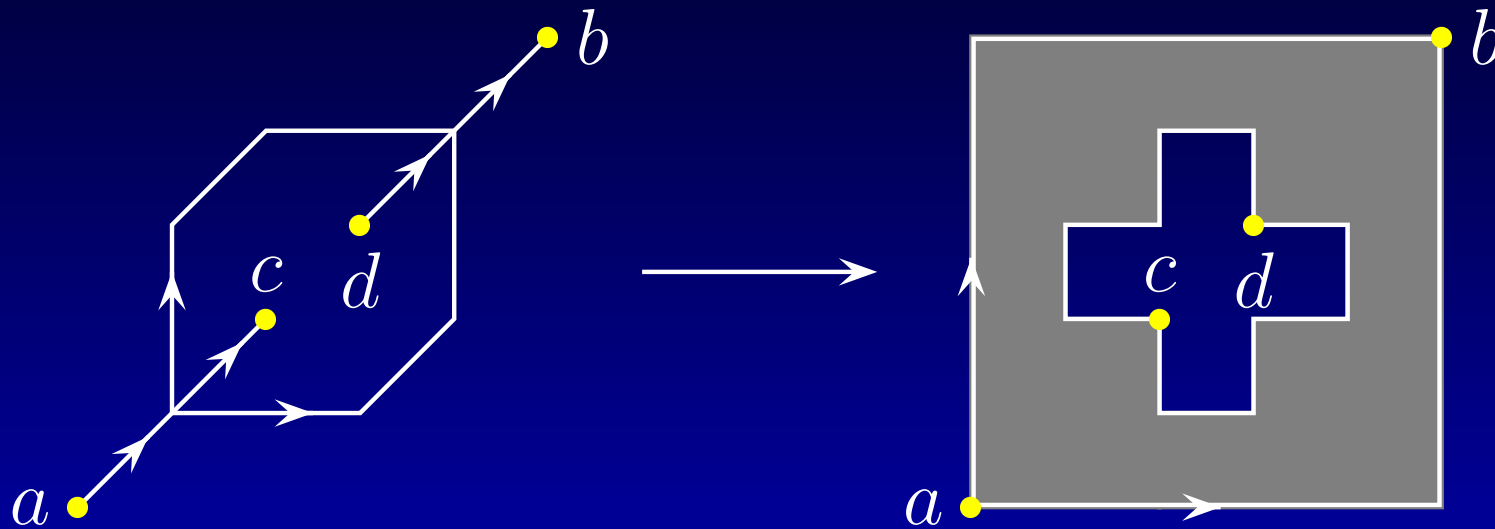
$\vec{I} \times \vec{I}$  with a square removed and two marked points



is dihomotopy equivalent to its boundary.

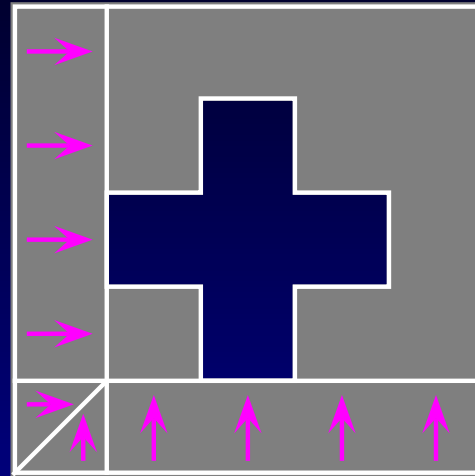
# Context for the Swiss flag

Let  $A = \{a, b, c, d\}$ . Then the inclusion



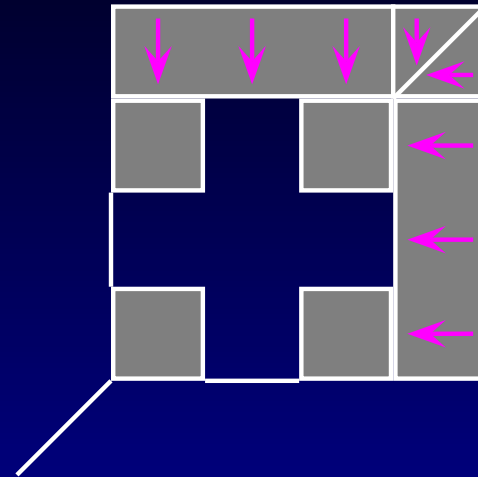
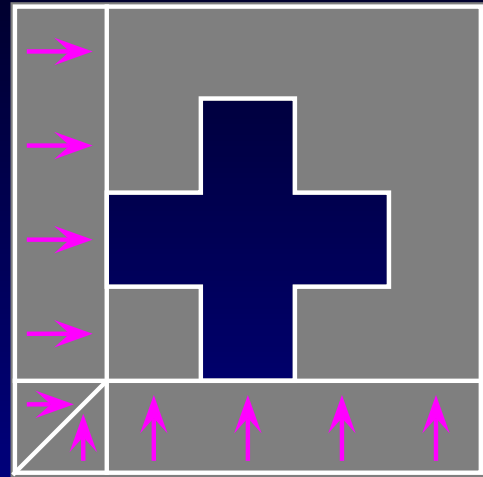
is a dihomotopy equivalence in the context of the four marked points.

# Sketch of the proof

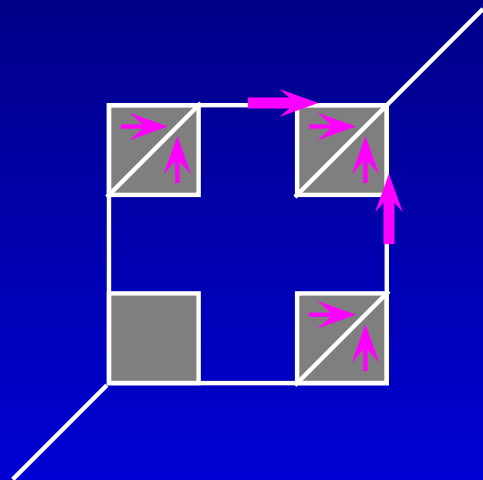
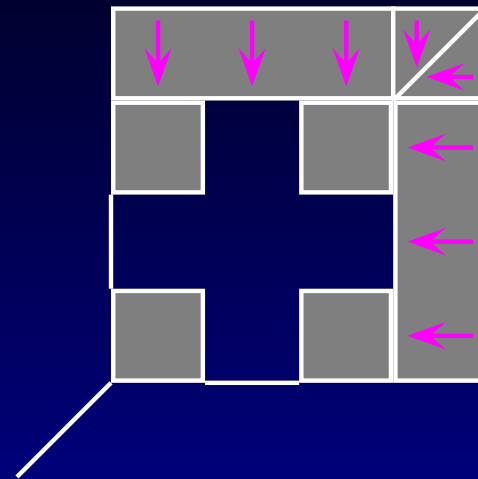
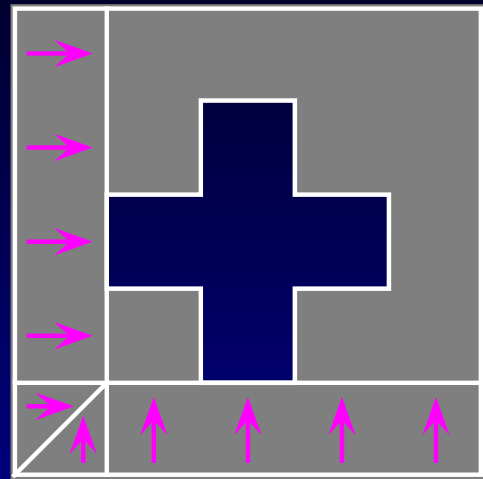




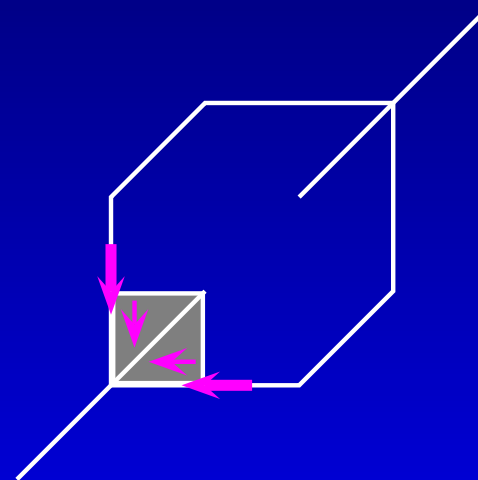
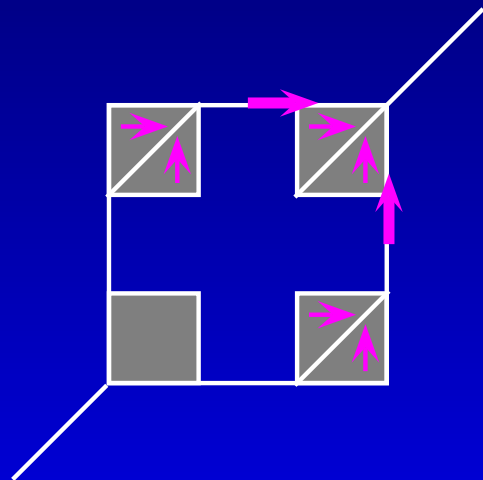
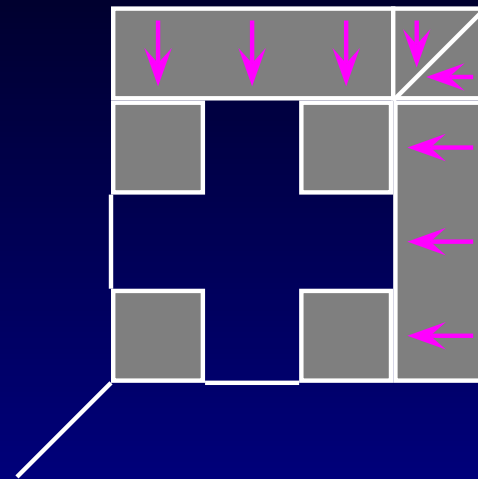
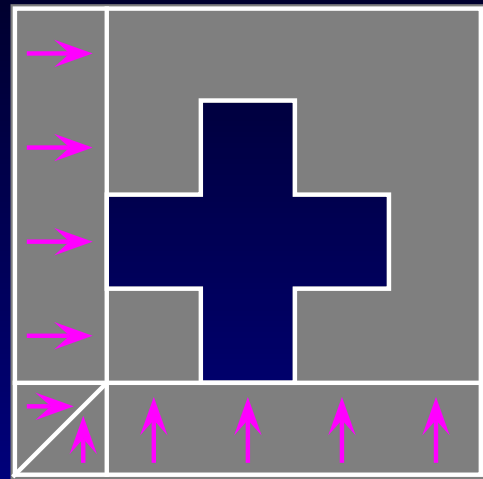
# Sketch of the proof



# Sketch of the proof



# Sketch of the proof



# Directed paths

**Definition:** Let  $x, y \in$  the po-space  $B$ .

- A **dipath** is a dimap  $\vec{I} \rightarrow B$ .
- Dipaths are **dihomotopy equivalent** if they are so in the context of their endpoints.
- Let  $\vec{\pi}_1(B)(x, y)$  be the set of dihomotopy equivalence classes of dipaths from  $x$  to  $y$ .

# Dipaths in equivalent spaces

**Notation:** Given a context  $A$  and po-spaces  $B, C$  together with  $i_B : A \rightarrow B$  and  $i_C : A \rightarrow C$ , if  $x \in A$  denote  $i_B(x)$  by  $x_B$  and  $i_C(x)$  by  $x_C$ .

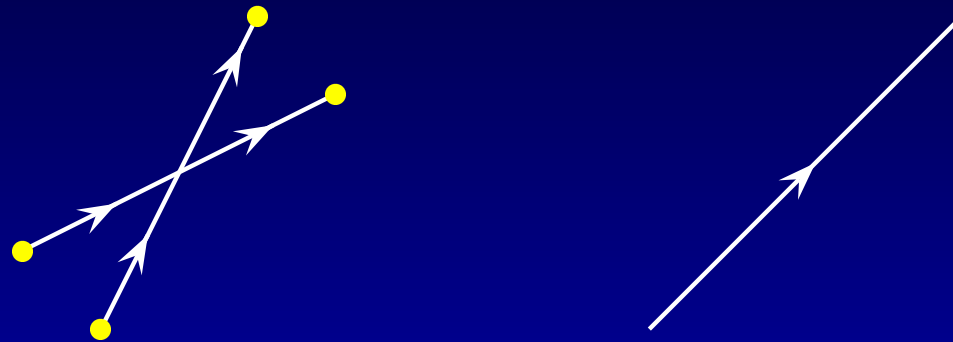
**Proposition:** Given a dimap  $f : B \rightarrow C$  respecting the context and  $x, y \in A$  there is an induced map

$$\vec{\pi}_1(f)(x, y) : \vec{\pi}_1(B)(x_B, y_B) \rightarrow \vec{\pi}_1(C)(x_C, y_C).$$

If  $f$  is a dihomotopy equivalence then it is an isomorphism.

# Example

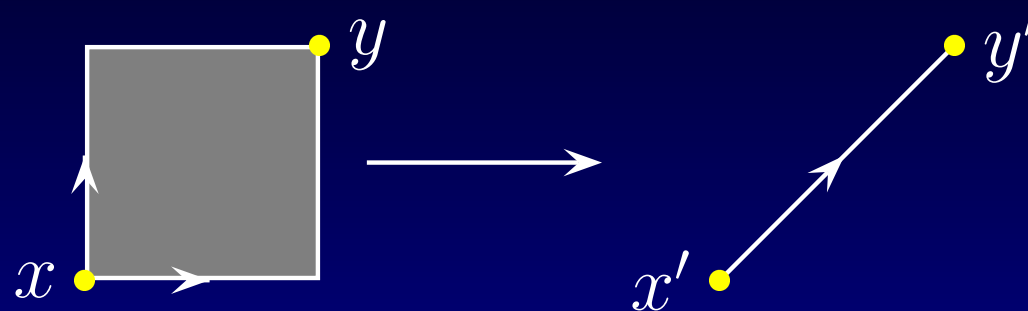
In the context of its four endpoints



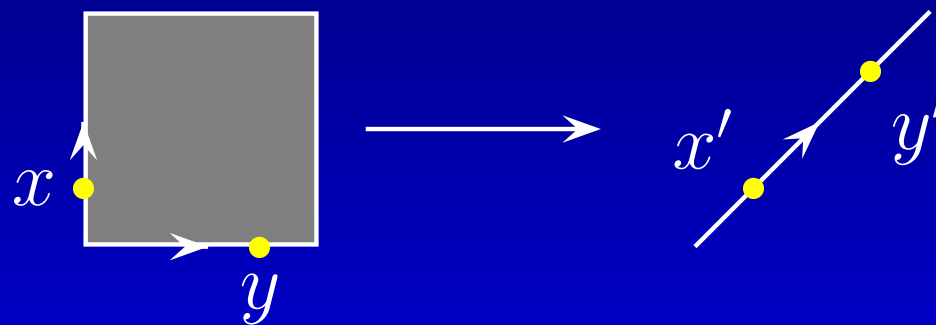
the left hand po-space is not dihomotopy equivalent to  $\vec{I}$ .

# Another example

Recall that there is a dihomotopy equivalence



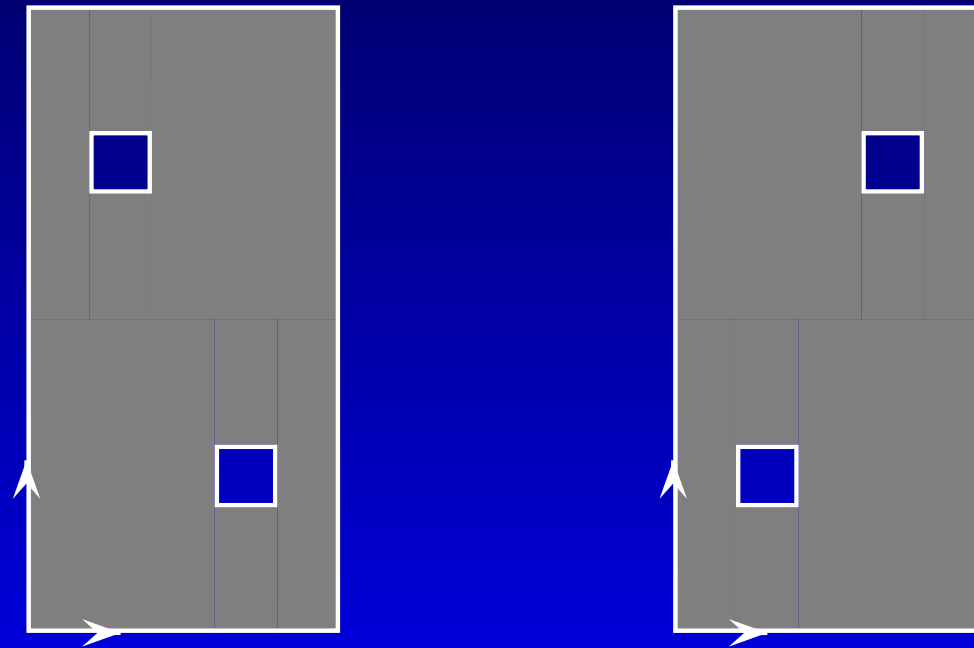
However there is no dihomotopy equivalence



Equivalence depends on the context!

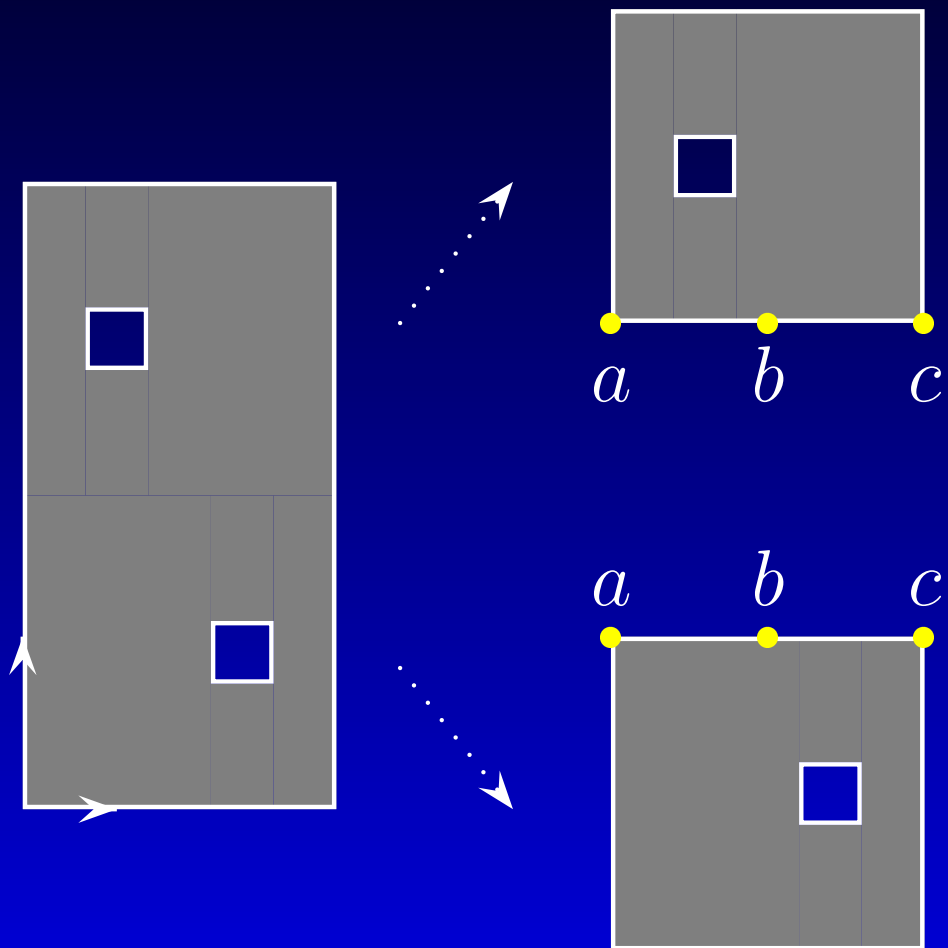
# Compound examples

We would like to find equivalent po-spaces to the following examples by analyzing them piece-by-piece.

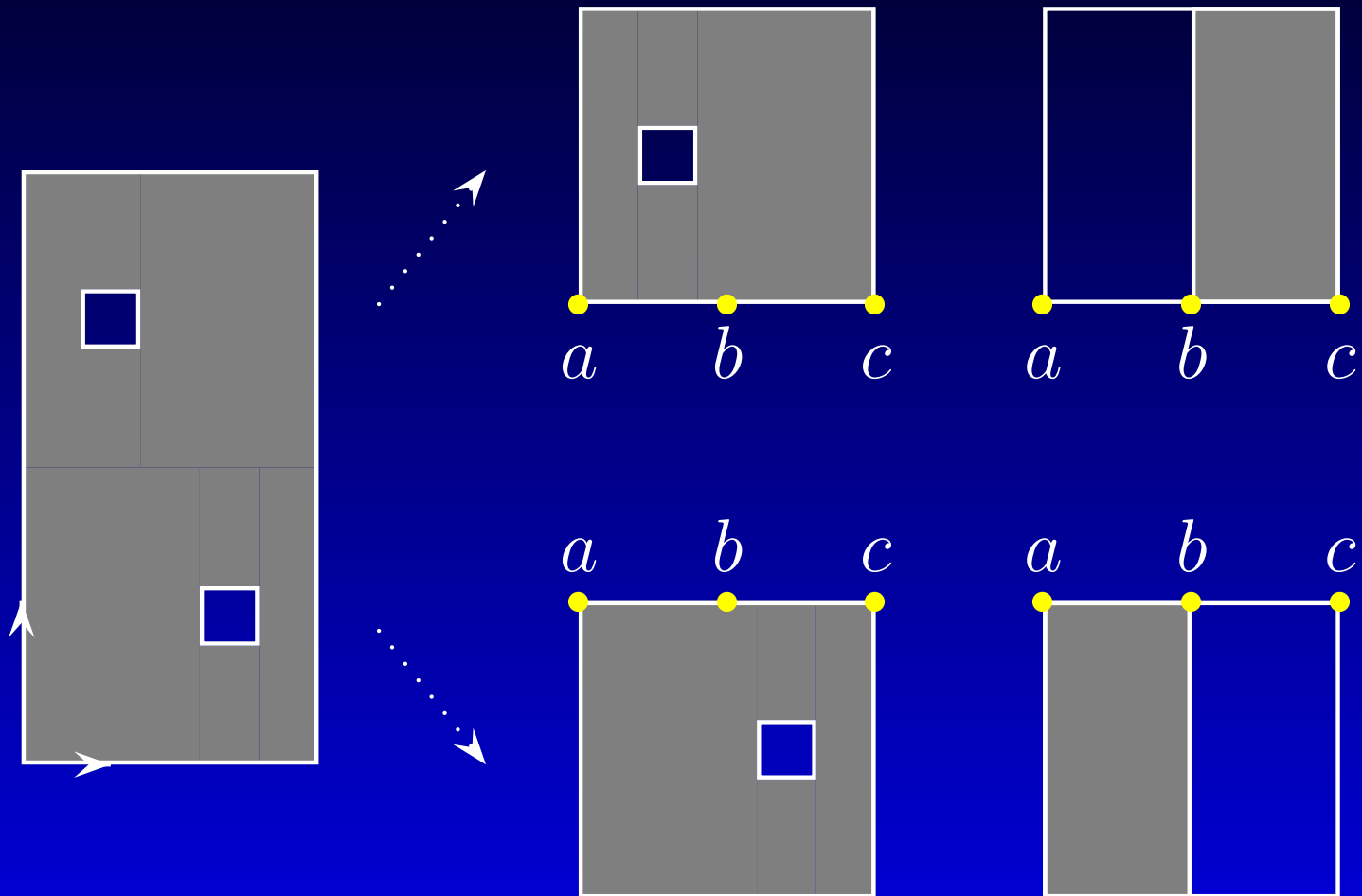




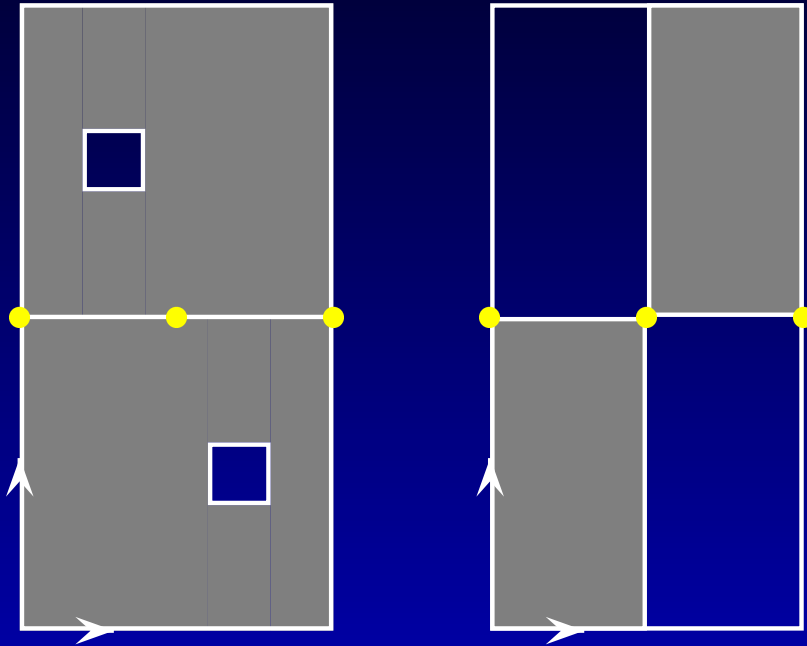
# Equivalences of pieces



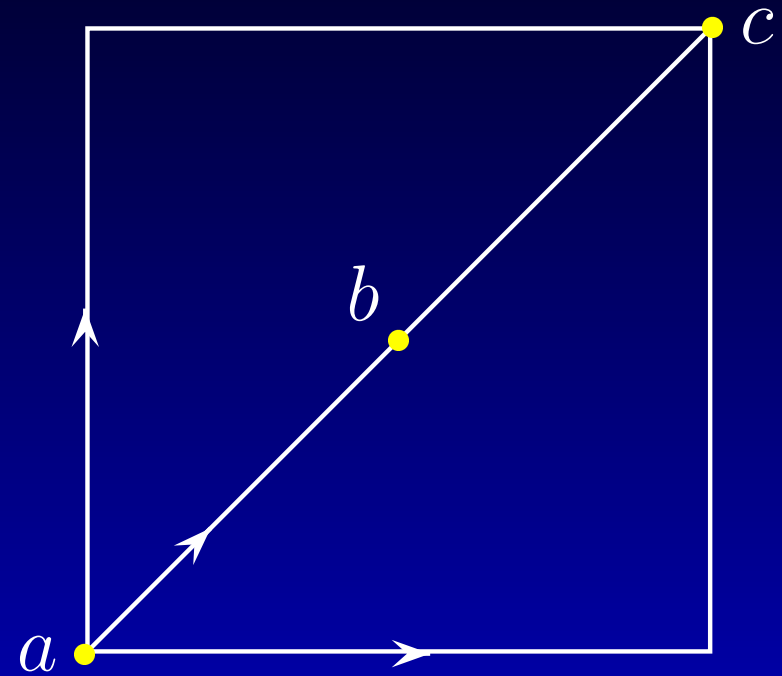
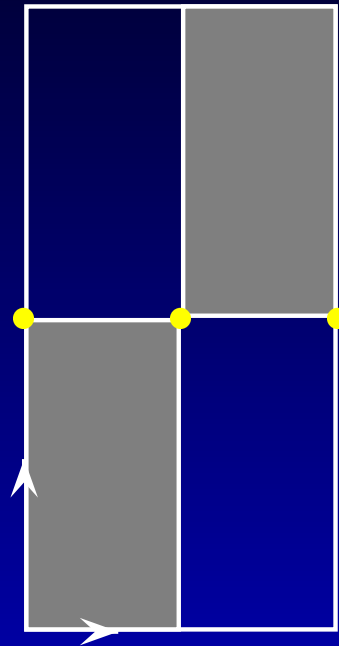
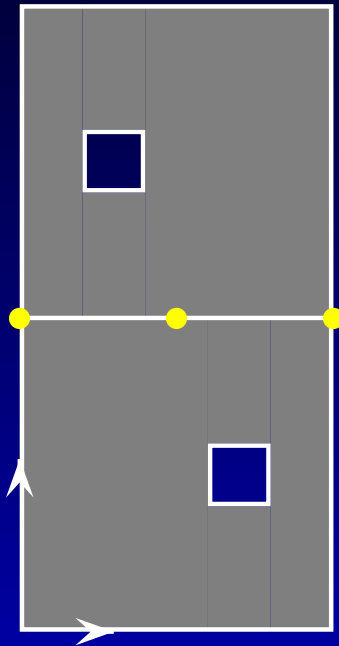
# Equivalences of pieces



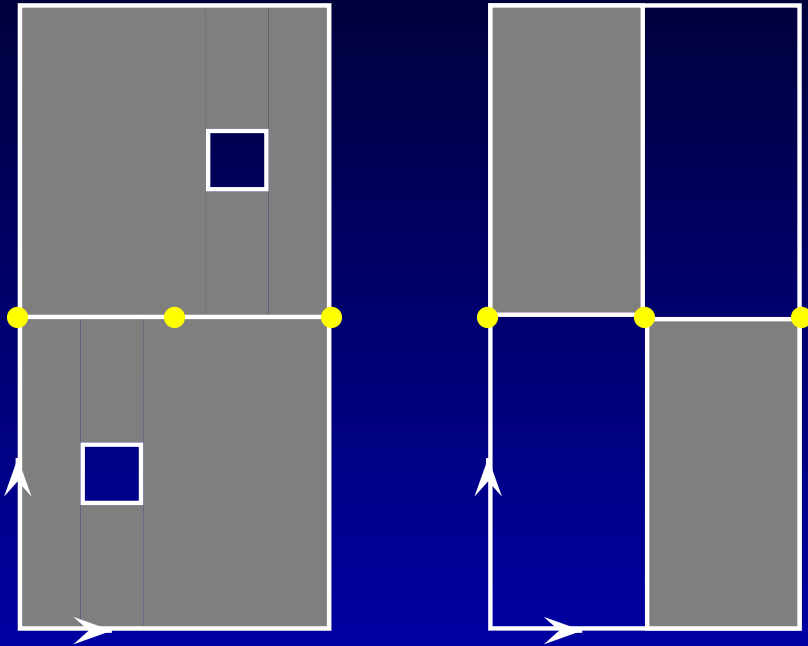
# Patching the pieces together



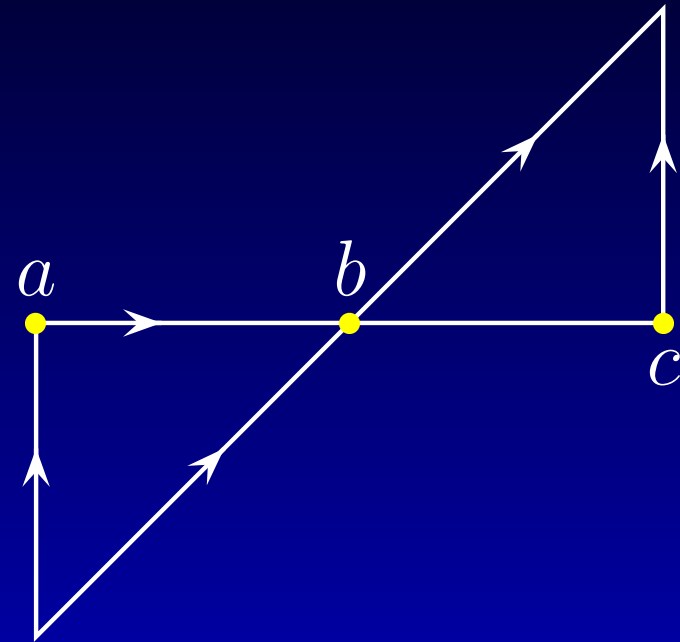
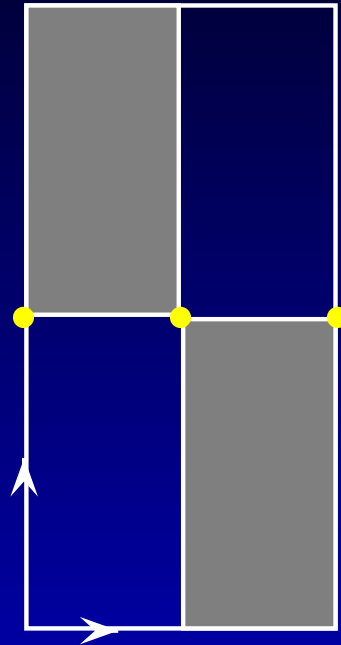
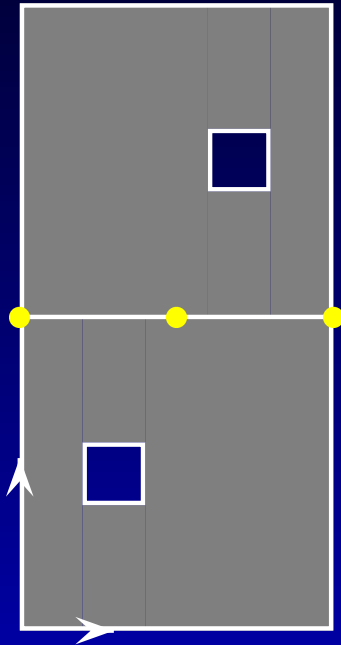
# Patching the pieces together



# Second example



# Second example



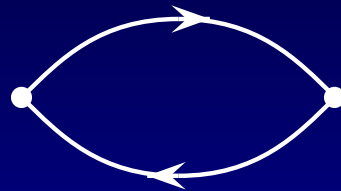
# 3. Modeling systems with loops

(joint work with Kris Worytkiewicz)

- Local po-spaces
- Model categories
- Equivalences for concurrent systems which may have loops

# A more general model

We would like to model execution loops such as



These cannot be modeled by po-spaces.

However they can be modeled by **local po-spaces**.



# Local po-spaces

## Definition:

- An **order atlas** is a open cover of po-spaces with compatible partial orders.
- A **local po-space** is a topological space together with an equivalence class of order atlases.
- A morphism of local po-spaces is a continuous map which respects the orders.

# Equivalences of local po-spaces

Just as with po-spaces, we can define local po-spaces under some context  $A$ , and we consider morphisms which respect the context.

We can also define dihomotopy equivalences using context exactly the same way as with po-spaces.

# Enter some machinery

A powerful framework for studying equivalences is given by **model categories**.

## **Definition:**

A **model category** is a category (with all small limits and colimits) and with three distinguished classes of morphisms: weak equivalences, cofibrations, and fibrations satisfying four simple axioms.

The structure of a model category allows one to apply the machinery of homotopy theory.

# Model category axioms

**M0.**  $\mathcal{C}$  has all small limits and small colimits

**M1.** 2 out of 3

**M2.** retracts

**M3.** lifting property

**M4.** factorization

# A model for concurrent systems

## Theorem [B-Worytkiewicz]:

The category of local po-spaces under a context  $A$  embeds into a model category such that

- the weak equivalences are the dihomotopy equivalences  $\text{rel } A$ ,
- the cofibrations are the monomorphisms, and
- pushouts of weak equivalence with cofibrations are weak equivalences

# Sheaves, Simplicial presheaves

## Definition:

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- There is a Yoneda embedding  $\mathbf{LoPospc} \hookrightarrow \mathbf{Set}^{\mathbf{LoPospc}^{\text{op}}} \hookrightarrow \mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$
- The **sheaves**  $\mathbf{Shv}(\mathbf{LoPospc})$  are the presheaves which are compatible with the topology

# Sketch of the proof

## Theorem[Jardine]:

Let  $\mathbf{C}$  be a small category with a Grothendieck topology. Then  $\mathbf{sSet}^{\mathbf{C}^{op}}$  the category of simplicial presheaves on  $\mathbf{C}$  has a (proper, simplicial) model structure in which

- the cofibrations are the monomorphisms, and
- the weak equivalences are the local weak equivalences.

Furthermore, if the Grothendieck topos  $\mathbf{Shv}(\mathbf{C})$  has enough points then the local weak equivalences are the stalkwise equivalences.

# Grothendieck topology

**Proposition[B-W]:**  $\mathbf{LoPospc}$  has a Grothendieck topology given by **open directed covers**.

# Enough points

Let  $Z$  be a local po-space. **Definition:** The category of directed étale bundles over  $Z$  has

- objects: dimaps  $E \rightarrow Z$  which are local homeomorphisms
- morphisms: maps  $E_1 \rightarrow E_2$  such that

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**Theorem[B-W]:** There is an equivalence between  $\mathbf{Shv}(Z)$  and  $\mathbf{Etale}(Z)$ .

**Corollary:**  $\mathbf{Shv}(\mathbf{LoPospc})$  has enough points.

# Stalkwise equivalences

Using the above results, it follows from Jardine's theorem that  $\mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$  has a model structure in which the weak equivalences are the stalkwise equivalences.

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**Proposition[B-W]:** The stalkwise equivalences coming from  $\mathbf{LoPospc}$  are the isomorphisms.



# Adding context

Given context  $A \in \mathbf{LoPospc}$  there is a induced model structure on

$$A \downarrow \mathbf{sSet}^{\mathbf{LoPospc}^{\text{op}}}$$

Finally one can localize with respect to the dihomotopy equivalences  $\text{rel } A$  to obtain the main theorem.  $\square$

## 4. Summary

- It would be useful to have a robust notion of equivalence in models of concurrency.
- To allow a piece-by-piece analysis we would like equivalences that remain equivalences even after additions are made to the model.
- Using **context** provides such equivalences.

# Summary

- Using po-spaces, local po-spaces, context, and model categories, we have a good mathematical framework for studying concurrent systems.
- In particular, using equivalences, this framework should allow for a piece-by-piece analysis of concurrent systems.

# Future Work

**Theoretical:** Consider all possible contexts in a single framework. (Stay tuned for Kris' talk!)

**Practical:** Use the current theoretical framework for analyzing real-world examples.

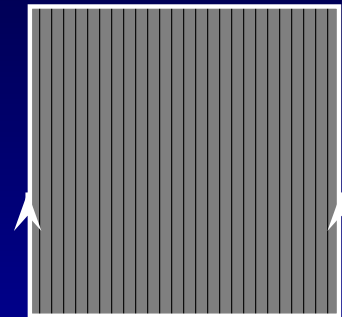
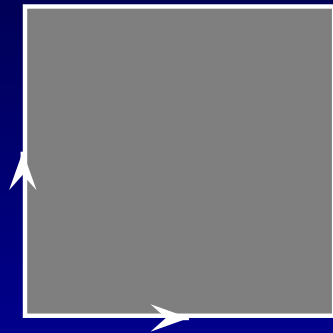
# Acknowledgments

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This talk and associated preprints are available at <http://igat.epfl.ch/bubenik/>

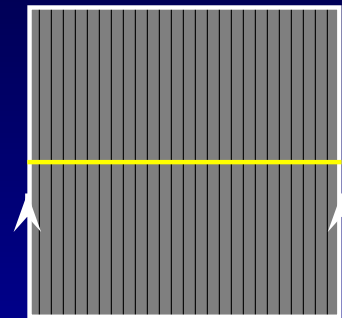
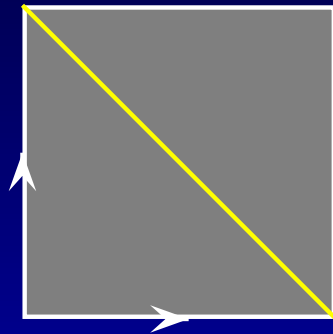
# Non-discrete context

Take  $\vec{I} \times \vec{I}$  and  $I \times \vec{I}$



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and glue them together along the yellow lines.