

CIS 632 / EEC 687

Mobile Computing

TCP in Mobile Networks

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- TCP in Mobile IP Networks*
TCP in Mobile Ad Hoc Networks

Wireless Mobile Transport Layer

- ❑ Wireless environments are characterized by *long latencies* and *frequent interruptions*
 - ❑ Problem: TCP has been optimized for wired networks
 - Wired Network: When a packet is lost, it's typically a sign of congestion → sender should slow down
 - Wireless Network: When a packet gets lost, it could be due to
 - Disconnects
 - Long latencies – slower transmission rates
 - IP tunneling while node moves to new link
- **What can be done? By whom?**

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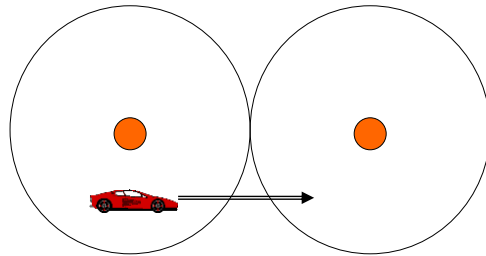
TCP Issues

- ❑ Reducing the transmission rate is often the wrong response over wireless links.
- ❑ The sender should know the network it is transmitting over to make the right decision
- ❑ TCP issues in Mobile IP networks
 - Handoff problem
 - High transmission error rate

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Hand-off Problem

- Hand-offs occur when a mobile host starts communicating with a new base station (in cellular wireless systems)



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Hand-off Problem

- Hand-offs may result in temporary **loss of route** to MH
 - with non-overlapping cells, it may take a while before the mobile host receives a beacon from the new BS
- While routes are being reestablished during handoff,
 - MH and old BS may attempt to send packets to each other, resulting in **loss of packets**

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Hand-off Problem

❑ Packet loss is mistaken as congestion

- Drops the transmission window size
- Slow start to restrict the windows growth rate
- Resets the retransmission timer to a backoff interval

- Thus, reduces the TCP throughput

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Hand-off Problem

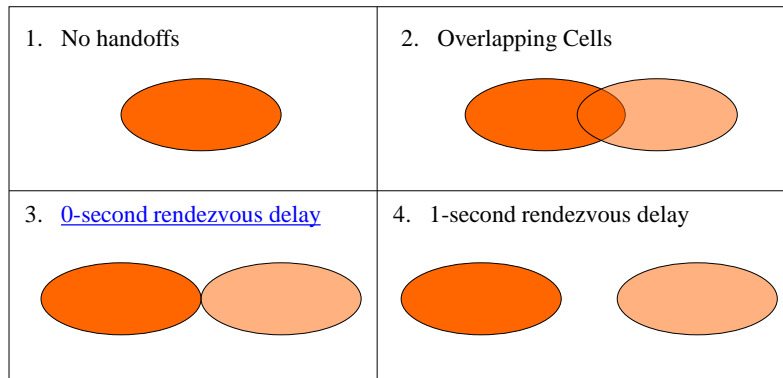
❑ During the long delay for a handoff to complete, a whole window worth of data may be lost

- After handoff is complete, acks are not received by the TCP sender
- Sender eventually times out, and retransmits
- If handoff still not complete, another timeout will occur

❑ Performance penalty

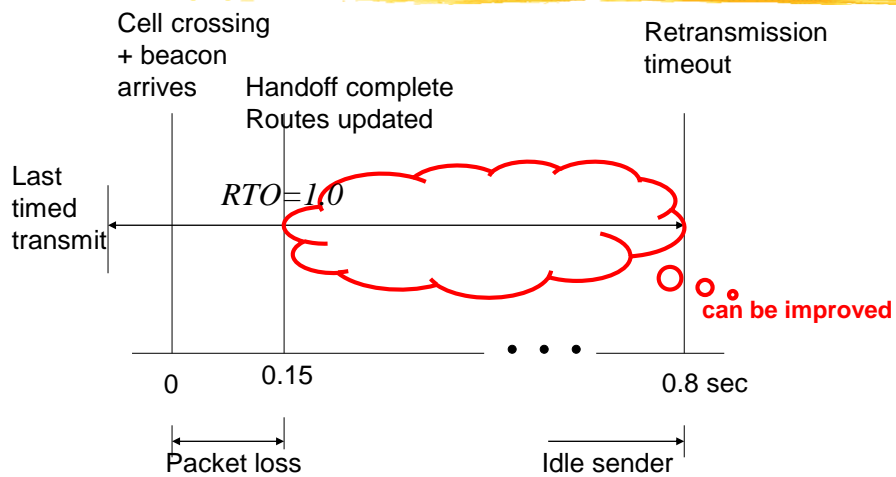
- Time wasted until timeout occurs
- Window shrunk after timeout

Mobility Scenarios



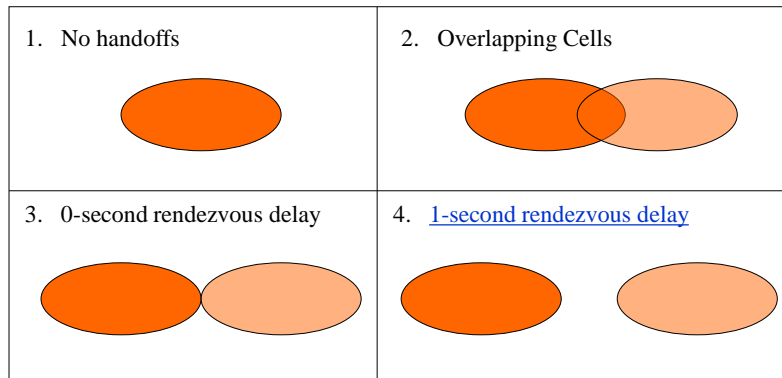
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0-second Rendezvous Delay : Beacon arrives as soon as cell boundary crossed



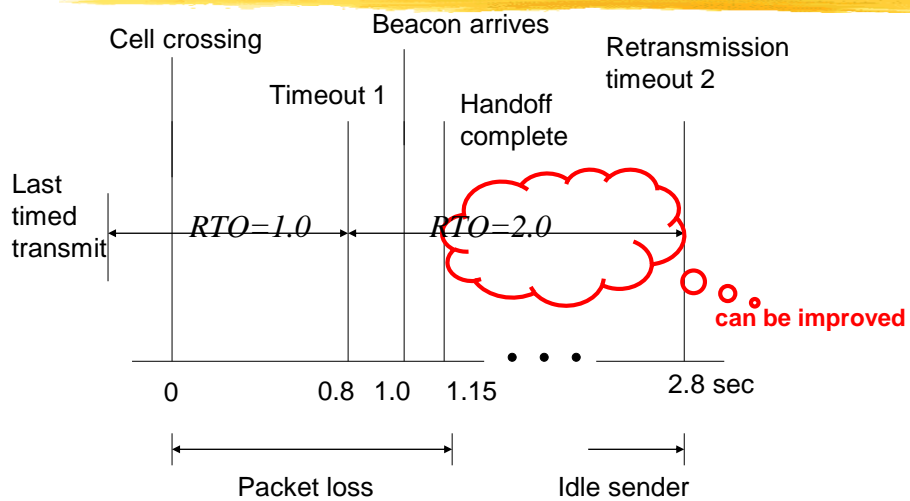
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Mobility Scenarios



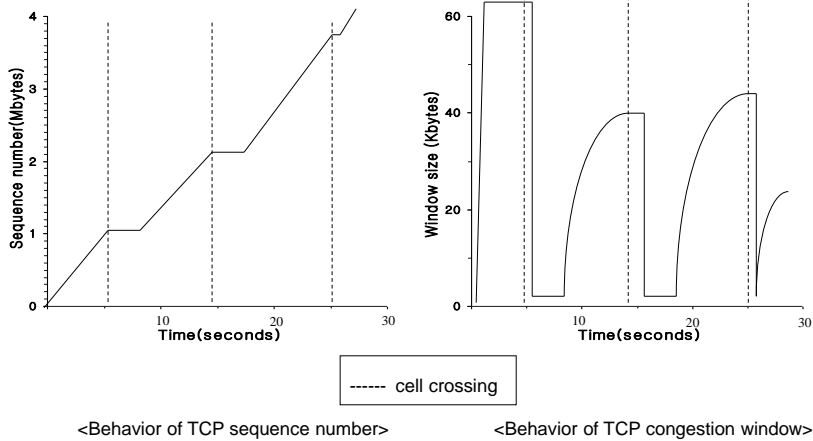
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1-second Rendezvous Delay : Beacon arrives 1 second after cell boundary crossed



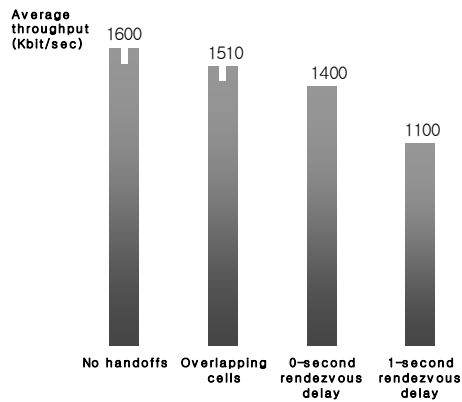
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TCP Performance (1-sec delay)



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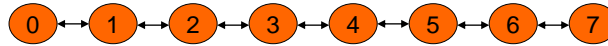
TCP Performance



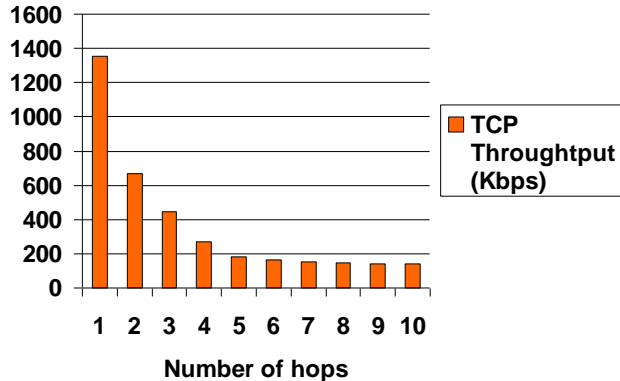
- MH switches cells every 8 seconds
- Throughput dropped significantly in the presence of motion
- Degradation in overlapping cells is due to *encapsulation and forwarding delay* during handoff
- Additional degradation in cases 3 and 4 **due to packet loss and idle time at sender**

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TCP in MANET



Connections over multiple hops are at a disadvantage compared to shorter connections, because they have to contend for wireless access at each hop

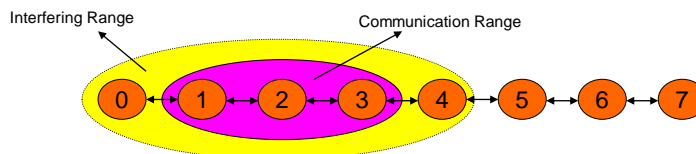


TCP Throughput using 2 Mbps 802.11 MAC
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Why Performance Degrades?

□ MAC Layer - IEEE 802.11 DCF

- Interfering range is a little more than two times of the communication range
- When node 2 transmits, nodes 0, 1, 3 and 4 are blocked



Approaches

- ❑ TCP issues in Mobile IP networks
 - Handoff problem - **Fast retransmission**
 - High transmission error rate - **Split connection**

- ❑ TCP issues in MANET (not today)
 - Impact of multiple-hop route
 - Interplay with 802.11 MAC

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Handoff Problem - Fast Retransmission

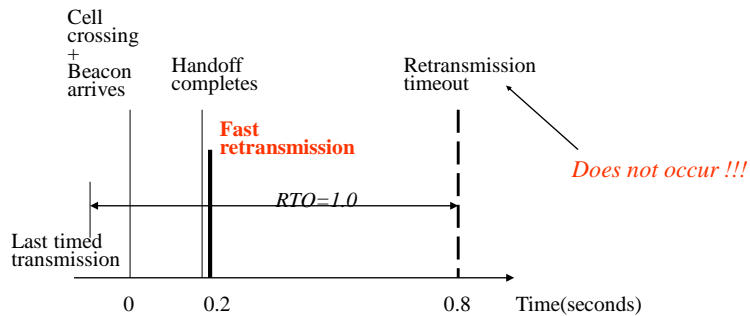
- ❑ When the packet loss is due to handoff, **who** can make the right decision? And **which action** can be taken?

- ❑ When MH is the TCP receiver: **after handoff is complete, it sends 3 dupacks to the sender**
 - this triggers fast retransmit at the sender
 - instead of dupacks, a special notification could also be sent

- ❑ When MH is the TCP sender: **invoke fast retransmit after completion of handoff**

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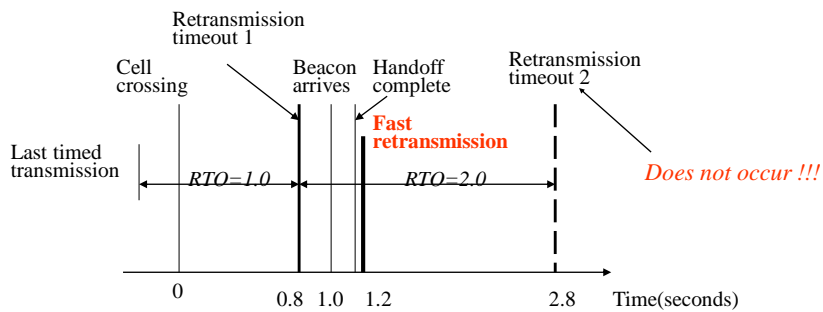
0-second Rendezvous Delay Improvement using Fast Retransmit



Fast retransmission after a handoff with a 0-second rendezvous delay

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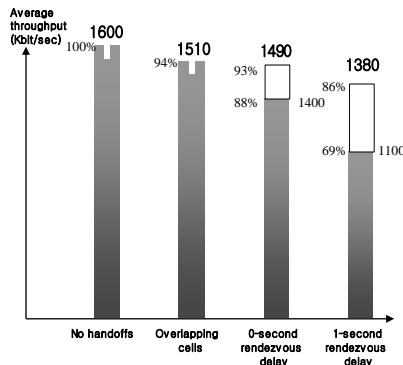
1-second Rendezvous Delay Improvement using Fast Retransmit



Handoff latency and related packet losses with a 1-second rendezvous delay

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TCP Performance Improvement



- ❑ No change in the first two cases as expected
- ❑ Improvement for non-overlapping cells
 - Some degradation still remains
 - fast retransmit reduces congestion window

❑ Do we need to change TCP software ?

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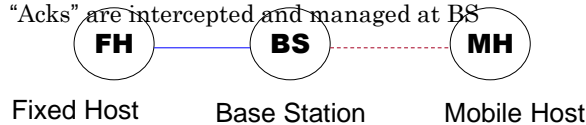
Approaches

- ❑ TCP issues in Mobile IP networks
 - Handoff problem - **Fast retransmission**
 - High transmission error rate - **Split connection**
- ❑ TCP issues in MANET
 - Impact of multiple-hop route
 - Interplay with 802.11 MAC

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High Transmission Error Rate - Split Connection Approach

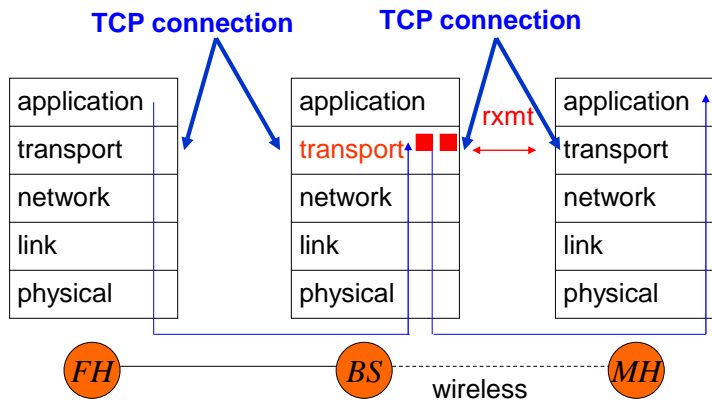
- ❑ How can we address the problem of high error rate over wireless links?
- ❑ End-to-end TCP connection is broken into one connection on the wired part of the route and one over wireless part of the route
- ❑ A single TCP connection split into **two TCP connections**
 - $FH-MH = FH-BS + BS-MH$
 - "Acks" are intercepted and managed at BS



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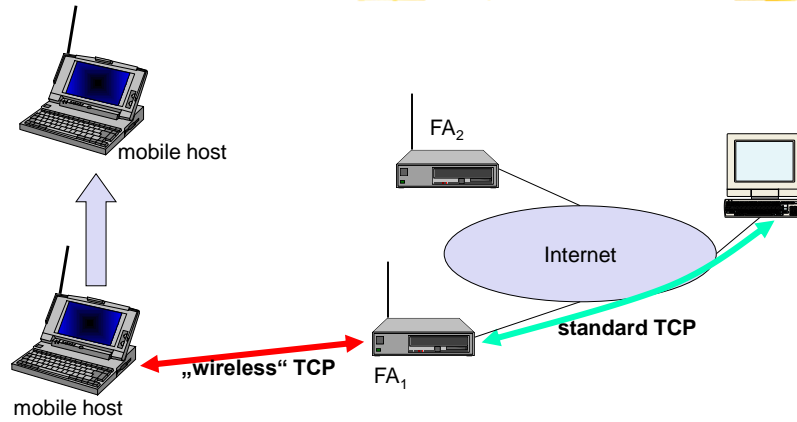
Split Connection Approach

■ Per-TCP connection state



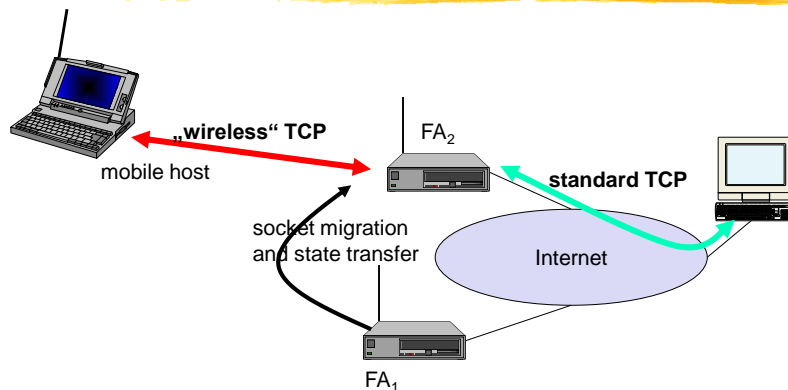
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Split Connection Approach: What happen if move



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Split Connection Approach: State Migration



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Split Connection Approach

- ❑ BS terminates the standard TCP connection acting as a proxy
- ❑ Old BS (FA) must migrate buffered packets (already acknowledged to FH) as well as **socket of the proxy** to new BS
- ❑ The socket contains the current state of the TCP connection
 - Sequence number, addresses, port number
 - Last packet transmitted to MH
 - Last packet acknowledged by MH
 - Next expected acknowledgement and expected number of duplicated acknowledgements
 - Round-trip time estimate of wireless link

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Split Connection Approach: Variations

- ❑ Indirect TCP
 - FH - BS connection : Standard TCP
 - BS - MH connection : Standard TCP
- ❑ Selective Repeat Protocol (SRP)
 - FH - BS connection : standard TCP
 - BS - MH connection : selective repeat protocol on top of UDP
- ❑ Asymmetric transport protocol (Mobile-TCP)
 - Low overhead protocol at wireless hosts such as header compression, simpler flow control, No congestion control
- ❑ Mobile-End Transport Protocol
 - BS-MH link can use any arbitrary protocol optimized for wireless link

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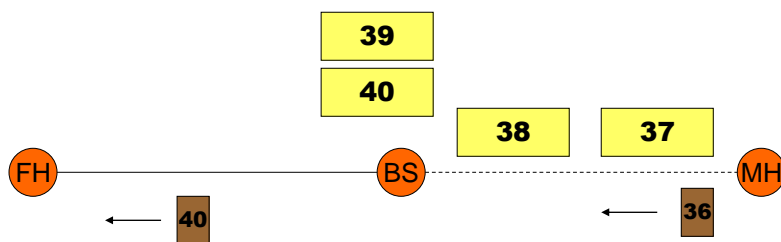
Split Connection Approach : Advantages

- ❑ BS-MH connection can be **optimized** independent of FH-BS connection
 - Different flow / error control on the two connections
- ❑ **Local recovery** of errors
 - **Faster** recovery due to relatively shorter RTT on wireless link
- ❑ **Good performance** achievable using **appropriate** BS-MH protocol
 - Standard TCP on BS-MH performs poorly when multiple packet losses occur per window (timeouts can occur on the BS-MH connection, stalling during the timeout interval)
 - **Selective acks improve performance for such cases**

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Split Connection Approach : Disadvantages

- ❑ End-to-end semantics violated
 - ack may be delivered to sender, before data delivered to the receiver



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Split Connection Approach : Disadvantages

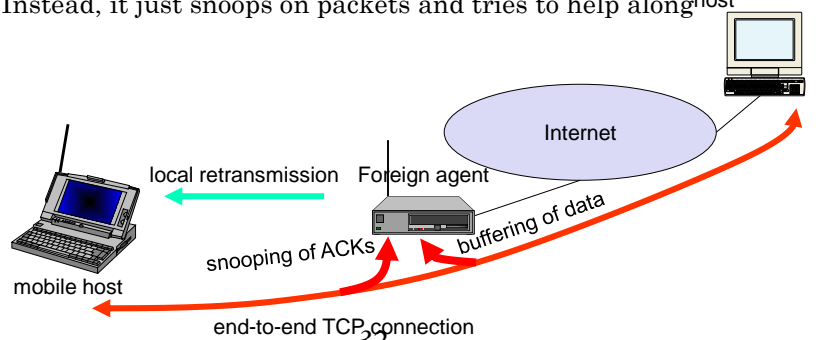
- ❑ BS retains hard state
 - BS failure can result in loss of data (unreliability)
 - If BS fails, packet 40 will be lost
 - Since it is ack'd to sender, the sender does not buffer 40
 - Hand-off latency increases due to state transfer
 - Data that has been ack'd to sender, must be moved to new base station
- ❑ Buffer space needed at BS for each TCP connection
 - BS buffers tend to get full, when wireless link slower (one window worth of data on wired connection could be stored at the base station, for each split connection)
- ❑ Extra copying of data at BS
 - copying from FH-BS socket buffer to BS-MH socket buffer (at TCP layer)
 - increases end-to-end latency

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Snoop Protocol: TCP-Aware Link Layer

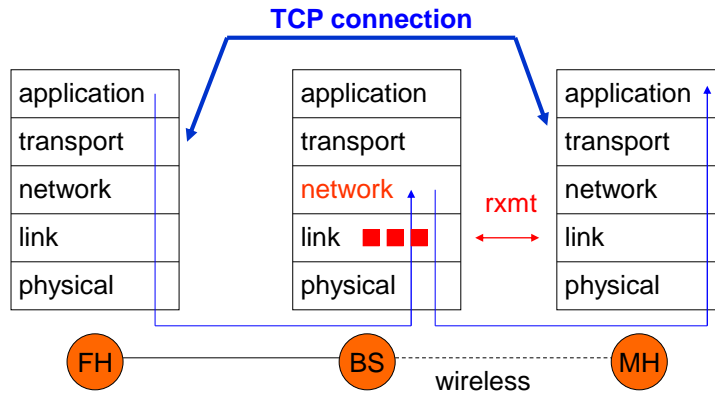
Tries to restore the end-to-end semantics of TCP:

- ❑ Foreign agent is not allowed to send acknowledgements
- ❑ Instead, it just snoops on packets and tries to help along



Snoop Protocol: TCP-Aware Link Layer

■ Per TCP-connection state



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